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JOHN WILSON

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THE CODE**
on the
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John Wilson

Cracking the code on the Sinclair ZX Spectrum

Pan Books London and Sydney

First published 1984 by Pan Books Ltd,
Cavaye Place, London SW10 9PG
in association with Personal Computer News

9 8 7 6 5 4 3 2 1

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ISBN 0 330 28665 X

Photoset by Parker Typesetting Service, Leicester
Printed and bound in Great Britain by
Richard Clay (The Chaucer Press) Ltd, Bungay, Suffolk

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Introduction

This book is intended for people with a reasonable knowledge of Sinclair BASIC and the Spectrum computer who now want to become proficient in machine code programming.

Machine code is like BASIC in that it is a language for communication with the computer, but it differs in the way that precise instructions have to be given in order to perform even the most simple of calculations and operations. These machine code instructions and their use will be introduced and explained with the aid of example programs, so that by the end of the book the reader and user (a lot depends on practice!) will be a competent machine code programmer.

The book starts by introducing the programmer to number representation and goes on to simple loading and storing techniques. It then proceeds through each set of instructions clearly and methodically, with plenty of examples.

After the explanation of the instruction set the use of a monitor is introduced and a full machine code monitor listing, which can be utilised to enter other machine code routines in this book, is provided. This is followed by a detailed breakdown of a machine code program.

Once these chapters have been digested the programmer can progress to dealing with more complex techniques. These involve using the ROM routines, screen handling, interrupts, and include a routine to handle sprites. Finally, we have a chapter which includes some useful machine code routines to enhance your own programs such as a sort, music, and pixel scroll routines. In the appendices are complete listings by op code and Mnemonic of the Z80 instruction set.

1 Chips, registers and numbers

The Spectrum's central processing unit (CPU) or main control chip is known as the Z80. This powerful little chip handles all the additions, subtractions and logical operations with which the Spectrum implements your BASIC and machine code programs. To communicate to this chip when calculations and operations need to be done the user can of course type instructions in BASIC. However, there are other languages which can be used. The fastest of these is machine code which acts directly on the Z80 chip and can be very efficient. Machine code consists of a set of simple instructions which the Z80 CPU understands and can execute, such as addition, subtraction and comparison. This particular chip has over 700 instructions that can be sorted into a collection of a few different types. These instructions act upon data in the form of memory addresses and numbers.

BASIC is a very easy language in which to program due to the fact that we write a line of BASIC almost as we would say it in English, so that;

```
LET X=X+20*2+1
```

means set the variable x equal to the correct value of x, plus twenty times two, plus one. In machine code programming, however, we have to give more precise instructions at a low level, and specify each individual operation needed to perform the calculation.

The example above could be broken down to the sequence:

'add 20 to itself'	(2*20)
'add 1 to that result'	(2*20+1)
'add x to that result'	(X+2*20+1)

'and put the answer back in x'

It should be noted that the above is not an example of machine code instructions but simply illustrates the precision with which machine code operations have to be specified. Why should programmers use this complex sequence of machine code instructions when BASIC is

so easy? Let us look at an example to answer this question. First type in this BASIC program and RUN it:

```
10 FOR X=16384 TO 22527
20 POKE X,255
30 NEXT X
```

When RUN, the program very slowly fills the screen with ink. Now try running an equivalent machine code program:

```
1 CLEAR 31999
10 FOR x=32000 TO 32014
20 READ a: POKE x,a
30 NEXT x
35 RANDOMIZE USR 32000
40 DATA 33,0,64,1,0,24,54,255,
35,11,120,177,32,248,201
```

This program POKES a sequence of machine code instructions into the Spectrum RAM. The DATA at line 60 is the machine code program equivalent to the BASIC version given above. Each number represents a certain instruction which the computer's 'brain', the 'Z80 chip, executes. (Don't try to understand the code yet, just type it in!)

RUN the program . . . but don't blink, otherwise you will miss what happens! As you can see from the example, machine code is incredibly fast. An efficient machine code program can execute up to 1000 times as fast as the BASIC equivalent. What's more, machine code is also compact. You can write machine code routines which occupy only a quarter of the memory that their BASIC counterparts would.

Why is BASIC so slow? Well, the reason lies in the fact that the Z80 chip (which does all the calculations for the Spectrum) can only understand machine code. In order for it to execute a BASIC program, it first has to look up each BASIC keyword or token every time it reads a line. It then takes this token and translates or interprets it to specify the equivalent ROM machine code routine so that it can then perform the operation. This all takes time. Machine code, however, is the Z80's 'Mother Tongue', so no translation is needed and the code is executed immediately.

Hexadecimal and binary

All of you should know that the Spectrum (or any other computer for that matter) stores data in terms of 'bytes'. A byte is an 8 bit binary number which can have a decimal value of 0 to 255. In a 48K

Spectrum there are 49152 locations in memory where bytes can be stored. The value 49152 is obtained by the calculation 48×1024 because $1K = 1024$ bytes.

The Z80 chip stores numbers in groups of 8 bits, so it is known as an '8 bit chip'. In this it is similar to the 6502 chip which is used in the BBC Micro, Oric and Commodore machines. Other microprocessor chips use 16 or 32 bits and are therefore known as '16 bit' or '32 bit' chips.

To address RAM the Z80 chip uses 2 bytes (or 16 bits) This means that it can access 65536 characters, since the number of combinations of 16 1's and 0's is 65536. These bits and how they represent numbers and characters are best explained by looking at the system known as the *binary system* (or 'base two system').

In the real world of handling money we count in a system known as *decimal* or 'base 10 system'. We have the digits 0,1,2,3,4,5,6,7,8 and 9 which we can write to represent certain quantities.

In the decimal system we can break down the number we are using into groups of powers of ten. That is units, tens, hundreds, thousands, ten thousands, and soon. For example, the number 3456 can be broken down to:

$$\begin{array}{r} 3 \times 1000 \\ + 4 \times 100 \\ + 5 \times 10 \\ + 6 \times 1 \end{array} \qquad \begin{array}{r} (3 \times 10 \uparrow 3) \\ (4 \times 10 \uparrow 2) \\ (5 \times 10 \uparrow 1) \\ (6 \times 10 \uparrow 0) \end{array}$$

In the binary system we use only two digits, these being 0 and 1. In order to represent large numbers therefore we can only write in a series of these two digits.

Remember that the Z80 chip represents information (numbers) in groups of 8 bits. Each of these bits may be 'off' (i.e. digit 0) or 'on' (i.e. digit 1). The bits in a byte are numbered 0 to 7, starting from the right.

In the binary system numbers are broken down in powers of two (that's why it is also known as the base two system). That is to say we break them down as factors of units (bit 0), two's (bit 1), four's (bit 2), eight (bit 3), sixteen (bit 4), thirty-two (bit 5), sixty-four (bit 6) and one hundred and twenty eight (bit 7).

Take for example the binary number 00011001, this represents the decimal number:

$0*128$	$(0*2 \uparrow 7)$
$+0*64$	$(0*2 \uparrow 6)$
$+0*32$	$(0*2 \uparrow 5)$
$+1*16$	$(1*2 \uparrow 4)$
$+1*8$	$(1*2 \uparrow 3)$
$+0*4$	$(0*2 \uparrow 2)$
$+0*2$	$(0*2 \uparrow 1)$
$+1*1$	$(1*2 \uparrow 0)$

25 decimal

The maximum number that can be represented in 8 bit (one byte) binary form is therefore 11111111, which represents 255 in decimal ($128+64+32+16+8+4+2+1$).

In order to deal with larger numbers the Z80 has some 16 bit instructions. All memory addressing is done with 16 bits, so the total number of individual bytes that can be pointed to in memory (*addressed*) should be equal to the total number of combinations of a 16 digit binary number. This will be equal to the maximum value +1 (since the value zero is a unique combination).

To obtain the maximum value possible in a 16 digit binary number we must evaluate 1111111111111111. This has a value of:

$1*32768$	$(1*2 \uparrow 15)$
$+1*16384$	$(1*2 \uparrow 14)$
$+1*8192$	$(1*2 \uparrow 13)$
$+1*4096$	$(1*2 \uparrow 12)$
$+1*2048$	$(1*2 \uparrow 11)$
$+1*1024$	$(1*2 \uparrow 10)$
$+1*512$	$(1*2 \uparrow 9)$
$+1*256$	$(1*2 \uparrow 8)$
$+1*128$	$(1*2 \uparrow 7)$
$+1*64$	$(1*2 \uparrow 6)$
$+1*32$	$(2*2 \uparrow 5)$
$+1*16$	$(1*2 \uparrow 4)$
$+1*8$	$(1*2 \uparrow 3)$
$+1*4$	$(1*2 \uparrow 2)$
$+1*2$	$(1*2 \uparrow 1)$
$+1*1$	$(1*2 \uparrow 0)$
$+1$	

65536
OR 64K (1K=1024 bytes)

number distinguishes between hex and decimal numbers) so that 20_H is clearly different from 20 decimal. This means $2*16+0*1=32+0=32$ decimal.

The number AF_H means $16*A+F*1$. Since A in hexadecimal notation is 10 decimal, and F is 15, this gives $16*10+15*1=160+15=175$ decimal.

The hexadecimal notation is widely used by machine code programmers since it makes numbers easier to remember than binary but more significant than decimal. Because the hexadecimal system is based on 16 (10_H), and 16 is 10000 in binary, there is a close relationship between the binary system that the Z80 chip uses and the hexadecimal notation that most programmers use. Unfortunately there is no simple relationship between the decimal and binary systems, as the table below should illustrate:

BINARY	DECIMAL	HEXADECIMAL
10101011	171	AB
00010010	18	12
10000001	129	81
11110000	240	F0

8 bit hexadecimal values have up to two digits. These each represent the value in one *nybble* of the byte. A nybble consists of four bits, either the leftmost four or the rightmost four and by taking the value of each nybble the Hexadecimal digit can be calculated. In the example above the binary value 10101011 is shown to have a hexadecimal value of AB. This can be illustrated by taking the high nybble (1010) which equals 10 decimal (A in hex) and the low nybble (1011) which equals 11 decimal (B in hex), then combining them in the same order to give AB_H.

So far we have only seen machine code entered by *POKE*ing numbers into memory. This method of writing machine code is tedious and makes it difficult to understand and debug the code, so the designers of the Z80 chip developed a standard set of *mnemonics* in which to write Z80 code.

These Mnemonics are English-like words which (hopefully!) signify the action a particular instruction performs. For example, the mnemonic RET means RETURN and is equivalent to the RETURN instruction in BASIC, ie. it tells the processor to continue with the main program after a subroutine was called.

In order to translate these mnemonics into data which the computer understands we will need to assemble them. This can be done by hand but more often by a utility known as an *assembler*. The

programmer first of all types in a program in standard mnemonics and then the program assembles these instructions into machine language. Most machine code programmers write assembly code and use an assembler to create their machine code

When an assembler translates the `RET` instruction it puts into memory the value for that instruction, which is 201 decimal, C9 hexadecimal or 11001001 in binary.

There are plenty of good assemblers for the Spectrum on the market ranging in price from around £7 to £14. Most of these will work on both the 16K and 48K models. The 'Devpac' package from Hi-Soft, as an example, is at the top of this price range, but is good value. In addition to the assembler it comes with another package known as a *monitor*. Alternatively, Chapter 7 provides you with your own monitor program for only the cost of wear and tear on the fingertips. This is a utility which will allow you to enter and experiment with the routines in this book.

A monitor program allows the machine code programmer to input and look at a program in hexadecimal form. Other features often included with it are utilities to set break points, look at the values held in the registers (the Z80 'variables') and to move, save and load blocks of memory. Both the Spectrum monitor provided in this book and the Devpac monitor have all these standard features. Devpac's also includes the capacity to move a single step at a time through a machine code program. There is also a *disassembler* in the Devpac package. This is a routine which is the opposite of an assembler for it converts machine code binary data into Z80 mnemonics.

When seeking an assembler for your Spectrum you are advised to buy one which allows you to assemble a program at different addresses in memory. Most assemblers have a command `ORC` (ORIGIN) which tells the assembler the start address from which to assemble the program. This is illustrated in the assembler listings included in this book.

There are certain features of assembler listings that need to be explained here otherwise confusion may occur. Assemblers have a feature which enables them to use what are known as '*pseudo*' operators. These are used to place strings or numbers in memory and are not standard Z80 mnemonics. They are only a feature used in certain assemblers, including the one used for the listings in this book.

DEFB Define Byte

Can sometimes be abbreviated to 'DB'. This places the following data in memory. For example:

```
DB 02H,04H
```

would place the number 2 followed by a 4 at the location where it is being assembled.

DEFW Define Word

This is similar to DEFB but is used to place a two byte number in memory. The low byte of the given number is placed in location where it is assembled. The high byte will follow, as we explained earlier when 16 bit values were introduced:

```
DEFW 7 (equivalent to DEFW 0007H)
```

is the same as:

```
DB 0,7 (equivalent to DB 00H,07H)
```

DEFS Define Space

The number following this Pseudo operator is the number of bytes which we want to reserve. So the operator:

```
DEFS 100
```

Would reserve 100 bytes.

EQU Equate

This instruction is used to give values to labels. The format is a label, followed by the EQU, followed by a number:

```
PLOT EQU 22E5H
```

The above would give the label PLOT the value 22E5 hex.

; Comment

In most assemblers the ; is used in the same manner as the BASIC REM to indicate a useful remark or comment. This is very useful because without helpful comments assemble code is harder to understand than BASIC because the operations are less immediately obvious.

Another feature of machine code assemblers is the facility to refer to memory addresses by means of labels. Instead of entering an instruction which says 'Jump to Address 31000', we can set a label at the address 31000. We could assign the label the name 'Fred', for example, and then give an instruction 'Jump to Address Fred'. This can greatly simplify our program structure and also enables meaningful label names to be assigned to sections of code.

If you use the appendices of this book you will be able to assemble your own machine code programs. The first thing you need to do is to write the assembly code (Mnemonics) for your program. I have provided an example below which will go into the printer buffer to avoid you having to CLEAR high memory space:

```

ORG 23296 ; Start the code at the printer buffer
LD HL,4000H
LD DE,4001H
LD BC,17FFH
LD (HL),0
LDIR
RET

```

The effect of this program is to remove all the ink from the screen. How it does so is not important currently because it is serving only to demonstrate how you can get machine code to work without buying an assembler program.

The ORG is not a part of the machine code but it shows where in memory the machine code must be stored. This is the address into which we will start to POKE the data.

To obtain the data for each of the mnemonics above you will need to look them up in Appendix 2. As an example, the entry for LD HL, 4000H will read:

Mnemonic	Decimal	Hex
LD HL,XXXX	33 XXX XXX	21 XX XX

In order to get the hex for LD HL,4000H the 4000H must be converted into two bytes and reversed in order (due to the LSB/MSB storage convention explained earlier).

So, LD HL,4000H will assemble to 21 00 40 in hex or 33 0 64 in decimal. Since we will be using a BASIC program to POKE the code you will need to calculate the decimal values to be placed in the data statement. I have calculated the example for you but try to follow through the procedure to make sure you understand the principles involved.

HEX	DEC	MNEMONICS
		ORG 23296
21 00 40	33 0 64	LD HL,4000H
11 01 40	17 1 64	LD DE,4001H
01 FF 17	1 255 23	LDBC,17FFH
36 00	54 0	LD (HL),0
ED B0	237 176	LDIR
D9	201	RET

Now to enter this machine code program the following BASIC program could be used:

```

10 FOR I=0 TO 703:PRINT CHR$(32+INT (128*RND));:NEXT I
20 LET A=23296
30 READ B:IF B=-1 THEN GOTO 50
40 POKE A,B:LET A=A+1:GOTO 30
50 PRINT #0: "PRESS A KEY TO CLEAR":PAUSE 1:PAUSE 0
60 RANDOMIZE USR 23296
70 DATA 33,0,64,17,1,64,1,255,23,54,0,237,176,201,-1

```

As you can probably see, this would be a reasonable way to write small programs of up to about 100 bytes but to write your first full machine code 48K mega-game you will need an assembler to shorten the development time. Another considerable advantage of using an assembler program is that you can save the source code (assembly code or mnemonics). It than can be loaded back from tape or microdrive and errors can be corrected in the machine or object code.

Registers

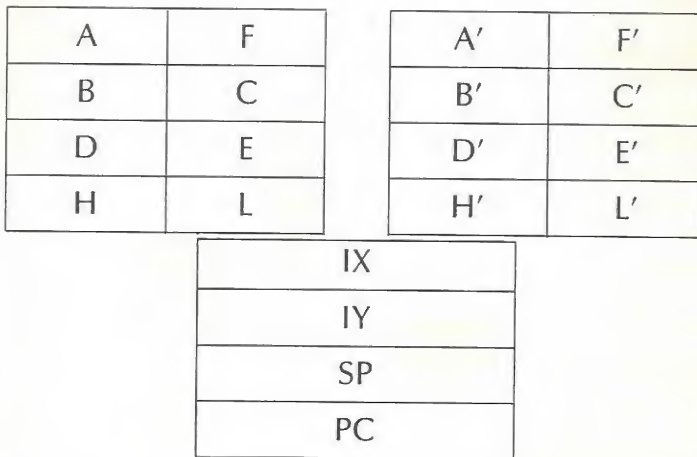
The Z80 CPU has several registers available to the programmer. These can be used to hold numeric values similar to BASIC variables but the programmer is limited to 22 registers. Some of these registers can be used in pairs to hold 16 bit values. The older chips such as the 6502 are unable to do this. The Z80 registers are referenced by the names:

A,B,C,D,E,F,H,L,IX,IY,SP,PC,I and R

From these A, B, C, D, E, H, L can all hold 8 bit values and IX, IY, SP and PC will hold 16 bit values. Registers SP, PC, F, I and R have specific functions which will be explained later and are not used for holding user data. In addition to these there is a second set of A, B, C, D, E, F, H

and L registers which are usually referred to as A', B', C', D', E', F', H' and L'. These two sets of registers cannot be used at the same time, so in order to access the alternate set a special instruction has to be used, 'EXX' (Exchange) which flips from one register set to the other. (Two exceptions here are the A' and F' registers which are exchanged using EX AF, AF')

As mentioned earlier, some of these 8 bit registers can be paired off to form one sixteen bit register. The diagram below demonstrates how this can be done:



Now let's take a more detailed look at each of the registers that we have just introduced and their functions.

IX and IY Registers

These are known as the Index registers. The IX and IY registers are often used to point to tables of data and are extremely powerful tools for accessing arrays of data by a method known as indexing. On the Spectrum great care must be taken before using the IY index register in your own machine code programs. A number of ROM routines require that IY contains the value 5C3AH (23610 decimal) otherwise they will not work correctly. The Interrupt routine also requires this value to be in IY. Therefore if you must use IY in your machine code, disable the interrupts and make sure that IY=5C3AH before calling any ROM subroutines or returning to BASIC. Disabling interrupts and ROM subroutines are dealt with later in this book.

I and R Registers

The I or Interrupt register is used in conjunction with a technique known as *vectored interrupt programming*. This is the Z80's pointer

for alternative interrupt routines and is described in detail in Chapter 11.

The R or Refresh register is used to refresh any dynamic ram connected to the Z80. The only purpose it serves for Spectrum programmers is its use in random number generation.

A Register

The A register is known as the Accumulator and is the main register for performing 8 bit arithmetic and logical operations.

F Register

The F or Flag register indicates the state of certain arithmetic conditions after particular groups of instructions have been executed. A large number of Z80 instructions set flags depending on the values in various registers (usually A). When a flag is set, a bit in the F register is set to 1. For example if the result of a subtraction was zero the z flag would be set to 1. That is, bit 6 of the F register would be on. There are other instructions that will only work if a particular flag is set. One example of this would be RET Z. This means if the z flag is set then RET (return from subroutine), otherwise do nothing.

7	6	5	4	3	2	1	0
S	Z	X	H	X	P/V	N	C

The Flag register has 8 bits which can be either high or low (1 or 0). Each of these bits is set if certain conditions exist, although bits 3 and 5 are not actually used. If you want to see the mnemonics for each instruction and how the flags are affected you can find them in Appendix 3.

Carry flag

The Carry flag indicates whether there was an overflow from bit 7 of a register. It is mostly affected by addition, subtraction or shift instructions. By overflow we mean that, for example, adding 250 to 250 would give a value of 500. However the maximum value that can be held in 8 bits is 255 so the actual value left would be 244. Since the Carry flag would be set we know that the real value is 244+256 (500). The same applies to 16 bit values where a result would exceed 65535.

Some sample instructions that use the result of this flag are:

RET C; RETURN IF CARRY FLAG SET

JP NC, ADDRESS; JUMP to Address if CARRY NOT SET

N flag

The N flag, know as the add/subtract flag, cannot be used directly by

the programmer. It is used by the Z80 chip to record whether the last operation was a subtraction or an addition.

Parity/overflow flag

This is a dual purpose flag. When used to indicate parity the Parity flag is set (i.e. 1) if there is an even number of bits in the byte set to one. It is reset (i.e. 0) if the number of bits set to one is odd.

The flag can be used to represent overflow, if it is set when an arithmetic overflow occurs during an arithmetic operation. This might happen in an addition or subtraction operation involving two numbers with the same sign (i.e. both positive or both negative) and it changes the sign in the result.

H flag

The Half carry flag is used to indicate a carry from bit 3 of a byte to bit 4 of a byte.

The H and N flags are used by the CPU in order to do something known as binary coded decimal arithmetic (more about this later!)

Zero flag

The Zero flag is set by certain instructions when the result of that execution is zero.

Sign flag

The Sign flag is set by certain instructions which show the sign of a result i.e. if the result was negative then the Sign flag would be set. If the result was positive then the Sign flag would be reset.

HL' Register pair

These are the alternate H and L registers working as a 16 bit HL' register. It is included here just to serve as a warning about using HL' in USR subroutines. HL' is used to point to the calculator stack during USR subroutines and BASIC will probably crash if you RETURN TO BASIC with HL' changed.

PC Register

The PC, or Program Counter is a 16 bit register that holds the address in memory of the instruction currently being executed.

The SP Register

The SP or Stack Pointer is another 16 bit register. This one points to the current address at the top of the stack. Unlike the term *queue*, which indicates that literally the first item in is the first item out, the *stack* is a term used to represent data held in the reverse order, in which the last item placed will be the first item out. (This is sometimes known as a LIFO 'Last In First Out' list).

Imagine a pile of books onto which more books are placed. In order to get to the bottom of the pile the last book placed on top will have to be the first one removed. This analogy is very similar to the way in which the stack works on the Z80.

If we wish to call a routine in machine code we use an instruction `CALL` (This is similar to the `GOSUB` instruction in BASIC). When the Z80 executes a `CALL` instruction it places the return address onto the stack. The return address is always `PC+3`, because the `CALL` instruction is three bytes long and the subroutine must `RETURN` at the start of the next instruction after the `CALL`. It then gets the `CALL` address and puts this into the Program Counter (PC register). You will need to remember that the Program Counter points to the location of the instruction currently being executed, so the program will carry on running from that address. When the Z80 meets a `RET` instruction (`RETURN`) the chip then `POPS` the return address from the stack and places it back into the PC register.

This is very similar to what happens in a BASIC program when it executes the `GOSUB` command and then `RETURNS`. As well as saving return addresses, the stack can also be used to save data. (This can prove useful when you start to run out of registers.) For example we can save the HL register pair by using the instruction:

`PUSH HL`

This means 'PUSH the HL register pair on the stack'. We could now use the register pair for other calculations if we wanted to, knowing that we have a copy on the stack. To retrieve data from the stack we use the instruction:

`POP HL`

This means 'POP the data on top of the stack into the HL register pair'. It is important, however, to note the order in which we `PUSH` and `POP` data. For example, if we use the instructions:

`PUSH HL`

`PUSH BC`

we must remember to POP the data in the reverse order to that in which we originally pushed them. So to place the data back into the same registers we would need to use the instructions:

```
POP BC
```

```
POP HL
```

If we popped the data from the stack with:

```
POP HL
```

```
POP BC
```

then it would become apparent that the register pairs had been changed over. This can be a useful way of moving data within the chip but care must be taken when using the stack. Problems will arise when a PUSH or POP instruction is missing because a RET could POP some data and RETURN to the wrong address. A large proportion of machine code 'crashes' are caused by programmers wrongly using the stack in this way. *Remember 'Last In First Out'*

Let us examine the following code:

```
LD HL,0
```

```
PUSH HL
```

```
RET
```

The first instruction tells the computer to load the HL register pair with the number 0. The second is the PUSH instruction which places the HL pair onto the stack and leaves the number 0 on the top of the stack. The last instruction is the RETURN instruction which retrieves the last 16 bit number on the stack and places it into the program counter. Since the top of the stack contains 0 the program will start to run from address 0000 — Bad news if you have not SAVED your program!

Loading and storing

In order to manipulate information from one register to the other, from RAM to registers and vice-versa, we need to use what is known as loading operations. These operations can be used on both 8 bit and 16 bit registers and constitute the major part of the Z80 instruction set. So learn them well!

First let us look at a few 8 bit LOAD operations:

3E 16 LD A,22

The above instruction means 'LOAD the A register with the value 22 decimal'. It does precisely what it says: it puts the value 22 into the A register. The two digits on the lefthand side of the operation are its hexadecimal equivalent, which are **POKED** into memory or typed in using a monitor. (An assembler does automatically.) We can also LOAD other 8 bit registers with data.

Examples

```
06 16 LD B,22 ;LOAD B register with 22 decimal
06 22 LD B,22H ;LOAD C register with 22 hex
2E 04 LD L,4 ;LOAD L register with 4 decimal
0E 0C LD C,12 ;LOAD C register with 12 decimal
16 10 LD D,10H ;LOAD D register with 10 hex
1E FF LD E,255 ;LOAD E register with 255 decimal
26 56 LD H,56H ;LOAD H register with 56 hex
```

Here too the hexadecimal translation is given on the lefthand side of the mnemonic.

If you look at the first two examples, which LOAD the B register, you might notice something similar in their hexadecimal output. The first byte (06H) is the same in both instances. It is not a coincidence. The first byte of the instruction is known as the Op code and tells the computer which register we are dealing with. The second byte is the actual data which we are **LOADING** into the register. It is important to note that it is not possible to have an instruction such as:

LD A,289
'Load A register with 289 decimal'

This is because the number 289 takes more than 8 bits to represent it. We can however LOAD register *pairs* with 16 bit numbers.

16 bit LOADS

As we have mentioned before the Z80 chip has the facility for pairing off registers, a feature which gives access to some powerful 16 bit commands.

Let us recap which registers can be paired off together:

AF	AF'
BC	BC'
DE	DE'
HL	HL'

You can see from the diagram that the registers (with the exception of the Accumulator and Flag registers) are paired off in alphabetical order. The IX, IY, SP and PC registers have not been included in the diagram as these are true 16 bit registers and are not split into two like the others.

Let us now take a look at some 16 bit LOAD operations.

```
21 00 40 LD HL,16384
```

This means 'Load the HL register pair with 16384 decimal' (4000H) If you look at the hex translation, this time there are 3 bytes to represent the instruction. The first is the Op code for 'LD HL' and the last two are the data. The low part of the data is the second byte and the high part is the third byte. (Remember that the Z80 stores 16 bit values in the opposite way to which you would write them!)

Other examples of 16 bit LOAD operations are given below. (HH is the high byte of a number in hex while LL is the low byte).

```
01 LL HH LD BC,HHLL
11 LL HH LD DE,HHLL
31 LL HH LD SP,HHLL
DD 21 LL HH LD IX,HHLL
FD 21 LL HH LD IY,HHLL
```

Loading from one register to another

As well as LOADING 8 bit and 16 bit numbers into registers it is also possible to transfer information from one register into another.

Consider these examples:

```
78 LD A,B
79 LD A,C
6B LD L,E
```

The first example reads 'Load the A register with the B register'. If, for example, we had the instructions:

```
06 02 LD B,2 ;load B register with 2
```


and then added the following instruction:

```
78 LD A,B
```

we would find that the A register would take the contents of the B register, thus ending up with the value 2.

The Z80 chip does not have 16 bit instructions such as:

```
LD HL,DE ;load HL pair with DE pair????
```

so in order to achieve the same effect it is necessary to use a couple of 8 bit transfers, like this:

```
62 LD H,D ;load H register with D register
6B LD L,E ;load L register with E register
```

Easy, isn't it!

The only 16 bit register to register load operations allowed in Z80 code are the following which deal exclusively with the stack pointer.

```
F9 LD SP,HL
DD F9 LD SP,IX
FD F9 LD SP,IY
```

The next mode of addressing data is very similar to the way in which the BASIC instruction PEEK and POKE work. We are going to look at examples which load and store from locations in RAM and ROM.

```
3A 00 40 LD A,(16384)
```

The instruction above reads 'LoaD the A register with the contents of the address 16384 (4000H). You can think of it as being similar to the BASIC instruction:

```
LET x=PEEK(16384)
```

The number at the location 16384 is put into the A register. We could also put the contents of the A register into RAM by the instruction:

```
32 LL HH LD (ADDRESS),A
```

If we used the following instructions:

```

3E FF LD A,255
32 00 40 LD (16384),A

```

the first instruction would load the A register with the value 255 and the second would put the value of this register into the address 16384.

The Accumulator is the only 8 bit register which allows us to do this kind of addressing. There are *no* instructions such as:

```
LD (16384)B, ;load the address 16384
with b?
```

One way to get over this problem would be to use the instruction:

```

78 LD A,B ;let A register=B register
32 00 40 LD (16384),A ; put A register in 16384

```

Sixteen bit addressing in this mode is quite extensive; here are some examples of the instructions allowed.

```

ED 4B LL HH LD BC,(HHLL)
ED 5B LL HH LD DE,(HHLL)
ED 6B LL HH LD HL,(HHLL) ;most assemblers use the faster form
of this instruction which is 2A LL HH
DD 2A LL HH LD IX,(HHLL)
FD 2A LL HH LD IY,(HHLL)
ED 7B LL HH LD SP,(HHLL)

```

These instructions are 16 bit load instructions so they read two bytes from a given address. We could use:

```
2A 53 SC LD HL,(23635)
```

which reads 'load the HL register pair with the contents of address 23635(5C53H). This would take the contents of the address 23635 and place them in the L register (low byte first). Finally it would take the contents of 23635+1 (i.e. 23636) and place it in the H register.

It is also possible to save the contents of registers at a given address, as follows:

```
ED 63 00 40 LD (16384),HL ;most assemblers would use the more
efficient 22 00 40 form of this
instruction
```

'LoaD at the location 16384 the value in the HL register pair.'

This instruction will put the value of the L register at the address 16384 and then put the value of H at the address 16385.

```
21 AA 22 LD HL,22AAh ;load HL with 22AA hex
22 00 40 LD (16384),HL
```

The two instructions above would load AA hex at location 16384 and 22 hex at the location 16385.

Now suppose we wanted to load a value into the A register from an address which we did not directly know. The address can be worked out from a calculation. We would address that value by a method known as *register indirect addressing*. Sounds complicated, doesn't it? Don't worry, it's all very easy. All this means is that instead of giving an address directly to load from we have that address pointed to by a register pair, as you will see.

```
7E LD A,(HL)
```

The instruction above reads: 'LoaD the A register with the contents pointed by the address in the HL register pair'. If HL contained 16384 then the contents of that address would be put in the A register.

It is also possible to save using register indirect addressing, as follows:

```
77 LD (HL),A
12 LD (DE),A
36 22 LD (HL),22h ;load 22h at the address in HL
```

The last instruction here is unique to the HL register pair. It is one of the most important and powerful register pairs available on the Z80 chip.

Last, but by no means least, is the powerful *index addressing mode*. These use the IX and IY registers and are extremely useful in accessing arrays of data.

The index modes are in the form:

DD RR NN	LD r,(IX+nn)
FD RR NN	LD r,(IY+nn)
DD RR NN	LD (IX+nn),r
FD RR NN	LD (IY+nn),r
DD 36 NN dd	LD (IX+nn),d
FD 36 NN dd	LD (IY+nn),d

where RR depends upon the register being used and dd represents the data. r is any of the registers A,B,C,D,E,H,L. nn is an offset with the value of 0 to 127 & 0 to -128. This is derived from the signed binary value of the number, which is added to the value of the index register. The store or load is then done at the resultant address. d is a byte value which can be loaded and stored directly.

Consider the following:

DD	21	00	60	LD	IX,6000h
DD	4E	05		LD	C,(IX+05)
DD	36	00	03	LD	(IX+00),03
addressdata					
	6000	00			
	6001	02			
	6002	04			
	6003	05			
	6005	06			
	6006	07			

After executing the first line the IX register is pointing to the portion of RAM/ROM at the address 6000 hex. When the second instruction is executed, the offset value 05 is added to the value of the IX register, which equals 6005 hex, and the contents of this location are put into the C register. Thus, the C register will contain the value 06. Note that the address in the IX register is not changed in any way. After executing this instruction it merely accesses the contents of that address. The last instruction:

LD (IX+00),03

goes through the similar process of working the offset address which is $6000+0=6000$ hex and this time stores the value 03 at that address. The IX register works in a similar way . . . but a word of warning! If you are using the IX register on the Spectrum be very careful when mixing machine code with BASIC, as the Spectrum uses the IX register to point to the system variables. The procedure, as explained earlier when you were introduced to the IX register, would have to be applied when using the IX register in your own programs.

2 Number crunching

So far we have looked at the way the Z80 stores data and how it can transfer values and control from one address to another. In this chapter we come to the actual number crunching instructions used in addition and subtraction. As was pointed out before, the main advantage of the Z80 chip over other 8 bit microprocessors is that it can handle 16 bit numbers directly, making addition and subtraction operations that much easier. To begin with let's take a look at the 8 bit arithmetic instructions.

The two simplest number crunching instructions are DEC, 'DECREMENT register' and INC, 'INCREMENT register'. These two instructions respectively subtract or add 1 to the value in a specified register. We are allowed to use single registers A,B,C,D,E,L and H, with these instructions, so the range of possible commands is:

DEC A	INC A
DEC B	INC B
DEC C	INC C
DEC D	INC D
DEC E	INC E
DEC H	INC H
DEC L	INC L

The Accumulator or A register is one of the main registers in the Z80 chip and allows 8 bit arithmetic operations which can work directly with other registers and numbers. To add to the A register a value held in another register we use the instruction:

ADD A,r

which means 'Add to the Accumulator the value in register r', where r can be any register of A,B,C,D,E,H or L.

If we wanted to add numbers directly to the Accumulator we could use the instruction:

ADD A,N

where N is any 8 bit number. So for example, `ADD A,5` would add 5 to the Accumulator.

We can also use the `ADD` instruction in conjunction with something known as *indirect addressing*. The `HL` register pair contains an address where the actual number which we wish to add to the Accumulator is stored:

```
ADD A,(HL)
```

The above instruction actually performs the operation 'add to the Accumulator the contents of the location pointed to by the register pair `HL`'. Take for example the following code:

```
LD A,8
LD HL,6000H
ADD A,(HL)
```

We'll assume we have the following data stored in memory from address `6000` hex onwards:

Address	Contents
<code>6000H</code>	<code>02</code>
<code>6001H</code>	<code>03</code>
<code>6002H</code>	<code>06</code>
<code>6003H</code>	<code>07</code>

The first instruction would set the Accumulator to 8 decimal. The `HL` register pair is then set to point to the address `6000` hex. The final instruction then gets the value from the address at `HL` (i.e. `6000` hex) and adds it to the Accumulator. This leaves it with the value 10 decimal.

Pursuing the indirect method even further, it can also be used with indexing utilising the `IX` or `IY` registers.

Using the same data and starting at address `6000` hex let us run through the following example to demonstrate this:

```
LD A,0
LD IX,6000H
ADD A,(IX+0)
ADD A,(IX+3)
```

The first and second instructions are simple enough. These set the `A` register to zero and the `IX` register to the address `6000` hex.

ADD A,(IX+0)

The instruction above adds the index to the address in the IX register. This new address is then used to point to the data which we wish to use. Since our index is zero, the address calculated is $6000H+0H=6000H$. Therefore the contents are taken from this address and added to the accumulator, leaving it with a value of 2 after the first ADD instruction. The second addition is similar but uses the index 3, which means that the data to be added is stored at the address $6000H+03H=6003H$ and has the value 07. When this is added to the Accumulator the final result will be 9.

Subtraction works on the same registers as the ADD instruction, the mnemonic being SUB. Again every operation is done on the A register but the actual format of the mnemonic is slightly different as it does not actually mention the A register. The operand follows the SUB instruction directly. For example, to add the B register to the A register we would write:

ADD A,B

but to subtract the B register from the Accumulator we would write:

SUB B

Not too confusing, hopefully!

ADD A,A	SUB A
ADD A,B	SUB B
ADD A,C	SUB C
ADD A,D	SUB D
ADD A,E	SUB E
ADD A,H	SUB H
ADD A,L	SUB L
ADD A,(HL)	SUB (HL)
ADD A, (IX+d)	SUB (IX+d)
ADD A, (IY+d)	SUB (IY+d)
ADD A,N	SUB N

It is useful to note that ADD A, A is a quick and efficient instruction for doubling the value in the A register. SUB A is a quick way of setting the A register to zero (it works nearly twice as fast as LDA,0 and only takes up one byte instead of two).

Using the Carry flag

There is another set of 8 bit arithmetic instructions which take into account the state of the Carry flag. These are known as the `ADC` (Add with Carry) and `SBC` (Subtract with Carry).

In the case of addition the `ADC` adds the state of the carry flag as well as the given register or data. So, for example, if the `A` register contained 5 and the Carry flag was high (i.e. set to 1), if we ran the instruction:

`ADC A,2`

the answer left in the `A` register would be $5+2+1=8$. On the other hand, if the Carry flag were to be reset we would return with the answer 7 as with the normal addition.

When it comes to subtraction, we subtract the state of the Carry flag from the Accumulator. So, if we had 5 in the `A` register and the Carry flag was set, the instruction:

`SBC A,3`

would leave the answer in the Accumulator as $5-2-1=2$

<code>ADC A,A</code>	<code>SBC A</code>
<code>ADC A,B</code>	<code>SBC A,B</code>
<code>ADC A,C</code>	<code>SBC A,C</code>
<code>ADC A,D</code>	<code>SBC A,D</code>
<code>ADC A,E</code>	<code>SBC A,E</code>
<code>ADC A,H</code>	<code>SBC A,H</code>
<code>ADC A,L</code>	<code>SBC A,L</code>
<code>ADC A,(HL)</code>	<code>SBC A,(HL)</code>
<code>ADC A,(IX+d)</code>	<code>SBC A,(IX+d)</code>
<code>ADC A,(IY+d)</code>	<code>SBC A,(IY+d)</code>

The 16 bit increment and decrement instructions work in exactly the same manner as their 8 bit equivalents, but on pairs as opposed to single registers. The instruction `DEC BC` subtracts 1 from the value held in the `BC` register pair, while the instruction `INC DE` adds 1 to the `DE` pair. Because we are dealing with 16 bit operations we also have the option to increment or decrement the `IX`, `IY` and `SP` registers.

INC BC	DEC BC
INC DE	DEC DE
INC HL	DEC HL
INC IX	DEC IX
INC IY	DEC IY
INC SP	DEC SP

16 bit addition is quite versatile on the Z80. It allows the user to add (with or without Carry) other 16 bit registers to the HL, IX or IY register pair. Subtraction, however, is limited to subtracting the registers BC, DE, HL and SP from the HL pair and we only have the use of the Subtract with Carry instruction.

ADD HL,BC	ADC HL,BC
ADD HL,DE	ADC HL,DE
ADD HL,HL	ADC HL,HL
ADD HL,SP	ADC HL,SP
ADD IX,BC	ADD IY,BC
ADD IX,DE	ADD IY,DE
ADD IX,IX	ADD IY,IY
ADD IX,SP	ADD IY,SP

16 Bit subtraction.

SBC HL,BC
SBC HL,DE
SBC HL,HL
SBC HL,SP

Let us look at a few examples using some of these instructions.

```
LD HL,0432H
LD BC,0536H
ADD HL,BC
```

The above instructions would result in the HL pair containing $0432H + 0536H = 0968H$.

The ADD HL,HL instruction has the same effect as multiplying by 2. Combined with additional instructions it could be used to multiply a number by a power of two. For example, suppose we wished to multiply the contents in the DE pair by 32. First we transfer DE into HL, then we do five ADD HL,HL instructions in order to multiply by 32, and finally we transfer the answer back into DE like this:

```
;multiply DE pair by 32
```

```

EX DE,HL      ;SWOP DE AND HL
ADD HL,HL     ;TIMES BY 2
ADD HL,HL     ;TIMES BY 4
ADD HL,HL     ;TIMES BY 8
ADD HL,HL     ;TIMES BY 16
ADD HL,HL     ;TIMES BY 32
EX DE,HL      ;SWOP DE AND HL
               ;ANSWER IS NOW IN DE

```

The first and last instructions `EX DE,HL` mean 'exchange the `DE` and `HL` registers'. What they actually do is simply to swop the contents of the `DE` pair for the contents of the `HL` pair.

As we mentioned earlier, the Add with Carry instruction `ADC` takes into account the state of the Carry flag. For example, if the Carry flag were set and we used the instruction:

```

LD HL,0432H
LD BC,0536H
ADC HL,BC

```

the `HL` pair would contain $0432H + 0536H = 0968H + 1$ (state of Carry) = $0969H$. It is worth repeating that the only form of subtraction available with the 16 bit set is using the `SBC` instruction which also subtracts the state of the Carry flag to give the final result. Therefore, it is sometimes necessary to clear or reset this Carry flag before executing an `SBC` instruction in order to obtain the correct result. The way to do this is very simple. We use the 1 byte instruction:

AND A

This means 'AND the Accumulator with itself'. This is known as a *logical operation*, a process which we will be looking at more closely in chapter 5. All you need to know for now is that one of the effects of this instruction is to reset the Carry flag. Thus in order to subtract `0432` hex from `0563` hex we could use the following piece of code:

```

LD HL,0536H   ;Put first number in HL
LD DE,0432H   ;Put second number in DE

AND A         ;clear the carry flag
SBC HL,DE     ;do the subtraction!

```

This should leave the result $0536H - 0432H - 0$ (state of Carry) = 104 hex

If we had not used the `ANDA` instruction as a precaution to clear the Carry flag and if the Carry flag was set after the execution of a previous instruction, the result would be `0536H-0432H-1` (state of Carry)=`103H`.

Jumping and calling

In Spectrum BASIC we transfer control from one part of a program to another using the BASIC instructions `GOTO` and `GOSUB`. In order to implement transfers in machine code we use the `JUMP` and `CALL` instructions.

The simplest of these instructions is the `JUMP to address` command:

```
C3 00 60      JP 6000H
```

The above example reads 'JUMP to the address `6000` hex' and it loads the program counter with `6000` hex from where it will continue to execute the machine code.

We can also specify the address to `JUMP` to by the register pairs `HL,IX` and `IX`. For example, if we had the instruction:

```
JP (HL)
```

This would in effect load the program counter with the `HL` register pair. So if the `HL` pair contained `1601` hex the program would `JUMP` to the address `1601` hex.

In order to implement the equivalent of the BASIC statement 'IF condition THEN GOTO' we have to use something known as conditional jump instructions. There are eight conditions which can be identified, all of which are indicated by bits set in the flags register (F-register). Below we give all the conditional jump statements that are allowed:

```
JP NO,address ;'Jump if Carry flag reset (Non Carry)'
                ;to the address specified
JP C,address  ;'JUMP if Carry flag set (Carry)'
JP NZ,address ;'JUMP if Zero flag reset (non Zero)'
JP Z,address  ;'JUMP if Zero flag set (Zero)'
JP P,address  ;'JUMP if positive (Sign flag reset)'
JP M,address  ;'JUMP if minus (Sign flag set)'
JP PO,address ;'JUMP if Parity odd (Parity reset)'
JP PE,address ;'JUMP if Parity even (Parity set)'
```

Jump relative

There is another range of `JUMP` instructions available on the `Z80`, known as the *JUMP relative command*. This instruction allows us to specify an offset instead of an absolute address. The offset is a one byte number and allows us to jump backwards by up to 128 bytes and forwards up to 127 bytes, counted from the first byte after the instruction. This is because by using signed integer representation (see chapter 1) a byte can hold values between +127 and -128. The actual instruction is written as follows:

28 dd JR dd

JUMP relative dd bytes, where *dd* is the displacement to `JUMP`. For example, in the case below:

18 03	JR 03
00	NOP
00	NOP
00	NOP
3E 04	LD A,4

the code would load the `JUMP` past the two `NOP` (No operation) instructions to the instruction which `LOADS` the Accumulator with the value 4. The displacement `02` is added to the location after the `JUMP` instruction. Since the `JUMP` relative instruction is two bytes long the actual address to which the program is transferred is the address of the `JUMP` relative instruction *plus* the displacement *plus* 2:

$$\text{new address} = \text{old address} + \text{displacement} + 2$$

If you are using a monitor to type in a machine code program you will have to work out the displacement for yourself. However, most `Z80` assemblers will let you reference addresses as labels and will automatically work out the displacement needed. So you could write the code like this:

```

JR Here
NOP
NOP
Here LD A,4

```

When assembled the displacement would be placed with the appropriate value.

Like the absolute `JUMP` the relative `JUMP` also has conditional

options. However, these are limited to the testing of the carry and the zero flags:

```
JR C,dd      ;'JUMP relative on Carry (Carry flag set)'
JR NC,dd     ;'JUMP relative non Carry (Carry flag reset)'
JR Z,dd      ;'JUMP relative on Zero (Zero flag set)'
JR NZ,dd     ;'JUMP relative non Zero (Zero flag reset)'
```

The advantage of using the `JUMP` relative instructions as opposed to those of the `JUMP` absolute lies in relative addressing. This takes only two bytes as compared to the three needed for the absolute mode, making a routine smaller in size. It also allows some particular routines to be relocateable, that is, having the ability to be placed anywhere in memory without having to be re-assembled.

DJNZ

The `DJNZ` 'DECRement JUMP on non zero' is an extremely powerful instruction. It allows the programmer to effect a loop a specified number of times around a portion of code, very much like the 'FOR...NEXT' statements in BASIC. Take a look at the following machine code program:

```
                LD B,20H
                LD HL,5800H
                LD A,2
LOOP           LD (HL),A
                INC HL
                DJNZ LOOP
                RET
```

The first instruction `LD B,20H` LOADS the B register with the number 20 hex (32 decimal). The B register is used as a loop counter for `DJNZ`.

We then LOAD the HL register with the two byte number 5800 hex. This is the start of the attribute file:

```
                LD HL,5800H
```

The Accumulator is LOADED with the value 2, the colour code for red INK, black PAPER, BRIGHT 0 and FLASH 0. The next three instructions form the main part of the loop:

```
LOOP           LD (HL),A
                INC HL
                DJNZ LOOP
```

The value in the Accumulator is placed at the address pointed by the HL pair. When executed the first time round, the loop will load the value 2 into the start of the attribute file. Next we have the instruction:

```
INC HL
```

This means 'increment the HL register pair by one' and adds one to the HL pair so that it points to the next address in the attribute file. Finally we have:

```
DJNZ LOOP
RET
```

The DJNZ instruction will subtract one from the B register. If the value after this subtraction is not zero then it will JUMP relative to the address specified. If it is zero then it will go on to the next instruction which is a RETURN.

As you can see, DJNZ is an extremely powerful instruction. It is very much like having two instructions in one – a subtraction on the B register and a JUMP relative on non zero.

Bearing in mind that the DJNZ instruction uses relative and not absolute addressing we can only use it if the portion of code we are looping around is no longer than 128 bytes.

Calling and returning

The second method of transferring the control of a program is by using the set of CALL instructions.

There are times when a program executes the same portion of code many times or when other portions of code closely resembling each other are run with different parameters. Instead of having these similar routines scattered around at various different places in memory, you could have just one copy of this code when necessary, call it as a subroutine, very much like setting up a subroutine in BASIC using the 'GOSUB' BASIC instruction.

You can call this piece of code by using the instruction CALL followed by an address. The flow of the program will transfer to this address after storing the address of the instruction following the CALL instruction. The program is then executed normally until it reaches a RET (return) instruction, when it returns to the next instruction after the address of the call.

The CALL instruction takes this syntax:

```
CD LL HH          CALL HLLL
```

HH is the high byte of the address and LL is the low byte. It is possible for CALLS to be *nested*, which means that one subroutine may CALL another subroutine. In fact the number of nested calls allowed is limited only by the amount of memory left to the programmer. A subroutine may also call itself, a function known as recursion which is too abstruse for us to pursue here in any depth.

Like JUMP, the CALL and RETURN instructions also have conditional counterparts. We can CALL a subroutine or RETURN from a subroutine depending on the conditions set in the Flags register:

CALL HHLL	RET
CALL Z,HHLL	RET Z
CALL NZ,HHLL	RET NZ
CALL C,HHLL	RET C
CALL NC,HHLL	RET NC
CALL PO,HHLL	RET PO
CALL PE,HHLL	RET PE
CALL M,HHLL	RET M
CALL P,HHLL	RET P

There is another range of calling instructions, known as the restart (RST) set. They differ from the others in that they are only one byte long and are limited to CALLING one of eight addresses: 00 hex, 08 hex, 10 hex, 18 hex, 20 hex, 28 hex, 30 hex and 38 hex.

As you have probably noticed, all these addresses are in the ROM memory map which you may not find much use as we cannot write any code there. Well that's true, but we can CALL some of the routines from our own programs. Below are the CALLS and the object of the particular routines.

RST 00H ;start boot up

This is a bit like typing NEW in BASIC, so is not very useful unless you wish to return to BASIC from a machine code program and protect the routine from prying eyes.

RST 08H ;error restart

This routine is used by BASIC to report error messages. The error number is the byte following the restart instruction. It will give the error report of the data plus one. Thus:

```
RST 08H
DB 08
```

will generate the error message 09 'STOP statement'.

```
RST 10H ;print a character
```

This is an extremely useful routine. It prints the character in the Accumulator to the current channel. A channel outputs to a 'device', which can be either the printer or various parts of the screen. We'll see more of this in Chapter 9. A simple example for now is:

```
LD A,66 ;print the character B to the
RST 10H ;current channel
```

```
RST 28H ;floating point calculator
```

The number crunching routine above allows us easily to implement complex floating point arithmetic routines in machine code using the ROM functions. The floating point calculator is explained in more detail in Chapter 9.

```
RST 30H ;make space
```

This is not a particularly useful routine. It simply creates space in the workspace area.

Finally:

```
RST 38H ;scan the keyboard
```

This routine updates the system variable LAST-K and can be used to ascertain which keys are depressed. It is called 50 times a second by BASIC. It is also sometimes known as the Mode 1 maskable interrupt routine. We'll be hearing more about this routine when we get to Chapter 11, which deals with interrupts and their uses.

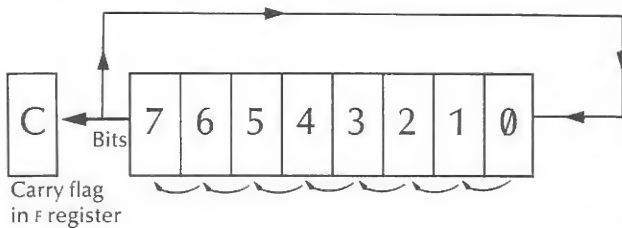
3 Rotating and shifting

Rotating and shifting operations provide the programmer with the means to manipulate the pattern of bits held in a register or a byte memory. These instructions, which are most useful for multiplication and division by powers of two, act on most of the 8 bit registers. They can use both indirect and index addressing modes. All the rotate instructions use the Carry flag (which is held in the F register) as a ninth bit, bit 8, therefore allowing the programmer to rotate this from the left or right through the register or memory. This should become clearer as we run through the available instructions.

Rotating

RLC Rotate Left Circular

This instruction rotates each bit of a given register or memory byte to the left by one bit. Bit 7 of the register or byte specified is rotated to the Carry flag and the same value is 'wrapped round' to bit 0:



For example, if the byte on which we were operating held 10101010 the following would occur after the RLC instruction was executed. The value of bit 7 (1) would be transferred to the Carry flag bit and to bit 0 of the byte with each of bits 0 to 6 shifted one place to the left. The result would be 01010101 stored in the byte, and the Carry flag set.

The RLC instruction can act on the registers A,B,C,D,E,H,L, as well as (HL) and (IY+INDEX) and (IX+INDEX). There is also a RLCA instruction which has the same effect as RLC A but is one byte shorter and twice

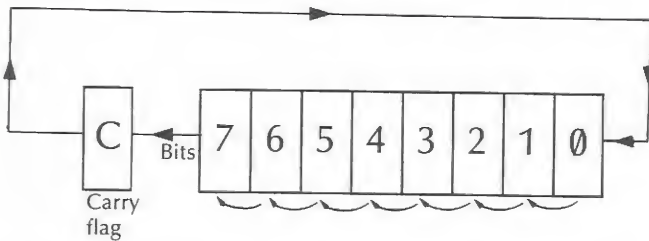
as fast to execute. These additional, short-form, rotate instructions on the Accumulator are available on all the rotate instructions. These are as follows (note that *d* indicates the index value where applicable):

RLCA	RLC (HL)
RLC A	RLC (IX+d)
RLC B	RLC (IY+d)
RLC C	
RLC D	
RLC E	
RLC H	
RLC L	

RL Rotate left

This instruction rotates the register left through all the nine bits, wrapping around the carry bit value to bit 0.

The effect of this instruction is to take the sequence of bits in the byte, add the Carry flag value as bit 8, and then shift all bits one place to the left. The Carry flag value then goes into bit 0. Thus if we had (1)01010101 before an RL instruction, we would end up with (0)10101011. This would produce a result which is the original value multiplied by two, plus the value of the Carry flag.

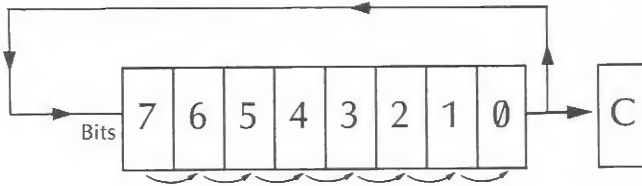


The available instructions are:

RLA	RL (HL)
RL A	RL (IX+d)
RL B	RL (IY+d)
RL C	
RL D	
RL E	
RL H	
RL L	

RRC Rotate Right Circular

The register or byte is rotated right from bit 7 through to bit 6 and so on. Bit 0 is then rotated to the Carry flag and bit 7. This is the reverse operation to that of RLC

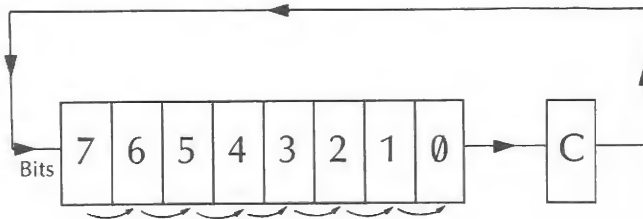


RRCA
RRC A
RRC B
RRC C
RRC D
RRC E
RRC H
RRC L

RRC (HL)
RRC (IX+d)
RRC (IY+d)

RR Rotate Right

The Rotate Right instruction has the opposite effect to that of the RL Rotate left instruction. Bit 0 of the register or byte is rotated to the right through the Carry, while the old Carry is rotated down to bit 7.



RRA
RR A
RR B
RR C
RR D
RR E
RR H
RR L

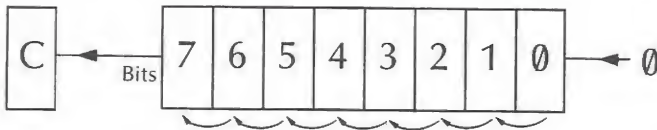
RR (HL)
RR (IX+d)
RR (IY+d)

Shifting

As well as the Rotate instructions, there is also available a set of shift instructions which can make registers shift either left or right. This differs from the Rotate instruction set in that there is no 'wrap around' effect. Therefore one bit at either end of the byte is lost and a zero goes into this bit. Like the Rotate set all the shifts can act on A, B, C, D, E, H, L, (HL) (IX+d) and (IY+d).

SLA Shift Left Arithmetic

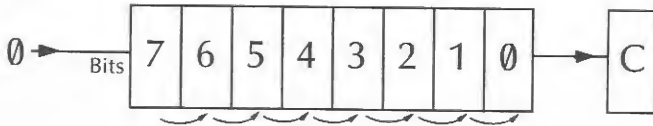
The content of the carry bit is lost and the whole byte or register shifts to the left. Bit seven is shifted into the Carry flag, and a 0 inserted in bit 0.



- | | |
|-------|------------|
| SLA A | SLA (HL) |
| SLA B | SLA (IX+d) |
| SLA C | SLA (IY+d) |
| SLA D | |
| SLA E | |
| SLA H | |
| SLA L | |

SRL Shift Right Logically

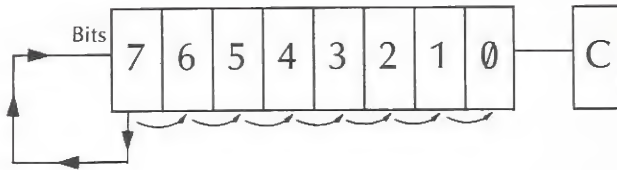
The SRL 'Shift Right Logically' shifts the bits from the left to the right, so is useful for dividing numbers by powers of two. Bit zero of the register/byte is shifted into the carry bit and a zero is placed into bit seven.



- | | |
|-------|------------|
| SRL A | SRL (HL) |
| SRL B | SRL (IX+d) |
| SRL C | SRL (IY+d) |
| SRL D | |
| SRL E | |
| SRL H | |
| SRL L | |

SRA Shift Right Arithmetic

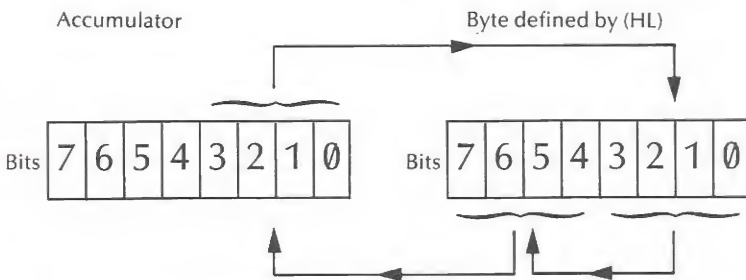
This is an odd instruction. Shift Right Arithmetic is identical to the SRL instruction apart from the fact that bit seven is left unchanged. This instruction is used to divide 'signed' numbers (i.e. numbers -127 to +128) by powers of two as it doesn't affect the sign bit.



SRA A	SRA (HL)
SRA B	SRA (IX+d)
SRA C	SRA (IY+d)
SRA D	
SRA E	
SRA H	
SRA L	

RLD (HL) ROTATE LEFT DECIMAL

This is a single instruction which acts on both the accumulator and the contents pointed to by the HL register pair. It actually moves 'half bytes' called 'nybbles' from the Accumulator to a RAM location and vice versa.



As you can see from the diagram the bottom four bits (bits 3-0) of the location pointed to by the HL register pair are shifted to the top four bits positions (7-4). The original top four bits are placed in the lower half of the accumulator with the original contents placed in the bottom four bits of the RAM location. If, for example, we had the HL pair containing 6000 hex, this byte holding CB hex, and the accumulator containing 2A:

Address	Contents	Accumulator
6000	2A	CB

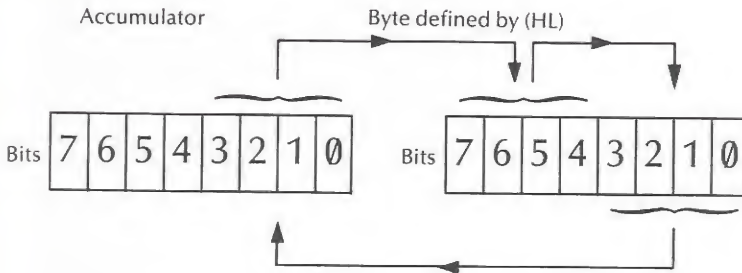
Then after executing the instruction

RLD (HL)

We would find that the contents of location 6000H and the accumulator would be changed to:

Address	Contents	Accumulator
6000	AB	C2

The instruction RRD (HL) has the opposite effect, as shown in the diagram below:



As already indicated, the shift instructions are very useful for multiplying and dividing by powers of two. If, for example, the Accumulator contained the value three and we had the instructions:

SLA A
SLA A

Then the result remaining in the Accumulator would be 12. Remember that the shift instructions affect the Carry flag, so if we had executed the instructions:

LD A,128
SLA A

the bit pattern for 128 is 1000000. Therefore when the SLA A instruction is carried out, the top bit of the Accumulator would be shifted into the Carry flag. Zero remains in the A register leaving the Carry flag and Zero flag set. It is easy to write small routines to multiply

registers by numbers which are not multiples of two. For example to multiply a number by 10 simply split the calculation into two parts. First multiply the number by eight and then add twice the original number.

MULT10

```
SLA A    ;LET A=2*A
LD B,A  ;LET B=A (2*original A)
SLA A    ;LET A=2*A (4*original A)
SLA A    ;LET A=2*A (8*original A)
ADD A,B ;LET A=A+B (10*original A)
```

The first two instructions:

```
SLA A
LD B,A
```

multiply the Accumulator by two and save the result in the B register. Remember, the instruction

```
LD B,A
```

has no effect on the Accumulator but copies its contents into the B register. Therefore at this point we have double the original number in both the A and the B registers.

```
SLA A
SLA A
```

The two other shift instructions multiply the number by eight. Finally, the last instruction:

```
ADD A,B
```

adds the contents of the B register, which contains twice our original number, to the A register. This leaves the desired answer.

This method of multiplication would only work for numbers in the range of 0 to 25. Any larger number would result in a number greater than 255 which we are unable to fit into an eight bit byte. To perform multiplication on two byte numbers, using shifts, we have to take into account that a Carry may occur from the lower half of a register. This must be shifted to the high part. Therefore to multiply the HL register pair by two we use the instructions:

```
SLA L           ;multiply lower part by two
RL H           ;rotate putting carry into bit
               ;0 in high register.
```

If we wanted to multiply the HL register pair by ten we could write:

```
SLA L
RL H           ;2*HL
LD E,L
LD D,H        ;Save in DE.ie DE=2*HL
SLA L
RL H           ;4*HL
SLA L
RL H           ;8*HL
ADD HL,DE    ;HL=8*HL+2*HL=10*HL
```

Of course it would be much easier to use the ADD instruction to perform the multiplication.

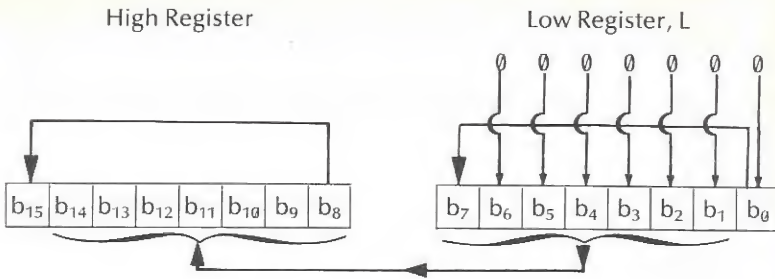
```
ADD HL,HL ;2*HL
LD E,L
LD D,H    ;DE=2*HL
ADD HL,HL ;4*HL
ADD HL,HL ;8*HL
ADD HL,DE ;10*HL
```

This piece of machine code is much faster and more concise to use than the previous example. However, this is not always the case. Suppose we wanted to multiply the HL pair by 128. The first thing that you thought of was probably to use the series of ADD instructions.

```
ADD HL,HL ;2*HL
ADD HL,HL ;4*HL
ADD HL,HL ;8*HL
ADD HL,HL ;16*HL
ADD HL,HL ;32*HL
ADD HL,HL ;64*HL
ADD HL,HL ;128*HL
```

A lot of instructions!

If we take a look at the bit pattern of a two byte number when we multiply by 128 we might be able to use shift and rotate instructions to our advantage. Let's look at the bit pattern we must get in order to multiply a two byte number by 128:



The top seven bits of the low byte need to be shifted into the bottom seven bits of the high byte. Bit 0 of the low byte will be shifted up to bit seven and bit seven of the high byte is lost.

If we represent the bit patterns by having h_n representing bit n of the high byte and l_n to represent bit n of the low byte then before we perform the multiplication we have the pattern:

h_7	h_6	h_5	h_4	h_3	h_2	h_1	h_0	high byte
l_7	l_6	l_5	l_4	l_3	l_2	l_1	l_0	low byte

After multiplying a two byte number by 128 we end up with the bit pattern:

h_0	l_7	l_6	l_5	l_4	l_3	l_2	l_1	high byte
l_0	0	0	0	0	0	0	0	low byte

Notice that the first seven bits of the low byte will always be set to zero. Looking at the pattern we can see that we can get the new high byte pattern by shifting the old low byte to the left one. Before we do this we can put h_0 into the Carry flag using the instruction:

SRL H

Now we have bit 0 of the H register i.e. h_0 in the Carry flag. We can now get the pattern we need for our new high byte in the low byte L register by using the instruction.

RR L

This puts the Carry (containing the old value h_0) into bit seven of the low byte. All the other bits are shifted to the right forcing the carry into the topmost bit. We now have the pattern we want for the H

register in the L register, so we transfer this by a simple LOAD command:

```
LD H,L
```

Finally, we set bit seven of the low byte to the contents of the Carry which is the bit 10. We do this by first setting the L register to zero and then rotating the Carry through to bit seven:

```
LD L,0
RR L
```

So our code for multiplying the HL register pair by 128 looks like this:

```
MULT128:    SRL H
             RR L
             LD H,L
             LD L,0
             RR L
```

The code is much smaller and faster to use than the series of ADD instructions. This portion of code is as much as 50% faster than the equivalent shown earlier. Arithmetic using bit manipulation is a little difficult to grasp at first but its implications are enormous. Screen addresses, for example, can be calculated much faster, giving games of infinite quality. Therefore it's very worthwhile to take some time to learn. Meanwhile, I'll end the chapter by giving you a routine to divide the HL register pair by 128 and let you find out how it works.

```
DIV128:     SLA L
            RL H
            LD L,H
            LD H,0
            RL H
```

4 Making comparisons and checking bits

The Compare Instruction

A COMPARE instruction operates in a similar fashion to a subtraction operation, except that the Accumulator is not changed. Instead, various flags are set or reset according to the result.

This instruction is most useful when used in conjunction with the Z80's conditional CALL and JUMP instructions and can be used to implement machine code equivalents of such BASIC statements as:

```
IF X>N THEN GOTO ADDR
IF X<N THEN GOTO ADDR
IF X=0 THEN GOTO ADDR
IF X<>0 THEN GOTO ADDR
```

The COMPARE instruction is limited to comparisons between eight bit numbers specified directly, indirectly, or contained in registers. Suppose we wanted to COMPARE the current value held in the Accumulator with the number 128 decimal.

We would use the instruction:

CP 128

which reads 'ComPare the current Accumulator value with the number 128'. This will set the Zero flag if the number in the Accumulator is equal to 128. The Carry flag will be set and the Zero flag reset if the number is greater than 128. Both the Carry flag and Zero flag would be reset if it was less than 128.

The following routine is a good example of the use of the COMPARE instruction with conditional branches to simulate the BASIC language IF . . . THEN . . . ELSE structure. The routine compares the A register with the B register and branches off to certain addresses, depending on whether the A and B registers are found to be equal, greater than, or less than A.

```

        CP    B
        JR    Z,EQUAL
        JR    C,BGREAT
LESSA:
        ..
        ..
        ..
        ..

```

The first instruction 'ComPare the B register with the A register' subtracts the B register from the A register. The actual result is not updated to the accumulator and only affects the flags.

If the A register was equal to the B register then the Zero flag would be set, causing the program to jump to the address labelled EQUAL. If they were not equal, then the program would carry on to the next instruction:

```
JR C,BGREAT
```

If the Carry flag was set this would indicate that the B register was greater than the Accumulator, causing a branch to the label BGREAT. If no branch occurred this would mean that the B register was less than the A register, causing the program to arrive at the label LESSA.

```

CP n           CP (HL)
CP A           CP (IX+dd)
CP B           CP (IY+dd)
CP C
CP D
CP E
CP H
CP L

```

Set Bit and Reset

There are other bit instructions in the Z80 set which allow us to set, reset or test individual bits in a byte.

The SET instruction

The 'SET' instruction allows us to SET a particular bit in a byte. We can test individual bits in a register or a RAM location. The format of the SET instruction can be any of the following forms:

```

SET n,r
SET n,(HL)
SET n,(IX+dd)
SET n,(IY+dd)

```

where n is the bit number we wish to test 0–7 and dd is an offset in the range –127 to 128

For example, the instruction:

```
SET 4,A
```

would set bit 4 of the Accumulator.

We can also use the indexing addressing mode to set and reset bits. If, for example, the IV register pointed to the address 6000H and the contents of its adjacent memory locations were as below:

Address	Contents
6000H	22H
6002H	00H
6002H	08H

Then the instruction SET 4, (IY+2) would have the following effect:

address	contents
6000H	22H
6001H	00H
6002H	18H

The contents of location 6002 hex are changed to 18 hex=24 decimal

The RES instruction

This has the opposite effect to the SET instruction; it RESETS a bit in a byte or RAM location.

The BIT instruction

The bit instruction allows us to test for individual bits of a register or byte. The results of the test are signified by resetting or setting the Zero flag. If the bit tested was zero then the Zero flag would be set and, if not, the Zero flag would be reset.

```
BIT 7,A
```

The above instruction would read 'test BIT 7 of the A register'. Therefore, if the A register contained 128, which is 10000000 binary,

then the instruction would reset the Zero flag as bit seven is set to 1. If, however, we used the instruction:

BIT 0,A

with the same contents in the accumulator the Zero flag would be set, as bit 0 is zero.

The BIT instruction is very useful because it does not corrupt anything we are testing. Similar to the COMPARE instructions, it affects the bits in the flag registers only.

Spectrum INs and OUTs

The Z80 chip needs to interface to other devices such as the keyboard and a cassette recorder so that the user can communicate with the computer. There are two methods what we can use to communicate to these devices. One is known as memory mapping, that is PEEKing or POKEing, the other is by PORT addressing. A PORT is a gateway to these devices which can be read by using the instruction 'IN' or written to by using 'OUT'. There are 256 of these PORTS on the Sinclair computer. Most can be used by electronics buffs, to link up to devices such as speech synthesisers and sound chips.

There are two instructions in BASIC, 'IN' and 'OUT' which allow us to gain access to these ports. Frequently, these instructions are used to scan the keyboard or output to the speaker to produce noises.

The keyboard is divided into 8 rows of 5 keys each and the actual syntax of Spectrum BASIC to read the keyboard uses a two byte number. For example:

LET X=IN 61438

scans the keys 0 to 6 on the top row of the keyboard. The other addresses and the keys they scan are given below:

ADDRESS	HEX	KEYS SCANNED
32766	7FFE	SPACE,SYMBOL SHIFT,M,N,B
49150	BFFE	ENTER,L,K,J,H
57342	DFFE	P,O,I,U,Y
61438	EFEE	0,9,8,7,6
63486	F7FE	1,2,3,4,5
64510	FBFE	Q,W,E,R,T
65022	FDFF	A,S,D,F,G
65278	FEFE	CAPS SHIFT,Z,X,C,V

So if we wanted to scan for the bottom row of keys from SPACE to the letter B we would use the BASIC instruction:

```
LET X=IN 32766
```

A value is returned in the variable X depending on which keys are pressed. There are five bits which represent the state of each row on the keyboard. If a particular bit is low (i.e. 0) then this means that a key is depressed, and if no keys were depressed then all bits would be high. In the table above the key values have been given in bit order, so if we were scanning the keys 1 to 5 then bit 0 would indicate the state of the key 1, bit 1 the state of key 2 and so on.

The other 3 bits returned when scanning the keyboard are not used and are unpredictable (mainly because of the different models of spectrums available) so it is wise not to compare the values read unless you mask out the first five bits. Masking means removing bits according to a pattern. To mask the top three bits you would have to do RES 7, r RES 6, r and RES 5, r to set them all to zero.

The IN instruction

In machine code to read a PORT we use the instruction:

```
IN A,(port)
```

The value port is the PORT address which is a one byte number in the range 0-255. This port is read and the value is returned in the Accumulator. How do we use this instruction to scan the various lines on the keyboard? Well, if you look closely at the address which you scan in BASIC to read a particular row you will notice that the low bytes of each address are all FE hex, 254 decimal. The port address and the high bytes all differ from each other.

To read a set of keys in machine code we first LOAD the Accumulator with the high byte of the line we wish to read and then execute the instruction:

```
IN A,(0FEH)
```

So for example, if we wanted to scan the keys 0 to 6, we would write the following code:

```
LD A,0EFH ;SELECT LINE 0-6
IN A,(0FEH) ;READ PORT
```

Now, if we wanted to test if the key 0 was pressed we could use the BIT instruction:

BIT 0,A ;Test for "0"

This would set the Zero flag if the key was pressed or reset the flag if it was not pressed.

It is an easy matter to read a set of keys by using the compare instruction:

```
LD A,0FBH ;select keys Q to T
IN A,(0FEH) ;read key board port
AND 31 ;mask off lower 5 bits
```

This portion of code sets the Zero flag if all the keys Q,W and E are pressed.

There is another form of the IN instruction which allows us to specify the port by the value in the C register. The register in which the value is read can also be chosen from the set A,B,C,D,E,H or L.

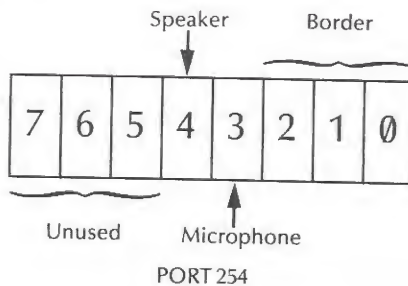
The code:

```
LD C,0FEH
LD A,0FBH
IN E,(C)
```

reads the line Q to T and places the value read in the E register.

The OUT instruction

The 'OUT' instruction is often used to generate sound and output to the cassette system. The port which controls the small pisa speaker is at the address 0FEh,254 decimal. This also has the ability to change the screen border colour. The values of output to this port is in the format shown below:



The first three bits (bits 0-2) of the byte are used for the border colour. Bit 3 is used to control the EAR and MIC sockets so that data

can be sent and read to and from a cassette unit. Bit 4 is used to pulse the (so-called!) speaker in the Spectrum.

Sound generation on the Spectrum is a simple matter of pulsing bit 4 of port 254 high and then low for a short period of time. We would set bit 4 (i.e. to 1) and hold it high for a short time and then set it low (0) and hold it at this level for the same period of time. The delay we have between 'flipping' bit 4 of the port determines the frequency or note we get from the speaker. A long delay produces a low frequency and a short delay a high frequency. To output the value in the Accumulator to a port we use the mnemonic:

```
OUT (addr),A
```

This simply reads 'OUTPUT the value in the Accumulator to the port address' So to turn the speaker on we would use the instructions:

```
LD    A,16
OUT  (0FEH),A
```

Notice how we first LOAD the Accumulator with 16. All this does is to set bit 4 high, which when sent to the port turns the speaker on. The following program demonstrates how the OUT instruction can work to generate sound. Both the assembler mnemonic listing and a BASIC listing have been given, with which the machine code can be loaded. Line 20 of the BASIC program changes the low bytes of the values for the duration and the frequency.

Assembler Listing

```

                ORG 28000D
                JP NOISE                ; JUMP TO MAKE A NOISE
SOUND:
                LD A,10H                ; MASK SPEAKER
                                                ; SO BIT 4 IS HIGH
                OUT (0FEH),A            ; TURN ON SPEAKER
                CALL DELAY              ; AND KEEP HIGH
                                                ; FOR A SHORT WHILE
                XOR A                    ; TURN BIT 4 OFF
                OUT (0FEH),A            ; TURN SPEAKER OFF
                CALL DELAY              ; AND KEEP IT OFF
                                                ; FOR A SHORT WHILE
                RET

```

```

DELAY:
    LD B, D                ;TRANSFER DE TO
                           ;BC REGISTER PAIR
    LD C, E                ;IE. PLACE DELAY IN
                           ;BC REGISTER

LOOP:
    DEC BC                 ;DECREMENT BC REGISTER PAIR
    LD A, B                ;AND REPEAT
    OR C
    JR NZ, LOOP           ;UNTIL BC PAIR IS ZERO
    RET                    ;RETURN AFTER FINISHING
                           ;DELAY

DELA EQU 100               ;DELAY
DURAT EQU 100             ;DURATION

NOISE:
    LD DE, DELA            ;GET DELAY
    LD HL, DURAT          ;GET DURATION
BUZZ:  CALL SOUND         ;MAKE A SOUND USING
                           ;DELAY AND DURATION
    DEC DE                 ;SUBTRACT ONE OFF DELAY
    DEC HL                 ;SUBTRACT ONE OFF DURATION
    LD A, L
    OR H
    JR NZ, BUZZ           ;REPEAT SOUND UNTIL
                           ;DURATION IS ZERO.

    RET

    END

```

BASIC Program Listing

```

10 FOR A=1 TO 100
20 POKE 28029,A: POKE 28026,A
30 FOR X=1 TO 20
40 RANDOMIZE USR 28000
50 NEXT X
60 NEXT A

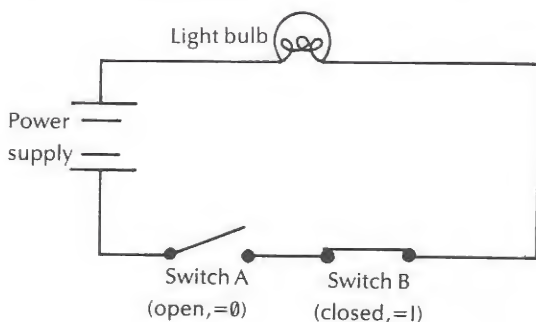
```

5 Operating logically

There are three logical operations available in the Z80 instruction set: AND, OR and XOR. These are all 8 bit operations which are best explained by looking at a series of diagrams and what are known as 'Truth Tables'. As explained in Chapter One numbers are represented in the computer by a series of 0's and 1's called bits and the Z80 groups 8 of these bits together to form a byte. Logical operations are performed on all 8 bits of a byte, and transform their values as described below.

AND Operations

The AND function operates on the corresponding bits of two bytes. If both corresponding bits were 1, then our result (another byte) after 'ANDing' these two bytes would set that bit to 1. If either or both of the bits were zero then the resulting bit would be zero. Look at the simple circuit diagram below:



This diagram has two switches, labeled A and B, a power supply and a light bulb. If we take the state of a closed switch as representing 1 and the open state as 0, then in order to switch the light on we have to have both switches on, i.e. then are both set to 1. We can represent the possible combination of the switches and their effective result on the light bulb by a truth table. The final result is 1 if the light bulb lights and 0 if it doesn't.

AND TRUTH TABLE

A	B	A AND B
0	0	0
0	1	0
1	0	0
1	1	1

As you can see Switch A and switch B have to be closed to make the light bulb light up.

When we do a logical operation in machine code all the operations act on the Accumulator. The AND operator is useful for picking up bits which we want to examine. This is known as 'masking'. If we had the code:

```
LD A,01001010B
AND 00000111B
```

Since the data is in binary we require the 'B' suffix after the numeric value. The second operation is the AND function. It takes each bit of the data and AND's it with the corresponding bit in the accumulator:

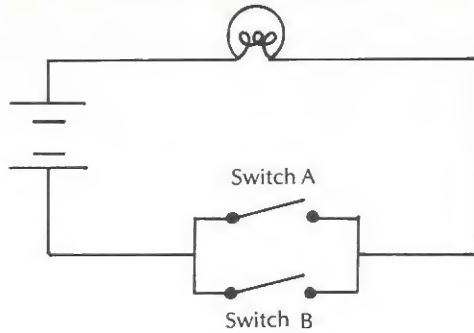
```
AND      01001010
         00000111
         -----
         00000010=2 decimal
```

AND instruction set

```
AND n    AND (HL)
AND A    AND (IX+dd)
AND B    AND (IY+dd)
AND C
AND D
AND E
AND H
AND L
```

OR Operations

The OR function is analogous to the circuit diagram below:



As you can see either one of the switches can be on to set the light on. Now looking at the possible combinations of switching in this circuit we get the following truth table:

OR TRUTH TABLE

A	B	A OR B
0	0	0
0	1	1
1	0	1
1	1	1

The OR operator is useful to set a series of bits in the Accumulator.

```
LD A,10101101B
OR 11100000B
```

The above code would set the top three bits in the Accumulator.

```
OR      10101101
        11100000
        -----
        11101101=237 decimal
```

OR instruction set

```

OR n    OR (HL)
OR A    OR (IX+dd)
OR B    OR (IY+dd)
OR C
OR D
OR E
OR H
OR L

```

XOR Operations

The XOR 'exclusive OR' operator is a little more difficult to explain with the aid of a circuit diagram but is just as easy to understand. It is similar to the OR operator but either, not both bits, may be high to give a high output

XOR TRUTH TABLE

A	B	A XOR B
0	0	0
0	1	1
1	0	1
1	1	0

The XOR operator is used to complement bits in the Accumulator and is sometimes known as *toggling* . What was previously on would be turned off, and what was off would be turned on. This instruction would be ideal for turning lights on and off connected to a computer. If, for example, we had the location at the address labelled LIGHT linked to some hardware which turned on a light if it contained 1 and turned the light off if it contained a zero, then the following program would generate a flashing strobe:

```

LD B,0      ;set delay
LD A,1      ;set state of switch
TOG: XOR 1   ;toggle switch
LD (LIGHT),A ;turn light on or off
DELAY: DJNZ DELAY ;short delay
JR TOG

```

The B register is loaded with zero which is used as a counter in a delay loop to hold the light on or off for a short period of time.

Because the DJNZ instruction will decrement B before testing it, giving B a start value of 0 causes 256 loops round the DJNZ.

```
LD A,1
```

The Accumulator is then set to 1 and toggled, leaving zero in the Accumulator since $1 \text{ XOR } 1 = 0$

```
TOG:  XOR 1
      LD (LIGHT),A
```

The LOAD instruction turns the light on or off according to the result in A, so first time round this would turn the light off. We now leave the light in this state for a period of time using a DJNZ instruction which counts from zero to 255 then back to zero again:

```
DELAY: DJNZ DELAY
```

We now come to the last instruction which transfers program control back to the address TOG:

```
JR TOG
```

This time, with A register containing zero the exclusive OR function will set the A register to one, producing a strobing effect.

XOR instruction set

```
XOR n  XOR (HL)
XOR A  XOR (IX+dd)
XOR B  XOR (IY+dd)
XOR C
XOR D
XOR E
XOR H
XOR L
```

6 Block manipulation

Block instructions give the Z80 the ability to move or compare blocks of data automatically or semi-automatically. A feature not found on any other 8 bit microprocessor on the market today.

A common use of block move instructions is to reduce screen flicker in games. A technique I will show you later. First, let us look at the block compare instructions. There are four instructions concerned with searching for a particular value ('key') in a block of data. To search for this key we can use any one of the following instructions:

CPIR
CPDR
CPI
CPD

CPIR

The A register is LOADED with the value which we are searching for (the 'key'). The HL register pair is LOADED with the address of the start of the block we wish to search, and the BC register pair is set up to contain the number of bytes we want to search through. The CPIR instruction is used to automatically go through all the data, comparing the contents at each address until it either finds the key it is searching for or until it has exhausted the search. This is signified by the BC register containing 0.

If the key is found then the Zero flag is set and the HL register pair points to the next address after the key. So the CPIR can be thought of as three instructions INC HL, CP (HL) and DEC BC.

Take a look at the following example:


```

STRING:          DEFM  "ABCDEFGH"
SEARCHFORWD:    LD    HL,STRING
                LD    BC,8
                LD    A,"G"
                CPIR
                JR    Z,FOUND
NOTFOUND:      .
                .
                .
FOUND:          DEC   HL

```

The first line contains the assembler psuedo operator `DEFM` 'Define Message' which tells the assembler to place the string "ABCDEFGH" in memory when the program is being assembled. As you can see the `HL` is LOADED with the start of the string and `BC` the number of bytes we wish to search through. The Accumulator contains the key "G", which we want to seek. The `CPIR` instruction will find this key (as it is contained within the string) and cause the program to jump to the address at the label `FOUND`. At this point we subtract the `HL` register by one to point to the actual address where the key is.

CPDR

The `CPDR` instruction is similar to the `CPIR` instruction but the `HL` register pair points to the end of the block of data and the search is made backwards. This time when a key is found the Zero flag is set as before, but the `HL` register pair will point to one less than the address where the key was found.

```

STRING:          DEFM  "ABCDEFGH"
SEARCHBACK:     LD    HL,STRING+8
                LD    BC,8
                LD    A,"G"
                CPIR
                JR    Z,FOUND
NOTFOUND:      .
                .
                .
FOUND:          INC   HL

```

As you can see, our code for searching for a key is similar to the last subroutine `SEARCHFORWD`. However we start off with LOADING the `HL` register pair with the address of the end of the string:

```
LD HL,STRING+8
```

This kind of instruction is allowed on most assemblers and all it does is to add the offset (+8) to the address of the label to get the resultant address. When the key is found it arrives at the label FOUND. However, this time the HL register pair will point to the character "F" so to correct this we amend the HL register pair with instruction:

```
FOUND:      INC  HL
```

CPI and CPD

These are known as semi-automatic instructions. If we use a CPI instruction then it will compare the A register with the contents of the HL register. The HL register pair will be incremented and the contents of the BC pair will be decremented. Flags will be set according to the result of the comparison and the subtraction of the BC pair. The two most significant flags to test are the PO flag (Parity Odd) and the Zero flag. If the key is found then the Zero flag is set. On the other hand if the search is exhausted and the key is not found then the P/V flag is set. These instructions are useful when we are searching through non-continuous data. For example, if we wanted to search through a string and the data we are seeking occurs every three bytes of the string, then we could use the following code:

```
LD  HL,STRING      ;SET OF START OF STRING
LD  BC,LENGTH      ;SET LENGTH
LD  A,KEY          ;KEY TO SEARCH FOR
LOOK: CPI          ;COMPARE (HL) WITH A REG
                          ;DECREMENT BC
                          ;AND INCREMENT HL
JP  Z,FOUND        ;FOUND KEY
JP  PO,NOTFOUND    ;EXHAUSTED SEARCH,NOT FOUND
INC HL             ;SKIP PASS
INC HL             ;UNWANTED DATA
JR  LOOK          ;KEEP SEARCHING
```

The routine will exit to the memory address specified by the label FOUND if the key is in the string or to that given by label NOTFOUND if the key is not in the string. Notice that there are only two INC HL instructions, not three. This is because the CPI instruction has already incremented the HL pair.

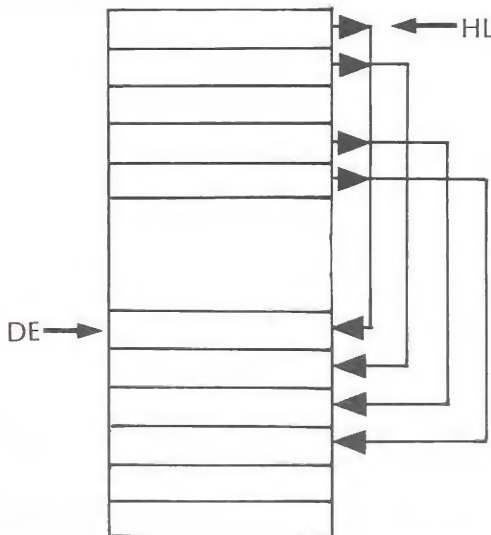
Block transfer

The Z80 is unique amongst 8 bit microprocessor chips in possessing a set of block transfer instructions. These allow blocks of data to be

moved around within memory utilising just four instructions. The HL register is LOADED with a 16 bit address which points to the start of the block to be moved. The BC register contains the numbers of bytes in that block which we wish to move. The DE register contains the destination address where the first byte of data is to be stored. After setting up these registers we could use any one of four block move instructions. Like the Block Compare set of instructions the Block Transfer possibilities allow for two automatic and two semi-automatic instructions. The two instructions LDIR (LOAD INCREMENT and REPEAT) and LDDR (LOAD DECREMENT and REPEAT) are the two automatic block instructions. They allow us to move whole blocks of memory simply by executing the instruction once.

LDIR

The LDIR instruction is used to move data which is held in a continuous sequence of memory locations. The HL registers are set up to point to the start of the data block, the DE pair is set up to point to the start of the destination and the BC register the number of bytes to move. When executing the LDIR instruction the contents of the location pointed to by the HL pair is copied to the location pointed to by the DE pair. Both the HL and the DE register pairs are incremented to point to the new data and destination locations while the BC register pair is decremented. This transfer continues until the BC register pair reaches zero, then the Z80 goes onto the next instruction.

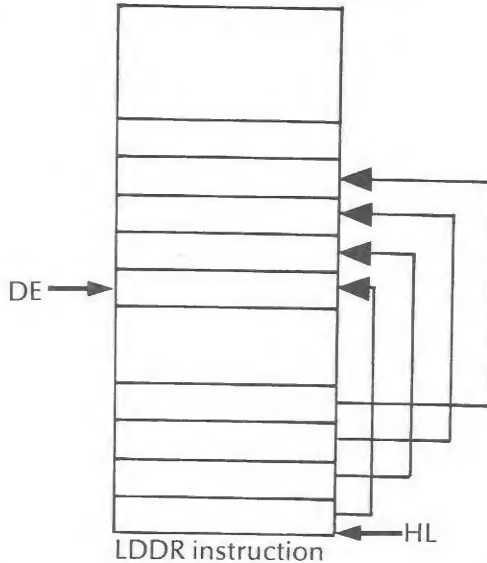


Block move instruction LDIR instruction

LDDR

The LDDR instruction is similar to the LDIR but the HL register pair and

DE register pair point to the end of the block and destination addresses. The transfers are made as with the LDIR instruction but the HL and DE pairs are decremented to point to the new data and destination addresses. Again the BC is decremented and the transfer continues until the BC register has reached zero.



Most games programmers use the block move instructions to move vast amounts of data to the screen. When a lot of information is needed to be drawn to the screen this can result in the TV display flickering. To reduce this flicker it is possible to draw the data on a dummy screen unseen by the player. This dummy screen can then be moved to the actual screen using the block move instruction.

If you imagine that we have set up a screen full of data at the address C000 hex and we wish to move it to the screen which is at the address 16384 or 4000 hex. We might use the following code:

```
;set up dummy screen
```

```
LD    HL,C000H ;HL points to dummy screen
LD    DE,4000H ;DE points to real screen
LD    BC,1B00H ;BC contains number of bytes
LDIR                    ;move it!
```

The HL register points to the start of the dummy screen which is C000 hex and the DE is set up to point to the start of screen. The BC register is set up to contain 1B00 hex or 6912 decimal, the number of bytes

contained in the display file. When we execute the LDIR instruction it moves 6912 bytes starting from the location C000 hex to the screen, reducing screen flicker to a minimum.

The two semi-automatic instructions LDI (LOAD and INCREMENT) and LDD (LOAD and DECREMENT) are used similarly to the CPI and CPD instructions when the data is non-continuous or we wish to stop moving data on certain conditions.

The parity odd flag is affected by the two instructions indicating that BC has reached zero when the PARITY ODD flag is set. If we wanted to write a routine which would move a block of data to the screen from C000 hex. until we reach a zero byte, then the following code could be used:

```

                LD    HL,C000H ;point to start address
                LD    DE,4000H ;point to destination address
                LD    BC,1B00H ;maximum number of bytes to move
MOVE:          LDI                    ;move one byte (HL)——>(DE)
                ;DE=DE+1:HL=HL+1:BC=BC-1
                RET    PO              ;PARITY ODD FLAG set, all done
                LD    A,(HL)          ;get next contents
                AND   A                ;test for zero
                JR    NZ,MOVE
                RET                    ;return we have reached a zero!

```

The routine will move at least one byte as the test for a zero byte is made after the LDI instruction. The transfer is complete if the PARITY ODD flag is set, indicating that we did not encounter a zero byte in the dummy screen and the BC register reached zero. It will exit when we reach the first zero byte. The AND A is used here to test for zero. A more obvious method would be CP0 (COMPARE with zero) but AND A is more efficient as it is faster to execute and uses less memory space. The AND A will leave the contents of A unchanged since any bit that is AND'ed with itself will remain unchanged (see the Truth Table for AND in chapter 5). The AND instruction will set the flags according to the eventual contents of A so this is an easy way of setting flags. OR A would be equally suitable for this purpose but XOR A would clear the A register to zero.

Miscellaneous instructions

The next (and last) batch of five instructions that will be explained in this book all operate on the Accumulator or flag register, so they have been grouped

CPL Complement accumulator

The Complement instruction simply replaces 0's for 1's and vice versa.

For example:

```
LD A,187 (10111011 binary)
CPL
```

The A register will contain 68 (01000100 binary) after executing the Complement instruction.

NEG Negate accumulator

The Negate instruction has the effect of multiplying the number by -1. It changes the number's sign (not just the sign bit!)

For example:

24 hex (36 decimal) becomes DC hex (-36 decimal)

This instruction performs the 'two's complement' on the contents of the Accumulator. It is directly equivalent to the pair of instructions:

```
CPL
INC A
```

CCF Complement Carry Flag

This changes the Carry flag to a 1 if it was a 0 and vice versa.

SCF Set Carry Flag

This instruction forces the Carry flag to a 1.

DAA Decimal Adjust Accumulator

This instruction is used to add numbers which are represented in Binary Coded Decimal (BCD) form. The decimal numbers 0 to 99 can be represented in one byte by splitting it into two sets of 4 bits each, called nybbles. The left nybble is the number of tens in the number and the right nybble represents the number of units. For example, the number 29 can be represented by the BCD number 0010 1001 (the 8 bit binary number has been split into two nybbles to make it easier to read.)

When we want to add or subtract two BCD numbers we use the normal ADD or SUB instructions followed by the DAA instruction. The H and N flags are used by the DAA to adjust the result to BCD.

For example:

```

LD    A,29H    ;LOAD A with 29 hex 41 decimal 29 BCD
LD    B,24H    ;LOAD B with 24 hex 36 decimal 24 BCD
ADD   A,B      ;ADD B register to A register
DAA                   ;decimal adjust

```

This piece of code would leave the result 53 hex in the A register (not 4D hex as with normal addition).

The way that DAA works is that after an arithmetic operation it checks whether the low nybble is in the range 0 to 9. If this is not the case, it will add 6 to the low nybble, which causes the high nybble to be incremented. Then the high nybble is checked. If it exceeds 9 then 6 is added to the high nybble, which will overflow into the Carry flag.

In the example the Accumulator will hold the value 4DH before DAA is executed.

Here is the flow of logic for DAA in this case:

(a) Low nybble=DH

(b) This is greater than 9 so add 6 to the low nybble to give temp. result of 13H.

(c) replace the low nybble of temp. result to leave 10H or the high nybble equal to 1.

(d) Add the high nybble of the temp. result to the high nybble of the accumulator to give:

$$4+1=5$$

(e) 5 is less than 9 so replace high nybble in accumulator.

(f) The Accumulator now contains 53H, this is the correct BCD result of 29+24

7 A Spectrum monitor

This chapter presents a program which will allow you to write, run and debug machine code programs. It can be entered into your Spectrum using the BASIC machine code loader given as Listing 1 below. After keying in and SAVING the BASIC program on tape, begin the program by entering the start address at which the machine code program will start to be built up. Then INPUT the hexadecimal data which makes up the program as given in the hexadecimal listing (Listing 2). If at any time you wish to correct a mistake there are edit facilities to help you (see below). Typing \$\$ when prompted for a hexadecimal byte allows you to change the address at which the next piece of data is to be placed. To quit the program, type in a double hash ## when prompted for hexadecimal data. After this has been done the program allows you to store the machine code using the SAVE command.

For anyone interested in the way in which the monitor was written the corresponding assembly mnemonic listing has been included in Appendix 3. At the moment don't worry about understanding how it works, just type in the data. Once it is up and working you can use it to enter the other machine code programs which are given. In this book each functional program, other than illustrative examples, has two listings. One is in mnemonic form which is easier to read and follow. This can be used by those of you that have full assembler programs available. The second listing is the hexadecimal equivalent. This is the portion of memory of your Spectrum which holds the program. It is displayed as a hexadecimal dump and can be reproduced by keying in the appropriate value for each memory location using the hex monitor. You could, of course, use the BASIC monitor to input the machine code listings. However, as you will see this is not as powerful as its machine code counterpart.

The monitor program has been assembled at the address 25500. This allows us to write machine code programs higher up in memory, giving more free space in which to RUN both BASIC and machine programs. The monitor offers the machine code programmer eleven functions. These are: Dump, Edit, Fill, Goto, Hunt, Identify, Load, Move, Print, Register and Save.

More commands can easily be added by changing a command table to point to the routine which deals with the new command.

After the monitor has been typed in using the BASIC loader the machine code can be *SAVED* by typing:

```
SAVE "SMON" CODE 25500, 1500.
```

The Spectrum is then cleared by switching the machine off and on again. Next the monitor can be loaded by typing in:

```
CLEAR 24999:LOAD "" CODE:RANDOMIZE USR 25500.
```

The monitor should then welcome you with a '>' prompt, inviting the input of one of the eleven commands.

*D*address (Hexadecimal dump)

Type in 'D' followed by a two byte hexadecimal number. The monitor then displays the contents of memory from the given address in a hexadecimal format. The routine will keep dumping the memory contents until a key is input other than a carriage return. For example:

```
>D 0100
```

The above command will display 64 bytes of memory from the address 0100 hex.

*E*address (Edit)

This allows you to *edit* or modify a byte in memory. To execute the command, type in 'E' followed by a two byte address which you wish to start modifying. The monitor will then show the address which is being modified and the contents of that address. Then type in a one byte hexadecimal number to change that location. After the modification has been given the monitor will automatically go onto the next location to be edited. The routine can be exited by typing in a non-valid hexadecimal digit. eg:

```
>E C000
C000 FF 3E
C001 00 2A
C002 00 C9
C003 00 <ENTER>
```

Pressing <ENTER> will exit from the edit command.

F start address end address byte (FILL memory with byte value)

To FILL a block of memory with a given byte, given the start address, end address and byte value to where the byte is to be filled. For example:

```
>F 4000 5800 2A
```

This will fill the memory from 4000 hex to 5800 hex with the byte 2A hex.

G address or *G* address, breakpoint address (GOTO address)

This command allows execution of a portion of code from a given address. A second parameter can be given which allows you to give a break point where the register values will be displayed.

To give a breakpoint type in a ',' after the first address and then type in the breakpoint address.

```
>G C000
```

The above example will cause the monitor to execute from address C000 and will RETURN back to the monitor after a RET instruction is met.

```
>G C000 ,C004
```

This second example will cause the monitor to execute from the address C000 with a breakpoint at C004. If the code flows through this address then it will return to the monitor displaying the Breakpoint address and the contents of the registers. Such an example could be as follows:

```
*C004
```

AF	BC	DE	HL	IX
2A3E	22AA	0000	DEF0	DDFE
0000	0000	2232	2312	

H start address end address byte value (HUNT for a byte)

The HUNT command will allow you to search for a specific byte through a given set of addresses. Type in 'H' and then give the start address, the end address and the byte value for which you wish to search. The routine will then display each address where that byte is found, pausing for you to type in ENTER. To exit from this routine before the search is exhausted, any other key may be pressed.

```
>H 0000 0100 2A
0008
0018
0031
003A
0068
007A
>
```

The above example will search for the byte 2A hex. From the address 0000 to 0100 hex. When the search is exhausted the monitor returns with the prompt.

I string (Identify file name)

This command is used in conjunction with the SAVE and LOAD command identifying the file name to be LOADED or SAVED.

Type in 'I' and then input a filename consisting of no more than ten letters of the alphabet. Lower case letters will be ignored.

```
>I SPECMON
```

The example above will set the identifier to the string 'SPECMON'.

L start address, number of bytes (Load file)

The LOAD command will wait for the filename given by the identifier (see the 'I' command) and once found on the tape will start to LOAD it at the given address. The second parameter (also a two byte number) specifies the number of bytes to LOAD. The LOAD command will read in each file of the tape and display its header to the screen.

```
>I SPECMON
>L C000 0100
Waiting for SPECMON
```

The commands above will LOAD the file 'SPECMON' to the address C000 hex. The number of bytes to be LOADED is 256 (i.e. 0100 hex). To exit from the LOAD command at any time press both the <CAPS SHIFT> and the <BREAK> keys.

M start address, end address, destination address (MOVE Block)

This command will MOVE blocks of data to a given address. You need to specify the start address, end address and the destination address.

```
>M C000 DB00 4000
```

The above example will MOVE data from the address C000 up to DB00 hex to the address 4000 hex.

Paddress (*PRINT ASCII*)

The PRINT command is similar to the DUMP command except that it displays the contents of the memory in ASCII code form rather than hexadecimal. For example,

```
>P 0690
```

will PRINT the ASCII contents from the address 0690 hex. To continue the listing press <ENTER>, otherwise press any other key.

R or **R r** or **R 'r** (*REGISTER modify*)

The REGISTER command allows you to examine or modify any register r, where r can be any of the following:

- A modifies the AF register pair
- B modifies the BC register pair
- D modifies the DE register pair
- H modifies the HL register pair
- X modifies the IX register pair

To examine the contents of the registers type in the command 'R' and the <ENTER>.

```
>R ENTER
```

AF	BC	DE	HL	IX
FF3E	0000	0000	F22A	0000
0000	0000	0000	002A	

The second value under each register pair is the contents of the alternative register set. The IX register is not shown as this is used by the BASIC system.

To modify a register, you can simply type in the register pair you wish to modify, after typing in the command 'R'. The monitor will then display the current value of the register pair, waiting for an INPUT of the new value. To modify the alternative set, type a ' before entering the register pair. For example:

```
>R H 002A 2FFF
>R 'D 0000 2A33
```

This will change the contents of HL register pair to the value 2FFF hex, and the alternate DE register pair to 2A33 hex.

S start address ,number of bytes (SAVE)

The SAVE command is used to SAVE machine code to a tape recorder. The start address and the number of bytes to SAVE are the parameters that you will need to specify with this command. The monitor will then prompt you to get the tape recorder ready and SAVE the portion of code on tape. The file name is set by using the 'I' (Identify) command. For example:

```
>I FRED
>S 4000 1B00
Press any key when ready
```

These will SAVE the 6912 bytes from the address 16384 to the tape. The file will be SAVED as 'FRED'.

The listings follow below.

Listing 1: BASIC Hex Monitor

```
10 CLEAR 25499
20 CLS : GO SUB 160
30 LET a=x
40 GO SUB 320: PRINT x$;" ";
50 FOR z=1 TO 8: GO SUB 90: IF
a$="$$" THEN GO TO 10
60 LET a=a+1: NEXT z
70 PRINT
80 LET x=a: GO TO 40
90 INPUT "hex :"; LINE a$: IF
LEN a$<>2 THEN GO TO 90
100 IF a$="$$" THEN RETURN
110 IF a$="##" THEN STOP
120 GO SUB 250
130 IF e=1 THEN GO TO 90
140 POKE a,x: PRINT a$;" ";
150 RETURN
160 INPUT "addr:"; LINE b$: IF
LEN b$<>4 THEN GO TO 160
170 GO SUB 250
180 IF e=1 THEN GO TO 160
190 LET x=t*256+x: LET a=x
200 RETURN
210 REM two byte hex input
220 LET a$=b$(1 TO 2)
```

```

230 GO SUB 250: LET t=x
240 LET a$=b$(3 TO 4)
250 REM one byte hex input
260 LET e=0
270 LET l=FN x(2): IF l>15 THEN
LET e=1
280 LET h=FN x(1): IF h>15 THEN
LET e=1
290 LET x=h*16+l
300 RETURN
310 DEF FN x(n)=CODE a$(n)-48-(
CODE a$(n)>57 AND CODE a$(n)<71)
*7-(CODE a$(n)>96 AND CODE a$(n)
<103)*39
320 REM two byte input
330 LET h=INT (x/256): LET l=x-
h*256
340 LET x=h: GO SUB 370
350 LET x=l: GO SUB 380
360 RETURN
370 LET x$=""
380 LET p=INT (x/16): GO SUB 39
0:LET p=x-(INT (x/16))*16
390 REM hex
400 IF p>9 THEN LET a$=CHR$ (p+
CODE "A"-10)
410 IF p<=9 THEN LET a$=CHR$ (p
+CODE "0")
420 LET x$=x$+a$
430 RETURN

```

Listing 2: Spectrum Monitor Hexadecimal Listing

639C	C3	36	64	3E	FF	CD	C2	04
63A4	C9	37	3E	FF	CD	56	05	C9
63AC	11	11	00	DD	21	06	69	AF
63B4	37	CD	56	05	3A	11	69	4F
63BC	3E	24	32	11	69	CD	CD	68
63C4	11	07	69	CD	2C	64	3E	20
63CC	CD	22	64	EB	71	23	23	23
63D4	7E	CD	92	66	2B	7E	CD	92

63DC	66	2B	3E	20	CD	22	64	7E
63E4	CD	92	66	2B	7E	CD	92	66
63EC	CD	CD	68	C9	11	11	00	DD
63F4	21	06	69	AF	CD	C2	04	C9
63FC	E5	C5	D5	CD	A7	64	3A	3B
6404	5C	CB	6F	28	F9	CB	AF	32
640C	3B	5C	D1	C1	E1	C9	CD	FC
6414	63	3A	08	5C	CD	22	64	C9
641C	3E	02	CD	01	16	C9	F5	F5
6424	AF	32	8C	5C	F1	D7	F1	C9
642C	1A	FE	24	C8	CD	22	64	13
6434	18	F6	31	6E	69	CD	1C	64
643C	11	B0	65	CD	2C	64	11	44
6444	65	CD	2C	64	CD	CD	68	31
644C	6E	69	3E	08	32	6A	5C	21
6454	6F	69	36	4B	23	36	64	2B
645C	22	3D	5C	21	5F	64	E5	CD
6464	CD	68	3E	3E	CD	22	64	CD
646C	12	64	D6	41	D8	FE	13	D0
6474	87	21	81	64	5F	16	00	19
647C	5E	23	56	EB	E9	83	66	83
6484	66	83	66	1A	66	B3	64	E1
648C	66	17	67	8F	68	5E	68	83
6494	66	83	66	FC	64	33	68	83
649C	66	83	66	5A	66	83	66	E2
64A4	67	B5	65	01	00	80	11	00
64AC	40	21	00	40	ED	B0	C9	CD
64B4	EE	64	CD	8D	66	3E	20	CD
64BC	22	64	7E	CD	92	66	3E	20
64C4	CD	22	64	E5	CD	B1	66	E1
64CC	77	23	CD	CD	68	18	E0	CD
64D4	D6	66	E5	3E	20	CD	22	64
64DC	CD	D6	66	E5	D1	E1	C9	CD
64E4	D3	64	E5	D5	CD	EE	64	D1
64EC	E1	C9	3E	20	CD	22	64	CD
64F4	D6	66	E5	C1	CD	CD	68	C9
64FC	3E	20	CD	22	64	CD	D3	64
6504	22	40	65	7B	B2	CA	AE	66
650C	ED	53	42	65	11	7D	65	CD
6514	2C	64	11	17	69	CD	2C	64

651C	CD	AC	63	11	07	69	21	17
6524	69	06	0A	1A	CB	AF	4E	CB
652C	A9	B9	20	EC	23	13	10	F3
6534	ED	5B	42	65	DD	2A	40	65
653C	CD	A5	63	C9	00	00	00	00
6544	0D	2A	53	42	55	47	2A	20
654C	28	43	29	20	4A	6F	68	6E
6554	20	57	69	6C	73	6F	6E	20
655C	31	39	38	34	2E	0D	24	0D
6564	50	72	65	73	73	20	61	6E
656C	79	20	6B	65	79	20	77	68
6574	65	6E	20	72	65	61	64	79
657C	24	0D	57	61	69	74	69	6E
6584	67	20	66	6F	72	20	24	0D
658C	52	4F	55	54	49	4E	45	20
6594	4E	4F	54	20	49	4D	50	4C
659C	45	4D	45	4E	54	45	44	24
65A4	0D	2A	2A	45	52	52	4F	52
65AC	2A	2A	0D	24	16	01	01	0D
65B4	24	3E	20	CD	22	64	CD	D3
65BC	64	22	13	69	7B	B2	CA	AE
65C4	66	ED	53	11	69	11	63	65
65CC	CD	2C	64	CD	A7	64	CD	A7
65D4	64	CD	A7	64	CD	12	64	3E
65DC	03	11	06	69	21	17	69	12
65E4	13	01	0A	00	ED	B0	CD	F0
65EC	63	CD	A7	64	CD	A7	64	CD
65F4	A7	64	DD	2A	13	69	ED	5B
65FC	11	69	CD	9F	63	C9	E5	CD
6604	92	66	3E	20	CD	22	64	E1
660C	C9	06	08	7E	CD	02	66	23
6614	10	F9	CD	CD	68	C9	3E	20
661C	CD	22	64	CD	EE	64	0E	08
6624	CD	8D	66	3E	20	CD	22	64
662C	CD	22	64	CD	0D	66	0D	20
6634	EF	CD	CD	68	CD	CD	68	CD
663C	12	64	FE	0D	28	E0	C9	06
6644	15	7E	FE	20	38	04	FE	80
664C	38	02	3E	2E	CD	22	64	23
6654	10	EF	CD	CD	68	C9	3E	20

665C	CD	22	64	CD	EE	64	0E	08
6664	CD	8D	66	3E	20	CD	22	64
666C	CD	22	64	CD	43	66	0D	20
6674	EF	CD	CD	68	CD	CD	68	CD
667C	12	64	FE	0D	28	E0	C9	00
6684	D5	11	8B	65	CD	2C	64	D1
668C	C9	7C	CD	92	66	7D	5F	CB
6694	3F	CB	3F	CB	3F	CB	3F	CD
669C	A1	66	7B	E6	0F	C6	30	FE
66A4	3A	FA	AA	66	C6	07	CD	22
66AC	64	C9	C3	4B	64	CD	12	64
66B4	CD	C8	66	5F	CD	12	64	CD
66BC	C8	66	CB	23	CB	23	CB	23
66C4	CB	23	B3	C9	A7	DE	30	FE
66CC	0A	D8	A7	DE	07	FE	10	30
66D4	D9	C9	CD	B1	66	F5	CD	B1
66DC	66	6F	F1	67	C9	3E	20	CD
66E4	22	64	CD	D3	64	3E	20	CD
66EC	22	64	E5	EB	A7	ED	52	DA
66F4	AE	66	CA	AE	66	E5	C1	E1
66FC	E5	D1	13	E5	D5	CD	B1	66
6704	D1	E1	77	ED	B0	CD	CD	68
670C	C9	D1	21	4A	67	E5	D5	CD
6714	64	67	C9	3E	20	CD	22	64
671C	CD	D6	66	E5	3E	20	CD	22
6724	64	CD	12	64	FE	0D	28	E1
672C	FE	2C	C2	AE	66	CD	D6	66
6734	E5	11	01	69	01	03	00	ED
673C	B0	E1	36	CD	23	36	7E	23
6744	36	67	CD	64	67	C9	ED	73
674C	04	69	31	01	69	08	D9	E5
6754	D5	C5	F5	D9	08	DD	E5	E5
675C	D5	C5	F5	ED	7B	04	69	C9
6764	ED	73	04	69	31	EF	68	F1
676C	C1	D1	E1	DD	E1	08	D9	F1
6774	C1	D1	E1	08	D9	ED	7B	04
677C	69	C9	CD	4A	67	E1	2B	2B
6784	2B	CD	CD	68	3E	2A	CD	22
678C	64	CD	8D	66	EB	21	01	69
6794	01	03	00	ED	B0	CD	B4	67

679C	C9	06	04	5E	23	56	E5	EB
67A4	CD	8D	66	3E	20	CD	22	64
67AC	CD	22	64	E1	23	10	EC	C9
67B4	CD	CD	68	11	D3	68	CD	2C
67BC	64	21	EF	68	CD	9D	67	CD
67C4	CD	67	CD	CD	68	CD	9D	67
67CC	C9	7E	F5	23	7E	CD	92	66
67D4	F1	CD	92	66	23	3E	20	CD
67DC	22	64	CD	22	64	C9	3E	20
67E4	CD	22	64	CD	12	64	FE	27
67EC	20	05	CD	12	64	C6	08	21
67F4	2A	68	01	09	00	ED	B1	C2
67FC	B4	67	2B	11	2A	68	A7	ED
6804	52	11	EF	68	CB	25	19	23
680C	3E	20	CD	22	64	7E	CD	92
6814	66	2B	7E	CD	92	66	23	3E
681C	20	CD	22	64	CD	B1	66	77
6824	2B	CD	B1	66	77	C9	41	42
682C	44	48	58	49	4A	4C	50	3E
6834	20	CD	22	64	CD	E3	64	E5
683C	A7	ED	52	30	04	E1	EB	18
6844	F6	EB	C5	D5	C1	D1	F1	E5
684C	A7	ED	52	E1	38	03	ED	B0
6854	C9	09	2B	EB	09	2B	EB	ED
685C	B8	C9	3E	20	CD	22	64	CD
6864	12	64	FE	00	C8	FE	41	DA
686C	AE	66	21	17	69	06	0A	0E
6874	20	71	23	10	FC	21	17	69
687C	06	09	77	05	C8	23	CD	12
6884	64	FE	0D	C8	FE	41	DA	AE
688C	66	18	EF	3E	20	CD	22	64
6894	CD	D3	64	E5	EB	A7	ED	52
689C	DA	AE	66	CA	AE	66	E5	C1
68A4	E1	3E	20	CD	22	64	E5	D5
68AC	CD	B1	66	D1	E1	BE	F5	20
68B4	0D	CD	CD	68	CD	8D	66	CD
68BC	12	64	FE	0D	20	09	23	0B
68C4	78	B1	28	03	F1	18	E6	F1
68CC	C9	3E	0D	CD	22	64	C9	41
68D4	46	20	20	20	20	42	43	20

8 Program production

In this chapter we will go step-by-step through an example machine code program looking at the different sets of instructions and how they are assembled into machine code by the assembler. You may remember that in the first chapter we explained that there are several ways of writing machine code programs. One method is to use a professional assembler such as HI-SOFT'S 'DEVPAC'. Another is to use a monitor which allows us to INPUT machine code by its hexadecimal values.

Using an assembler is the best method of writing machine code as it is written in the mnemonic type instructions which are so easy to learn. When a program is in the process of being assembled the assembler goes through each mnemonic converting it to its machine code equivalent into another part of memory. Most assemblers on the market provide the option of seeing the mnemonic files translated to their hexadecimal equivalents and the addresses where each particular instruction is to be stored in memory.

Below is an example of an assembled listing:

```
                                ORG  16384
4000    3E 21    LD    A,33
4002    C9     RET
                                .END
```

The first instruction 'ORG' (ORIGIN) is not a Z80 instruction but is used to tell the assembler at what address to place the first instruction in memory. Our example shows that the origin is set at 16384 decimal so that the first instruction will be assembled at 16384 or 4000 hex.

The assembled listing shows mnemonics to the right of the listing and the address and hexadecimal Op codes to the left. Why are hexadecimal values displayed and not decimals? Well why not? Hexadecimal numbers are used for convenience sake only. They use fewer digits to represent the decimal numbers 0-255 and are easier to read . . . well they should be easier to read after a bit of practice!

The first address is 4000 hex which is 16384 decimal. At this address the Op code for LD A, (LOAD the A register) is placed. The

content of the address 4001 hex (16385) is the data byte 33. This is shown by its hexadecimal equivalent 21 hex. Since the first instruction was two bytes long the next instruction will be placed at the address 4002 hex or 16386 decimal. This is the RETURN instruction which is only one byte long and has the value C9 hex or 201 decimal.

The last instruction 'END' is also not a Z80 instruction. This is used by most assemblers to signify that there are no more instructions to assemble. Now armed with this information we can now go through a short assembled listing looking at each instruction and the effect it has. Also listed is a hexadecimal dump of the program which can be used with the BASIC monitor or machine code monitor listed in this book. Instruction will be given on how to INPUT the machine code using the BASIC monitor. First of all let's go through the overall effect of the program.

The machine code routine enables you to enhance your programs by having a scrolling attribute BORDER along the edges of the screen. A coloured BORDER is produced along the screen by POKING the attribute file with random PAPER colours. Then the routine begins to move the whole BORDER in a clockwise direction. The whole program comes in two parts: a BASIC listing (listing 1) and the machine code (listing 2). The BASIC program does the easy work. It draws a random coloured BORDER along the Spectrum screen. Line 10 sets the variable x to the address of the start of the attribute file 22528 and this is used to place a random line of PAPER colours along the top of the screen. Lines 20 to 40 produce the top BORDER by POKING a random PAPER colour (The PAPER colour is produced by multiplying a random number from 0 to 7 by 8).

The lines 50 to 100 produce two coloured BORDERS along the side of the screen and finally, lines 110 to 140 produce the BORDER for the bottom of the screen. After the BORDER is produced the machine code routine is executed by the BASIC line 150 'RANDOMIZE USR 30000'. The USR command is used by BASIC to call a machine code routine in memory. The address following the USR instruction is the address to where BASIC will jump. It executes our machine routine which will shift the whole BORDER clockwise by one attribute. After the machine code routine has been executed and returns back to BASIC (by using a RET instruction) the line following the BASIC call will be executed. Line 160 is used to slow down the scroll by pausing for 1 second. Line 170 jumps back to 150 to call the machine code routine again and again until the user breaks out by pressing <CAPS SHIFT> and <BREAK SPACE> together.

Type in the BASIC program and SAVE it by typing SAVE "DEMO". DO NOT RUN it at this point, as we have not yet typed in the machine code routine. RUNNING the program will probably result in the Spectrum crashing and losing the BASIC program!

With the BASIC program safely SAVED on a cassette, you should now key-in the assembler listing on page 91. To do this you will, of course, have to LOAD your assembler first. When this listing has been entered into the Spectrum you should SAVE the source code, again onto a cassette, using the appropriate command for your assembler. Next, CLEAR the memory of the machine by switching it off and then on again.

Then LOAD up the BASIC 'DEMO'. Before re-LOADING the machine code program type in CLEAR 29000. This will re-set RAMTOP to protect our machine code program. LOAD the machine code routine from the tape into the Spectrum by typing LOAD "DEMOC" CODE and <ENTER>. After it has LOADED, RUN the program and we should get a coloured BORDER scrolling around the screen in a clockwise direction. If the program does not scroll the border then probably you have mis-typed the machine code routine. Re-type the routine again using the BASIC monitor.

Now let's look more closely at the machine code routine and see how it works:

```
ORG 30000
```

The first line tells the assembler where the origin of the machine code is to be assembled. Since our program is to be placed at the address 30000, this number is placed after the ORG instruction.

```
ATTRADD EQU 5800H
```

The next instruction is like the ORG instruction in that it is not a Z80 one. It is used by the assembler to produce a table of strings (symbols) which hold one or two byte numbers. The string 'EQU' stands for EQUATE and produces the string and gives it the value following the EQUATE. In this way the above line will produce the symbol ATTRADD with the number 5800 hex (22528 decimal).

The symbol ATTRADD now holds the address of the attribute file.

```
LD HL,ATTRADD+31 ;POINT TO RIGHT HAND SIDE OF ATTR
LD DE,ATTRADD+30 ;DE POINTS TO THE NEXT ATTR
```

The next two instructions are at the start of the machine code program. These use the value ATTRADD contained in the assembler's symbol table. The HL register pair is LOADED with the address of 5800H

plus 31 decimal and the DE pair is LOADED with the address 5800H+ 30 decimal. The assembler will automatically calculate the two results and place the address in the source output.

The HL register pair contains the address of the far right hand corner of the top row of the attribute file. The DE register pair holds the address of the attribute location to the left of the top right hand corner (it points to the left of the HL register pair).

The object of the first portion of code is to move each of the 31 attributes of the top line of the screen one character along to the right. This is done by repeatedly replacing the attribute byte pointed to by the DE pair and placing it in the location pointed to by the HL pair. Then make the two pointers point to the next locations.

The loop DOTOP, therefore, is used to scroll the top attribute line from the left to the right:

```
LD    A,(HL)    ;SAVE FIRST ATTRIBUTE
PUSH AF        ;ON THE STACK
```

Before going into this loop we have to SAVE the attribute in the top right hand corner. This will be over-written with the new attribute to its left.

The first part of the code LD A,(HL) LOADS the contents of the address pointed to by the HL pair. This is the top right hand corner of the attribute file and places the value into the A register. The second instruction PUSH AF pushes this attribute onto the stack, where it will stay until we need it. This has the effect of saving the first attribute onto the stack.

```
LD B,31        ;LOAD B WITH COLUMN COUNTER
```

We now set up a loop counter needed to scroll the top line of the attribute file. This is done by LOADING the B register with the value 31 decimal, which is the number of characters we have to scroll across.

```
DOTOP: LD    A,(DE)    ;GET NEXT ATTRIBUTE
        LD    (HL),A   ;AND PLACE IT TO THE ATTRIBUTE
                        ;TO THE LEFT
        DEC  DE        ;POINT TO THE NEXT ATTRIBUTE
        DEC  HL
        DJNZ DOTOP    ;REPEAT UNTIL ALL COLUMNS DONE
```

The content of the DE register pair is LOADED into the A register with

the instruction LD A, (DE). This instruction is used to force the new byte to be scrolled to the right. This new attribute is placed to the right by LOADING it to the address pointed to by the HL pair, performed by the instruction LD (HL),A (LOAD into the address pointed by the HL pair the contents of the A register). The attribute, pointers are then moved to the left by one attribute by subtracting one from each of the pointers. The DECREMENT instructions DEC DE and DEC HL are used to implement this. The last instruction of this portion of code is the branch instruction DJNZ 'DECREMENT and Jump if Not Zero'. This instruction takes the B register from our column counter and subtracts one from it. If the result is not zero (if the B register does not contain a zero) then a relative jump is made to the address DOTOP. Remember, a relative jump differs from an absolute jump in that an offset number of bytes is given to where the program must jump instead of a two byte address. The assembler has the job of calculating this offset. It does this by working out the number of bytes between the label DOTOP and the instruction DJNZ DOTOP.

After leaving this loop the program has to deal with the coloured borders along the lefthand side of the attribute file. Since the HL pair has been decremented 31 times in the loop it will now point to the start of the attribute file 5800 hex. The DE pair will point to the address 57FFH:

```
LD    BC,32    ;LET HL POINT ONE ROW DOWN
ADD   HL,BC    ;BY ADDING 32
INC   DE      ;DE NOW POINTS TO START
                ;OF ATTRIBUTE FILE IE. 5800H
```

The HL pair is adjusted to point to the second attribute row by the adding of 32. This is done by the two instructions LD BC, 32 (LOAD the BC pair with 32) and ADD HL, BC (ADD to the HL pair the contents of the BC pair). The DE is adjusted to point to the start of the Attribute file. Since its value is 57FFH then it is a simple matter of adding one to get the desired result. Therefore we use the instruction which increments the DE pair by one: 'INCD E'.

Now we come to the portion of code which deals with scrolling the attributes from the bottom left hand side to the top left hand side:

```
LD B,21
```

The B register is set with the row counter and the loop DOLEFT is then entered. This time the HL register pair points to the new attribute and the DE pair points one row up to the old attribute. Like the first

portion of code which dealt with the scrolling, the top row of the contents of the new attribute address is placed into the old attribute address. The pointers are then updated to point to the next attributes. Since we are going down the attribute file we must add an offset of 32 to both the DE and HL register pairs. The BC register pair is placed onto the stack, which saves the row counter from being corrupted. The BC pair is then LOADED with 32 which is the offset needed to point to the next row down. This is added to the HL pair so that it now points to the next row down. It is then exchanged with the DE pair so that it too can be updated. The instruction EX DE,HL (Exchange the DE pair with the HL pair) is used because there is no such instruction as ADD DE,BC. Therefore we swap the two pairs and update the other pointer with a second ADD HL,BC instruction. To restore the registers to their new values we have to use the EXCHANGE instruction once more. The row counter is then restored by the instruction POP BC (POP the top of the stack to the BC pair) which is decremented and tested to see if we have moved 21 bytes in the DJNZ instruction.

```
DOLEFT: LD  A,(HL) ;GET ATTRIBUTE BELOW
        LD  (DE),A ;AND PLACE ON OLD ONE.
        PUSH BC ;SAVE ROW COUNTER
        LD  BC,32 ;NEXT ROW OFFSET
        ADD HL,BC ;HL NOW POINTS TO NEW ROW
        EX  DE,HL ;SWOP FOR DE
        ADD HL,BC ;WHAT WAS DE NOW POINTS TO NEW
        ;ROW
        EX  DE,HL ;RESTORE BACK TO NORMAL
        POP BC ;RESTORE ROW COUNTER
        DJNZ DOLEFT ;REPEAT UNTIL ALL ROWS DONE
```

After executing the above loop the DE pair points to the bottom lefthand side of the attribute file. We now need to scroll the bottom line from the right to the left. Therefore, we need the HL pair pointing to the attribute to the right of the DE pair.

```
LD  H,D ;PLACE DE INTO HL
LD  L,E
```

The DE pair is first copied into the HL pair by the two instructions LOADING the high part of the DE pair (the D register) into the high part of the HL pair (the H register). This is performed with the instruction LD H,D (LOAD into the H register the contents of the D register). Then the low part is copied by using the instruction LD L,E (LOAD the L register with the contents of the E register).

Now that the HL pair is also pointing to the bottom left hand corner, point it to the right of the DE pair by incrementing it by one using instruction INC HL (INcrement the HL pair by one).

```
INC HL ;HL POINTS ONE TO RIGHT OF DE
```

The next portion of code is very similar to the DOTOP but this time we are scrolling the attributes in the opposite direction. Notice that we are incrementing the pointers instead of decrementing them.

```
LD B,31 ;LOAD B REG WITH COLUMN COUNTER
DOBOT:LD A,(HL) ;GET DATA FROM THE ATTR ON RIGHT
LD (DE),A ;AND PLACE IT IN THE LEFT
INC HL ;POINT TO NEW ATTRIBUTES
INC DE ;TO THE RIGHT
DJNZ DOBOT ;REPEAT UNTIL DONE 31 TIMES
```

After scrolling the bottom portion of the attributes we now deal with the scrolling of the right hand side of the attributes.

First, we adjust the HL register pair to point to one row above the DE pointer. The instruction LD BC,-32 is the one used. The final portion of code is similar to the loop DOLEFT but this time we are only scrolling 20 rows as the second row from the top has its new attribute SAVED on the stack.

```
DEC HL ;RE-ADJUST HL BACK ONE
LD BC,-32 ;AND MAKE IT POINT TO ONE
ADD HL,BC ;ROW ABOVE DE
LD B,20 ;THIS TIME ONLY DO 20 TIMES
;AS LAST ATTRIBUTE IS HELD ON
;THE STACK
DORIG:LD A,(HL) ;GET THE ATTRIBUTE ABOVE
LD (DE),A ;AND PLACE IT TO THE ATTRIBUTE
;BELOW
PUSH BC ;SAVE ROW COUNTER
LD BC,-32 ;LOAD BC WITH OFFSET
ADD HL,BC ;MAKE HL POINT TO ONE ROW ABOVE
EX DE,HL ;SAVE TEMP IN THE DE PAIR
ADD HL,BC ;MAKE OLD DE POINT ONE ROW ABOVE
EX DE,HL ;AND RESTORE BACK DE AND HL
POP BC ;AS WELL AS THE ROW COUNTER
DJNZ DORIG ;REPEAT UNTIL ALL ROWS DONE
```

Finally, we 'POP' off the first attribute that we saved and place it to

the last attribute on the second row of the screen. The RETURN instruction then RETURNS control back to the BASIC program.

```

        POP AF      ;GET ATTRIBUTE VALUE ON STACK
        LD  (DE),A  ;AND PLACE IN NEW POSITION
        RET          ;AND RETURN

        END

        ORG      30000

        ATTRADD EQU 5800H

        LD      HL, ATTRADD+31 ;POINT TO RIGHT HAND
                                ;SIDE OF ATTR
        LD      DE, ATTRADD+30 ;DE POINTS TO THE NEXT ATTR
        LD      A, (HL)        ;SAVE FIRST ATTRIBUTE
        PUSH   AF              ;ON THE STACK
        LD      B, 31          ;LOAD B WITH COLUMN COUNTER
DOTOP:  LD      A, (DE)        ;GET NEXT ATTRIBUTE
        LD      (HL), A       ;AND PLACE IT TO THE
                                ;ATTRIBUTE TO THE LEFT
        DEC    DE              ;POINT TO THE NEXT
                                ;ATTRIBUTES
        DEC    HL
        DJNZ   DOTOP          ;REPEAT UNTIL DONE
                                ;ALL COLUMNS

;HL NOW POINTS TO 5800H IE TOP
;LEFT HAND CORNER OF ATTRIBUTES

        LD      BC, 32        ;LET HL POINT ONE ROW DOWN
        ADD    HL, BC         ;BY ADDING 32
        INC    DE              ;DE NOW POINTS TO START
                                ;OF ATTRIBUTE FILE IE 5800H

        LD      B, 21

```

```

DOLEFT: LD      A, (HL)      ;GET ATTRIBUTE BELOW
        LD      (DE), A    ;AND PLACE ON OLD ONE.
        PUSH   BC          ;SAVE ROW COUNTER
        LD      BC, 32     ;NEXT ROW OFFSET
        ADD    HL, BC      ;HL NOW POINTS TO NEW ROW
        EX     DE, HL      ;SWAP FOR DE
        ADD    HL, BC      ;WHAT WAS DE NOW POINTS
                          ;TO NEW ROW
        EX     DE, HL      ;RESTORE BACK TO NORMAL
        POP    BC          ;RESTORE ROW COUNTER
        DJNZ   DOLEFT     ;REPEAT UNTIL DONE ALL ROWS

```

```

        LD      H, D       ;PLACE DE INTO HL
        LD      L, E
        INC    HL          ;HL POINTS ONE TO
                          ;RIGHT OF DE

```

```

        LD      B, 31     ;LOAD B REG WITH COLUMN
                          ;COUNTER

```

```

DOBOT:  LD      A, (HL)   ;GET DATA FROM THE ATTR
                          ;ON RIGHT
        LD      (DE), A   ;AND PLACE IT IN THE LEFT
        INC    HL        ;POINT TO NEW ATTRIBUTES
        INC    DE        ;TO THE RIGHT
        DJNZ   DOBOT     ;REPEAT UNTIL DONE 31 TIMES

```

```

        DEC    HL        ;RE-ADJUST HL BACK ONE
        LD      BC, -32   ;AND MAKE IT POINT TO ONE
        ADD    HL, BC    ;ROW ABOVE DE

```

```

        LD      B, 20    ;THIS TIME ONLY DO 20 TIMES
                          ;AS LAST ATTRIBUTE
                          ;IS HELD ON THE STACK

```

```

DORIG: LD      A, (HL)      ;GET THE ATTRIBUTE ABOVE
        LD      (DE), A    ;AND PLACE IT TO THE
                                ;ATTRIBUTE BELOW
        PUSH   BC          ;SAVE ROW COUNTER
        LD      BC, -32    ;LOAD BC WITH OFFSET
        ADD    HL, BC      ;MAKE HL POINT TO ONE
                                ;ROW ABOVE
        EX     DE, HL      ;SAVE TEMP IN THE DE PAIR
        ADD    HL, BC      ;MAKE OLD DE POINT ONE
                                ;ROW ABOVE
        EX     DE, HL      ;AND RESTORE BACK DE AND HL
        POP    BC          ;AS WELL AS THE ROW COUNTER
        DJNZ   DORIG      ;REPEAT UNTIL DONE ALL ROWS

        POP    AF          ;GET ATTRIBUTE VALUE
                                ;ON STACK
        LD      (DE), A    ;AND PLACE IN NEW POSTION
        RET              ;AND RETURN

```

```

5 CLEAR 29998: BORDER 0:CLS
10 LET X=22528
20 FOR F=X TO X+31
30 POKE F,8*RND*7
40 NEXT F
50 FOR S=1 TO 21
60 LET A=X+32*S
70 POKE A,8*RND*7
80 LET A=A+31
90 POKE A,8*RND*7
100 NEXT S
110 FOR X=23201 TO 23201+30
130 POKE X,8*RND*7
140 NEXT X
150 RANDOMIZE USR 30000
160 PAUSE 10
170 GO TO 150

```

9 Using the ROM routines

The ROM (Read Only Memory) is a permanent program built into the Spectrum. It handles the interpretation and execution of BASIC and controls the operating system. It is contained in the first 16k of the Spectrum's memory map and enables you to RUN and edit BASIC programs. The ROM manages sound, graphics and communication between the cassette port and the keyboard.

It is worthwhile looking at some of the routines contained in the ROM as these can be used by the programmer when memory space is scarce. They can also be utilised if a complex arithmetic routine is needed or a particular function of the Spectrum is to be used. Each ROM routine is found at a particular memory location and we need to know the correct sequence of operations and the values that must be placed in selected registers.

PRINT and CHANNEL routines

The following routine allows the user to PRINT characters to what are known as *streams* via channels. A channel is the route by which input and output are effected to the various devices on a computer. Examples of such devices on the Spectrum are the keyboard, the printer and the screen. There are seven channels on the Spectrum, the most useful of which are given below:

Channel 0 or Channel K (used for input and output to bottom part of screen)

Channel 1 or Channel K (as Channel 0)

Channel 2 or Channel S (used for printing to the screen)

Channel 3 or Channel P (used for printing to the printer)

Before we PRINT to a channel we must indicate to the Spectrum which stream we wish to use. This is known as 'Opening a Channel'. When we use the ROM routines to PRINT a character the output will go to the currently selected channel until we open another channel. To open a channel we LOAD the A register with the channel number we wish to use. We then call the ROM routine at address 1601 hex which

will 'open' that channel. Any PRINTING done by means of the ROM routines would now send the output to the channel selected.

Therefore, if we wanted to start PRINTING to the screen we would open channel 2:

```
LD      A,2      ;select screen
CALL   1601H    ;open channel
```

Now to start PRINTING a character to the currently selected channel we can use a routine at the address 10 hex. The A register is LOADED with the ASCII value of the character we wish to PRINT. We use the call instruction:

```
RST    10H      ;print the character in A reg
                    ;to the current channel
```

The RST 10 subroutine is extremely useful when PRINTING to the screen because it automatically updates the system variables by updating the print co-ordinate of the next character position. It can also handle all the control characters, thus enabling us to simulate the PRINT AT ,TAB,INK,PAPER,OVER,INVERSE,BRIGHT and FLASH BASIC commands. Given below is a table of the control characters and their code values:

Character	ASCII value	(decimal)
INK	16	
PAPER	17	
FLASH	18	
BRIGHT	19	
INVERSE	20	
OVER	21	
AT	22	
TAB	23	

Therefore if we wished to PRINT the character 'A' on the screen at Y position 10 and X position 10, in blue INK yellow PAPER we would use the code:

```
LD      A,2      ;first open channel 2
CALL   1601H
LD      A,22     ;PRINT AT
RST    10H
LD      A,10     ;PRINT AT10
RST    10
```

```

LD      A,10
RST    10H      ;PRINT AT 10,10;
LD      A,16
RST    10H      ;INK
LD      A,1
RST    10H      ;LOAD A WITH CODE FOR BLUE
LD      A,17
RST    10H      ;BLUE INK
LD      A,6
RST    10H      ;PAPER
LD      A,'A'
RST    10H      ;YELLOW PAPER
LD      A,'A'
RST    10H      ;PRINT CHARACTER

```

PRINT STRING – 203C hex

This routine can be used to print a string of characters. The DE register pair is set up to contain the address of the start of the string that we wish to print, the BC register contains the number of characters in the string. As an example we have taken the last routine but this time all the characters have been put into a string:

```

LD      A,2
CALL   1601H    ;open channel 2
LD      DE,string ;point to string
LD      BC,8    ;number of chars in string.
CALL   203CH    ;print string.
RET
string: DB      22,10,10,16,1,17,6,'A'

```

Printing numbers*PRINT LINE NUMBER – 1A1BH*

This is a simple routine which is used by the ROM to PRINT line numbers in BASIC therefore it is limited to PRINTING numbers from 0 to 9999. The BC register pair is set up with the number we wish to PRINT;

```

LD      BC,400H
CALL   1A1BH    ;print 1024 decimal
RET

```

PRINT LINE NUMBER2 – 1A28H

This routine is identical to the one above except that the number is pointed to by the HL pair.

LD	HL,6000H	address	contents
CALL	1A28H	6000h	03H
RET		6001h	CAH

This section of code would result in the number 970 decimal being PRINTED to the current channel. You may notice that the two bytes that make up the number are stored in memory in the reverse manner to that usually used by the Z80 (i.e. MSB, LSB). This is due to the line numbers being held the 'wrong way' round in memory.

SCREEN ADDRESSING

Screen addressing routines in the ROM are used for PRINTING and PLOTTING. When we PLOT a point using the Spectrum BASIC the screen is split into a grid of 176 lines by 256 points with the co-ordinate 0,0 starting at the bottom left hand side. There are 16 remaining lines at the bottom of the screen which are used by the BASIC operating system for INPUT, error messages, etc.

These routines may be used if required. However if you wish to PLOT or PRINT to any part of the screen, including the 16 'unusable' lines, I suggest you use my PLOT routine described in chapter 10 on the display file.

CHARADD - 0E9BH

This routine can be used to find the address on the screen for a given character line number 1-24. The first line starts at the bottom of the screen. The B register is LOADED with the line number of which we want to find the address. After calling the routine the HL register pair is RETURNED with the address of the first character line.

```
LD      B,1
CALL   0E9BH    ;find address of line 1
```

PIXADD - 22AAH

This routine is used to find the address of a point on the screen. The B register is set with Y co-ordinate and the C register with the X Co-ordinate, the address of which we wish to find. On RETURNING from the subroutine the HL register pair contains the screen address and the A register the bit position of the screen which is 0-7. Note that this is inverted and refers to the bit sequence left to right, not right to left. Therefore if A returns 0 then the leftmost bit, bit seven is referred to and if A gives 2 then bit five (the third from the left) is referred to, etc.

Beeper routine - 03B5

The BEEPER routine is called by LOADING the DE register with the duration, which is the frequency of the note multiplied by the number of seconds we wish the note to last. The pitch is LOADED into the HL register pair and we call the BEEPER routine at address 03B5 hex. The notes and their corresponding pitch values are given in the table below. For example, if we wanted to play middle G sharp for half a second (i.e. G#4) we would LOAD the DE register pair with $0.5 * 415 = CF$ hex ($415 = 19F$ hex), LOAD the HL pair with 3FF, and then call the routine:

```

;play middle G sharp
;for half a second
LD DE,0CFH
LD HL,19FH
CALL 03B5H

```

Table of duration and pitch values for musical notes

Note	Freq (Hz)	Dur (sec) (hex)	Pitch (hex)	Note	Freq (Hz)	Dur (sec) (hex)	Pitch (hex)
C 0	16.35	10	6868	A 1	55.00	37	1EF4
C#0	17.32	11	628D	A#1	58.28	3A	1D34
D 0	18.35	12	5D03	B 1	61.74	3D	1B90
D#0	19.44	13	57CB	C 2	65.40	41	1A03
E 0	20.60	14	52D7	C#2	69.20	45	188C
F 0	21.83	15	4E2B	D 2	73.40	49	172A
F#0	23.12	17	49CC	D#2	77.76	4D	15DC
G 0	24.50	18	45A3	E 2	82.40	52	149F
G#0	25.96	19	41B6	F 2	87.32	57	1374
A 0	27.50	1B	3E06	F#2	92.48	5C	125C
A#0	29.14	1D	3A87	G 2	98.00	62	1152
B 0	30.87	1E	373E	G#2	103.84	67	1057
C 1	32.70	20	3425	A 2	110.00	6E	F6B
C#1	34.64	22	3137	A#2	116.56	74	E8B
D 1	36.70	24	2E72	B 2	123.48	7B	DB8
D#1	38.88	26	2BD6	C 3	130.80	82	CF2
E 1	41.20	29	295C	C#3	138.56	8A	C37
F 1	43.66	2B	2706	D 3	146.80	92	B86
F#1	46.24	2E	24D7	D#3	155.52	9B	ADF
G 1	49.00	31	22C2	E 3	164.80	A4	A40
G#1	51.92	33	20CC	F 3	174.64	AE	9AB

Note	Freq (Hz)	Dur (sec) (hex)	Pitch (hex)	Note	Freq (Hz)	Dur (sec) (hex)	Pitch (hex)
F#3	184.96	B8	91F	A 5	880.00	370	1D3
G 3	196.00	C4	89A	A#5	932.48	3A4	1B7
G#3	207.68	CF	81C	B 5	987.84	3DB	19C
A 3	220.00	DC	7A6	C 6	1046.40	416	183
A#3	233.12	E9	736	C#6	1108.48	454	16C
B 3	246.96	F6	6CD	D 6	1174.40	496	156
C 4	261.60	105	66A	D#6	1244.16	4DC	141
C#4	277.12	115	60C	E 6	1318.40	526	12D
D 4	293.60	125	5B3	F 6	1397.12	575	11B
D#4	311.04	137	560	F#6	1479.68	5C7	109
E 4	329.60	149	511	G 6	1568.00	620	F8
F 4	349.28	15D	4C6	G#6	1661.44	67D	E9
F#4	369.92	171	480	A 6	1760.00	6E0	DA
G 4	392.00	188	43D	A#6	1864.96	748	CC
G#4	415.36	19F	3FF	B 6	1975.68	7B7	BF
A 4	440.00	1B8	3C4	C 7	2092.80	82C	B2
A#4	466.24	1D2	38C	C#7	2216.96	8A8	A7
B 4	493.92	1ED	357	D 7	2348.80	92C	9C
C 5	523.20	20B	326	D#7	2488.32	9B8	91
C#5	554.24	22A	2F7	E 7	2636.80	A4C	87
D 5	587.20	24B	2CA	F 7	2794.24	AEA	7E
D#5	622.08	26E	2A1	F#7	2959.36	B8F	75
E 5	659.20	293	279	G 7	3136.00	C40	6D
F 5	698.56	2BA	254	G#7	3322.88	CFA	65
F#5	739.84	2E3	231	A 7	3520.00	DC0	5E
G 5	784.00	310	20F	A#7	3729.92	E91	57
G#5	830.72	33E	1F0	B 7	3951.36	F6F	50

As a demonstration the program below plays a well known tune which true Spectrum owners should recognise. The program reads a series of pitch and duration values which are passed to the BEEPER. To finish the tune we use the byte 0FF hex. If you are really into music a more sophisticated version of PLAY is given in Chapter 12.

Assembler Listing

```

        ORG      32000D
BEEPER EQU      03B5H

PLAY:
GETV:  LD        IX,FRERE          ;POINT TO START OF MUSIC
        LD        L,(IX+0)        ;LOAD L WITH LOW PART
                                       ;OF THE PITCH
        LD        H,(IX+1)        ;LOAD H WITH HIGH PART
                                       ;OF THE PITCH
        INC       H                ;IF H IS 0FF HEX THEN
                                       ;END OF MUSIC
        RET       Z                ;SO RETURN
        DEC       H                ;RESTORE HIGH PART OF PITCH
        LD        E,(IX+2)        ;LOAD E WITH LOW PART
                                       ;OF DURATION
        LD        D,(IX+3)        ;LOAD D WITH HIGH PART
                                       ;OF DURATION
        PUSH     IX                ;SAVE MUSIC POINTER ON
                                       ;THE STACK
        CALL     BEEPER           ;CALL BEEPER ROUTINE IN
                                       ;THE ROM
        POP      IX                ;RESTORE IX
        LD        DE,4            ;LOAD DE WITH 4
        ADD     IX,DE            ;AND POINT TO NEXT
                                       ;PIECE OF MUSIC
        JR      GETV             ;GET NEXT PIECE OF MUSIC

FRERE:
        DEFW     66AH
        DEFW     105H
        DEFW     5B3H
        DEFW     125H
        DEFW     560H
        DEFW     98H
        DEFW     5B3H
        DEFW     92H
        DEFW     66AH
        DEFW     105H

```

```
DEFW 66AH
DEFW 105H
DEFW 5B3H
DEFW 125H
DEFW 560H
DEFW 98H
DEFW 5B3H
DEFW 92H
DEFW 66AH
DEFW 105H
```

```
DEFW 560H
DEFW 137H
DEFW 4C6H
DEFW 15DH
DEFW 43DH
DEFW 188H
```

```
DEFW 560H
DEFW 137H
DEFW 4C6H
DEFW 15DH
DEFW 43DH
DEFW 188H
```

```
DEFW 43DH
DEFW 126H
DEFW 3FFH
DEFW 67H
DEFW 43DH
DEFW 0C4H
DEFW 4C6H
DEFW 0AEH
DEFW 560H
DEFW 98H
DEFW 5B3H
DEFW 92H
DEFW 66AH
DEFW 105H
```

```

DEFW 43DH
DEFW 126H
DEFW 3FFH
DEFW 67H
DEFW 43DH
DEFW 0C4H
DEFW 4C6H
DEFW 0AEH
DEFW 560H
DEFW 9BH
DEFW 5B3H
DEFW 92H
DEFW 66AH
DEFW 105H

DEFW 66AH
DEFW 105H
DEFW 89AH
DEFW 0C4H
DEFW 66AH
DEFW 20AH

DEFW 66AH
DEFW 105H
DEFW 89AH
DEFW 0C4H
DEFW 66AH
DEFW 20AH
DEFW 0FFFFH
END

```

Hexadecimal Listing

7D00	DD	21	21	7D	DD	6E	00	DD
7D08	66	01	24	C8	25	DD	5E	02
7D10	DD	56	03	DD	E5	CD	B5	03
7D18	DD	E1	11	04	00	DD	19	18
7D20	E3	6A	06	05	01	B3	05	25
7D28	01	60	05	9E	00	B3	05	92
7D30	00	6A	06	05	01	6A	06	05
7D38	01	B3	05	25	01	60	05	9E

7D40	00	E3	05	92	00	6A	06	05
7D48	01	60	05	37	01	C6	04	5D
7D50	01	3D	04	88	01	60	05	37
7D58	01	C6	04	5D	01	3D	04	88
7D60	01	3D	04	26	01	FF	03	67
7D68	00	3D	04	C4	00	C6	04	AE
7D70	00	60	05	9B	00	E3	05	92
7D78	00	6A	06	05	01	3D	04	26
7D80	01	FF	03	67	00	3D	04	C4
7D88	00	C6	04	AE	00	60	05	9B
7D90	00	E3	05	92	00	6A	06	05
7D98	01	6A	06	05	01	9A	08	C4
7DA0	00	6A	06	0A	02	6A	06	05
7DA8	01	9A	08	C4	00	6A	06	0A
7DB0	02	FF	FF	05	01	6A	06	05
7DB8	01	E3	05	25	01	60	05	9B
7DC0	00	E3	05	92	00	6A	06	05
7DC8	01	60	05	37	01	C6	04	5D
7DD0	01	3D	04	88	01	60	05	37
7DD8	01	C6	04	5D	01	3D	04	88
7DE0	01	3D	04	26	01	FF	03	67
7DE8	00	3D	04	C4	00	C6	04	AE
7DF0	00	60	05	9B	00	E3	05	92
7DF8	00	6A	06	05	01	3D	04	26

TAPE LOADING AND SAVING

The following routines can be used in programs to enable the user to SAVE and LOAD data to cassette. When you SAVE or LOAD data in BASIC it comes in two parts. The first part is a 'header' containing information about the file, its name, the length of the file, and its start address. The two main routines are SAVEDATA, at address 04C2 hex, and LOADDATA at 0556 hex. These routines are called with the IX register pair containing the address where the data is placed or SAVED from and the DE register containing the number of bytes we wish to SAVE or LOAD. We LOAD the A register with 0 if we are dealing with the header section or 0FF hex if we are dealing with a data block. When LOADING a file the Carry flag is set when actually LOADING data but it is reset if we want to VERIFY data.

Below is a set of routines which the programmer can use with the cassette.

SAVEBYTES:

```
LD    DE,NBYTES           ;number of bytes to save
LD    IX,START            ;start of block
LD    A,0FFH             ;saving a data block
CALL  SAVEDATA           ;save bytes
RET
```

SAVEHEADER:

```
LD    DE,17              ;save 17 bytes length of header
LD    IX,START-OF-HEADER ;point to start of header info
XOR   A                  ;signify header.A=0
CALL  SAVEDATA           ;save header
RET
```

LOADBYTES:

```
LD    DE,NBYTES
LD    IX,START
LD    A,0FFH
SCF                                     ;signify loading
CALL  LOADDATA
RET
```

LOADHEADER:

```
LD    DE,17
LD    IX,START-of-HEADER
XOR   A
SCF
CALL  LOADDATA
RET
```

VERIFYBYTES:

```
LD    DE,NBYTES
LD    IX,START
LD    A,0FFH
AND   A                            ;reset carry flag
CALL  LOADDATA
RET
```

VERIFYHEADER:

```
LD    DE,17
LD    IX,START-of-HEADER
XOR   A                            ;reset carry flag and
CALL  LOADDATA                       ;set A=0
RET
```


The following program uses these routines to make backup copies of most tapes. Please do not infringe copyright by using this to copy protected tapes!

Assembler Listing

```

                ORG    32000
                JP     START

MEM    EQU    23900
PRXSTRING    EQU    203CH

MESS1:  DEFM    22, 1, 0, 'FILENAME: '
SIZ1    EQU    13
MESS2:  DEFM    22, 3, 0, 'PROGRAM TYPE: '
SIZ2    EQU    17
MESS3:  DEFM    22, 6, 0, 'LENGTH: '
SIZ3    EQU    11
MESS4:  DEFM    22, 9, 0, 'START: '
SIZ4    EQU    10
;MESS5: DEFM    22, 12, 0, 'BASIC LENGTH:'
;SIZ5   EQU    16
WAITP:  DEFM    22, 10, 0, '    MAC WAITING FOR HEADER. '
        DEFM    22, 13, 0, ' COPYRIGHT J.K WILSON 1983. '
SIZWP   EQU    62

TBASIC: DEFM    'BASIC'
TNUM:   DEFM    'NUMER'
TCHAR:  DEFM    'CHARA'
TCODE:  DEFM    'BYTES'

SAVEQ:  DEFM    22, 21, 0, 'Do you want a copy?'
SIZQ    EQU    22
SAVER:  DEFM    22, 21, 0, 'Press ENTER when ready.'
SIZR    EQU    26
BLANKM: DEFM    22, 21, 0, '
SIZB    EQU    26

```

```

ERRXSP EQU 23613 ;SYSTEM VARIABLE
;ERROR STACK POINTER
LOADBTS EQU 0556H ;LOAD BYTES ROM ROUTINE
SAVEBTS EQU 04C2H ;SAVE BYTES ROM ROUTINE

NUMTA: ;BASE TEN TABLE
    DEFW 10000
    DEFW 1000
    DEFW 100
    DEFW 10
    DEFW 1

NUMB: DS 5 ;NUMBER BUFFER

OUTXNUM:
    LD IX, NUMTA ;POINT TO TABLE
    LD DE, NUMB ;DE POINTS TO BUFFER
DIGIT: LD C, (IX+0) ;GET LOW BYTE OF BASE 10
    LD B, (IX+1) ;GET HIGH BYTE OF BASE 10
    LD A, '0'-1 ;A REGISTER =30 HEX
    AND A ;CLEAR CARRY

FIN:
    INC A ;CALCULATE NUMBER OF
    SBC HL, BC ;MULTIPLES OF TENS UNTIL
    JR NC, FIN ;CARRY FLAG IS SET
    ADD HL, BC ;RESTORE NUMBER
    LD (DE), A ;PLACE ASCII NUMBER
    ;IN BUFFER
    DEC C ;TEST TO SEE IF FINISED
    INC DE ;BUMP BUFFER POINTER
    JR Z, OUTP ;FINISHED OUTPUT NUMBER
    ;TO CURRENT CHANNEL
    INC IX ;POINT TO NEXT
    ;MULTIPLE OF 10
    INC IX
    JR DIGIT ;FIND NEXT ASCII DIGIT
OUTP: LD DE, NUMB ;NUMBER BUFFER
    LD BC, 5 ;LENGTH OF STRING
    CALL PRXSTRING ;PRINT IT!
    RET ;AND RETURN

```

```
; HEAD GETS HEAD INFORMATION FROM TAPE
; 17 BYTES OF INFORMATION ARE PASSED TO STARTING ADDRESS IX
; NB. DE MUST BE LOADED WITH 17
```

```
;
; 0-1-2-3-4-5-6-7-8-9-10-11-12-13-14-15-16
```

```
-----
; T*****FILENAME*****|*LEN*|*STR*|*PRG*|
-----
```

```
; T=TYPE 0 : BASIC PROGRAM
;         1 : NUMERICAL ARRAY
;         2 : CHARACTER ARRAY
;         3 : BLOCK OF CODE
```

```
; FILENAME : NO MORE THAN TEN BYTES
```

```
; LEN=LENGTH OF CODE IF TYPE 3
```

```
; STR=START ADDRESS OF CODE OR LINE NUMBER
```

```
; PRG=LENGTH OF PROGRAM AREA
```

```
;
HEADER: DS      17
```

TYP	EQU	HEADER+0	; TYPE OF PROGRAM
FILE	EQU	HEADER+1	; FILENAME
LEN	EQU	HEADER+11	; LENGTH OF CODE
STR	EQU	HEADER+13	; STARTING ADDRESS

```
HEADIN: SCF                    ; SET CARRY FLAG
        LD      A, 0
        LD      IX, HEADER      ; IX POINTS TO HEADER BUFFER
        LD      DE, 17         ; 17 BYTES OF INFORMATION
        CALL    LOADBTS        ; LOAD BYTES
        RET
```

```
HEADOUT:
        LD      A, 0           ; SET A TO ZERO
        LD      IX, HEADER      ; POINT TO HEADER
        LD      DE, 17         ; 17 BYTES TO
        CALL    SAVBTS         ; SAVE
        RET
```

SAVECODE:

```

LD      A, 0FFH
LD      IX, MEM      ;POINT TO RAM
LD      DE, (LEN)    ;GET LENGTH FROM HEADER
CALL    SAVEBTS      ;AND SAVE
RET

```

LOADCODE:

```

SCF                                ;SET CARRY FLAG
                                           ;TO SIGNIFY LOADING
LD      A, 0FFH      ;A REG LOADED WITH
                                           ;TYPE OF DATA
LD      IX, MEM      ;POINT TO START OF CODE
LD      DE, (LEN)    ;PUT LENGTH OF CODE INTO DE
CALL    LOADBTS      ;DO LOADING
RET

```

CLS:

```

LD      HL, 4000H      ;CLEAR SCREEN
LD      DE, 4001H
LD      BC, 8*32*24-1
LD      (HL), 0
LDIR
LD      DE, HOME      ;AND PLACE CURSOR
LD      BC, 3
CALL    PRXSTRING     ;AT HOME
RET

```

```
HOME:  DB      22, 0, 0
```

DISPLAY:

```

CALL    CLS           ;CLEAR SCREEN
LD      DE, MESS1     ;PRINT FILENAME STRING
LD      BC, SIZ1
CALL    PRXSTRING
LD      DE, FILE
LD      BC, 10        ;PRINT
CALL    PRXSTRING     ;FILENAME

```

```

LD      DE, MESS2      ;PRINT TYPE STRING
LD      BC, SIZ2
CALL    PRXSTRING

LD      HL, TBASIC
LD      A, ( TYP )
LD      E, A           ;SAVE TYP IN E REG
SLA     A              ;2*TY
SLA     A              ;4*TYF
ADD     A, E          ;4*TYF+TYP=5*TYF
LD      E, A
LD      D, 0          ;PUT OFFSET IN DE
ADD     HL, DE
;HL POINTS TO STRING
EX      DE, HL
LD      BC, 5         ;NUMBER OF BYTES TO PRINT
CALL    PRXSTRING     ;TYPE

LD      DE, MESS3     ;PRINT LENGTH STRING
LD      BC, SIZ3
CALL    PRXSTRING
LD      HL, ( LEN )
CALL    OUTXNUM       ;NUMBER OF BYTES

LD      DE, MESS4     ;STARTING LINE/ADDRESS
LD      BC, SIZ4
CALL    PRXSTRING
LD      HL, ( STR )
CALL    OUTXNUM
RET

START:
LD      SP, STACK
LD      A, 2          ;OPEN CHANNEL 'S'
CALL    1601H

```

ERRORS:

```

LD      SP, STACK
LD      HL, ERRSP      ;ERROR STACK
LD      (HL), LOW(ERRORS)
INC     HL
LD      (HL), HIGH(ERRORS)
DEC     HL
LD      (23613), HL

```

NEXT:

```

CALL    WAITM          ;WAIT FOR HEADER MESSAGE

CALL    HEADIN         ;GET HEADER
CALL    DISPLAY        ;DISPLAY INFORMATION
CALL    LOADCODE       ;LOAD CODE
CALL    WANT           ;DOES HE WANT TO SAVE THIS?
JR      NZ, NEXT

ACOPY:  CALL    SAVEMESS ;ASK HIM IF HE IS READY
                ;TO SAVE ETC ETC.
CALL    HEADOUT        ;OUTPUT HEADER
CALL    PAUSE          ;WAIT
CALL    PAUSE
CALL    SAVECODE       ;SAVE CODE
CALL    WANT           ;DOES HE WANT TO MAKE
                ;ANOTHER COPY?
JR      Z, ACOPY       ;YES. WELL MAKE ANOTHER ONE
JR      NEXT

```

WAITM:

```

CALL    CLS           ;CLEAR SCREEN
LD      DE, WAITP     ;TELL THEM
LD      BC, SIZWP     ;WE ARE WAITING
CALL    PRXSTRING
RET

```

WANT:

```

LD      DE, SAVEQ     ;PROMPT FOR ANSWER
LD      BC, SIZQ
CALL    PRXSTRING

```

```

WAITK: CALL KEY ;GET KEYBOARD STATUS
        CP 'Y'
        JR Z,RSZ ;IF YES RETURN WITH A
        CP 'y'
        JR Z,RSZ ;ZERO
        CP 'N'
        JR Z,RSZ-1 ;NON ZERO
        CP 'n'
        JR NZ,WAITK
        AND A ;NON ZERO FLAG RESET
RSZ: PUSH AF ;SAVE FLAGS
BLK: LD DE,BLANKM ;BLANK OUT BOTTOM MESSAGE
      LD BC,SIZB
      CALL PRXSTRING
      POP AF ;RESTORE FLAGS
      RET

PAUSE: CALL PAUSE2 ;WAIT A WHILE
PAUSE2: LD HL,0
        LD DE,0
        LD BC,0FFFFH
        LDIR
        RET

SAVEMESS:
        LD BC,SIZR ;PRINT MESSAGE
        LD DE,SAVER ;TO PROMPT FOR ENTER
        CALL PRXSTRING
ENT: CALL KEY
      CP 0DH
      JR NZ,ENT ;WAIT FOR ENTER
      CALL RSZ ;BLANK OUT BOTTOM SCREEN
      RET

FLAGS EQU 23611 ;STATE OF KEYBOARD
LASTXK EQU 23560 ;LAST KEY PRESSED

```

KEY:

DEPR:

```

LD      A,(FLAGS)      ;LOOK AT STATUS OF KEYBOARD
BIT     5,A
JR      Z,KEY          ;WAIT FOR A KEY
                          ;TO BE PRESSED

RES     5,A
LD      (FLAGS),A
LD      A,(LASTXK)     ;GET KEY VALUE
RET
DS      100           ;STACK SPACE
STACK: DB      0
ERRSP:  DEFW    0      ;ERROR STACK SPACE

```

Hexadecimal Listing

```

7D00  C3 C3 7E 16 01 00 46 49
7D08  4C 45 4E 41 4D 45 3A 20
7D10  16 03 00 50 52 4F 47 52
7D18  41 4D 20 54 59 50 45 3A
7D20  20 16 06 00 4C 45 4E 47
7D28  54 48 3A 20 16 09 00 53
7D30  54 41 52 54 3A 20 16 0A
7D38  00 20 20 20 20 4D 41 43

7D40  20 57 41 49 54 49 4E 47
7D48  20 46 4F 52 20 48 45 41
7D50  44 45 52 2E 20 16 0D 00
7D58  20 43 4F 50 59 52 49 47
7D60  48 54 20 4A 2E 4B 20 57
7D68  49 4C 53 4F 4E 20 31 39
7D70  38 33 2E 20 42 41 53 49
7D78  43 4E 55 4D 45 52 43 48

```


7D80	41	52	41	42	59	54	45	53
7D88	16	15	00	44	6F	20	79	6F
7D90	75	20	77	61	6E	74	20	61
7D98	20	63	6F	70	79	3F	16	15
7DA0	00	50	72	65	73	73	20	45
7DA8	4E	54	45	52	20	77	68	65
7DB0	6E	20	72	65	61	64	79	2E
7DB8	16	15	00	20	20	20	20	20
7DC0	20	20	20	20	20	20	20	20
7DC8	20	20	20	20	20	20	20	20
7DD0	20	20	10	27	E8	03	64	00
7DD8	0A	00	01	00	00	00	00	00
7DE0	00	DD	21	D2	7D	11	DC	7D
7DE8	DD	4E	00	DD	46	01	3E	2F
7DF0	A7	3C	ED	42	30	FB	09	12
7DF8	0D	13	28	06	DD	23	DD	23
7E00	18	E6	11	DC	7D	01	05	00
7E08	CD	3C	20	C9	00	00	00	00
7E10	00	00	00	00	00	00	00	00
7E18	00	00	00	00	00	37	3E	00
7E20	DD	21	0C	7E	11	11	00	CD
7E28	56	05	C9	3E	00	DD	21	0C
7E30	7E	11	11	00	CD	C2	04	C9
7E38	3E	FF	DD	21	5C	5D	ED	5B
7E40	17	7E	CD	C2	04	C9	37	3E
7E48	FF	DD	21	5C	5D	ED	5B	17
7E50	7E	CD	56	05	C9	21	00	40
7E58	11	01	40	01	FF	17	36	00
7E60	ED	B0	11	6C	7E	01	03	00
7E68	CD	3C	20	C9	16	00	00	CD
7E70	55	7E	11	03	7D	01	0D	00
7E78	CD	3C	20	11	0D	7E	01	0A
7E80	00	CD	3C	20	11	10	7D	01
7E88	11	00	CD	3C	20	21	74	7D
7E90	3A	0C	7E	5F	CB	27	CB	27
7E98	83	5F	16	00	19	EB	01	05
7EA0	00	CD	3C	20	11	21	7D	01
7EA8	0B	00	CD	3C	20	2A	17	7E
7EB0	CD	E1	7D	11	2C	7D	01	0A
7EB8	00	CD	3C	20	2A	19	7E	CD

FLOATING POINT CALCULATION ROUTINES

The floating point calculator is used by the BASIC interpreter to handle calculations on floating point numbers and strings. Numbers in BASIC are represented by 5 bytes. The way these bytes represent a floating point number is shown in the Spectrum user's manual (chapter 24).

When the Spectrum is interpreting BASIC calculations it uses something known as a calculator stack which stores the interpreted calculation in a series of numbers, strings and operators. The way this line is interpreted is known as REVERSE POLISH NOTATION OR RPN.

Numeric expressions in BASIC consist of operands, variables and operators or functions. In an algebraic expression operators are placed between two operands. For example, in the expression $X+Y$ the operator is $+$, and the operands are X and Y . The principle of RPN is to place the operator after the operands. Thus $X+Y$ is written as $X Y +$.

the larger algebraic expression:

$$A*B+C*D/2$$

could be written as:

$$A B * C D 2 / * +$$

TO READ a RPN expression we use the following technique. READING the expression from left to right any values or operands would be pushed onto a stack. When we arrive at an operand or function the appropriate number of data is POPPED off the stack and that operation or function is executed with the data. The 'new' value is then pushed onto the stack. This is repeated until we have reached the end of the expression:

$$2 4 * 3 8 2 / * +$$

For example, for the above RPN expression we go through the operations shown below. On the lefthand side of the diagram I have given the operations we use as we read the expression from left to right. On the righthand side the current status of the stack is shown. The top of the stack is the rightmost digit.

Operation	Stack
stack 2	2
stack 4	2 4
operator *	8
stack 3	8 3
stack 8	8 3 8
stack 2	8 3 8 2
operator /	8 3 4
operator *	8 12
operator +	20

The answer left on the stack is 20.

The Spectrum has a calculator stack where numbers can be manipulated in the same way as we would deal with an RPN expression. Groups of numbers can be pushed onto a stack and routines can be called to do various operations such as add, subtract and multiply. We can push numbers onto the stack by calling routines within the ROM which allow us to SAVE the numbers in 1 byte, 2 byte or 5 byte form.

STACKA-2D28H

LD A,20

CALL 2D28H ;STACK NUMBER 20

This routine will convert the one byte number contained in the Accumulator to its five byte floating format which is then pushed onto the calculator stack.

There are two other routines which are similar to the last routine. They allow us to stack a two byte integer in the BC register pair and a five byte floating point number contained in the A,E,D,C,B registers. The routine at the address 2D2B hex (STACKBC) will convert a two byte integer into five byte floating point format and push this on the calculator stack. Likewise the routine at the address 2AB6 hex (STACK5) will push the floating point number in the registers A,E,D,C and B.

To retrieve numbers from the stack we have routines which can pop them off and convert them into one byte, two byte or the normal five byte form. The addresses are given below:

UNSTACKA-2D5DH will convert the floating 5 byte number on top of the calculator stack to its equivalent one byte integer and place it in the accumulator.

UNSTACKBC-2D2AH will remove a five byte floating point number from the stack, convert it and place it in the BC register pair.

UNSTACK5 – 2BF1H. This routine is used to place a floating point number from the stack into the registers A,E,D,C and B.

There are two routines which are very useful to the programmer when handling floating point numbers. The routine at the address 2DE3 hex (PRINTFP) will take the top number on the calculator stack and PRINT it to the current channel selected. The second routine at 2C9B hex (ASCTOFP) enables us to convert a number from a string to a floating point number which is pushed to the calculator stack. Look at the following program:

```
LD      HL,STRING ;put to start of string.
LD      (5C5DH),HL ;save in system variable CH-ADD
LD      A,(HL) ;get first character in A reg.
CALL    2C9BH ;convert ascii number to fp.
CALL    2DE3H ;print fp number
RET
DEFM    2.31693
DB      0DH ;carriage return.
```

The HL register pair is set up to point to the start of the string. This is stored in the system variable CH-ADD which usually holds the address of the next character to be interpreted when RUNNING Spectrum BASIC. The A register is LOADED with the first character and the routine ASCTOFP is called. This leaves the binary floating point number at the top of the calculator stack so when we call the routine PRINTFP (2DE3H) the number 2.31693 will be printed out. The end of the string is signified by a carriage return.

To start the floating point calculator we call the routine at address 28 hex with the one byte instruction RST 28H. The data following the call instruction indicates which operations the calculator must perform. The calculator goes through each operation automatically pushing and popping data until it reaches the data 38 hex which signifies the end of the calculation. Some of the most useful codes are given below:

data code (hex)	Name	Action
01	exchange	Swops the two topmost floating numbers.
02	delete	Deletes top number on stack.
03	subtract	Subtracts second number on stack from first
04	multiply	Multiplies the two topmost numbers.
05	divide	Divides the first over the second number.

data code (hex)	Name	Action
06	power	Raises the first to the power of the second
1B	negate	Changes the sign of the top number.
1F	sin	Calculates the sine of the top number.
20	cos	Calculates the cosine of the top number.
21	tan	Calculates the tangent of the top number.
22	arcsin	Calculates the Arcsine of the top number.
23	arccos	Calculates the Arccosine of the top number.
24	arctan	Calculates the Arctan of the top number.
25	log2	Calculates the Log of the top number.
26	exp	Calculates the exponential of the top number.
27	integer	Calculates the integer of the top number.
28	square root	Calculates the square root of the top number.
29	sign	Places the sign of the top number on stack
2A	absolute	Converts the top number to its absolute.
2B	peek	Places the contents of the address at top of stack
2C	in	Scans address at top of stack.
31	duplicate	Duplicate the top of the stack.
38	Endcalc	End of calculation.

The calculator has five constants available in the ROM used for calculating sines and cosines:

Data code (hex)	Name	Action
A0	stack 0	Place the number 0 on the stack.
A1	stack 1	Place the number 1 on the stack.
A2	stack 1/2	Place the number 1/2 on the stack.
A3	stack PI/2	Place half of PI on the stack.
A4	stack 10	Place the number 10 on the stack.

There are six memory locations used by the floating point calculator in order to SAVE numbers on top of the stack. The data codes C0 hex to C5 hex are used to SAVE the topmost number on the stack to one of the six memory locations. The other codes E0 hex to E5 hex are used to place numbers from one of the memory locations to the top of the calculator stack.

Data Name code (hex)	Action
C0 store 0	Place the number on the stack in memory 0.
C1 store 1	Place the number on the stack in memory 1.
C2 store 2	Place the number on the stack in memory 2.
C3 store 3	Place the number on the stack in memory 3.
C4 store 4	Place the number on the stack in memory 4.
C5 store 5	Place the number on the stack in memory 5.
E0 Stack mem 0	Place the contents of memory 0 on the stack.
E1 Stack mem 1	Place the contents of memory 1 on the stack.
E2 Stack mem 2	Place the contents of memory 2 on the stack.
E3 Stack mem 3	Place the contents of memory 3 on the stack.
E4 Stack mem 4	Place the contents of memory 4 on the stack.
E5 Stack mem 5	Place the contents of memory 5 on the stack.

The following program demonstrates how we use the floating calculator to multiply the numbers 2342 (926 hex) and 156 (9C hex).

```

LD    BC,926H
CALL  STACKBC    ;STACK NUMBER
LD    BC,9CH
CALL  STACKBC    ;STACK NUMBER
RST   28H        ;START CALCULATION.
DB    04         ;MULTIPLY
DB    38H        ;END OF CALC.
CALL  PRINTFP    ;PRINT ANSWER
RET

```

The second example shows us how we can convert fairly complex algorithms into machine code. It PLOTS a sine wave on the screen. The BASIC routine takes 17 seconds while the machine code equivalent takes 14 seconds. The reason why the machine code routine is not significantly faster is because the sine calculations inside the ROM are fairly slow.

```

10 FOR X=0 TO 255
20 LET Y=100+50*SIN(X*PI/20)
30 PLOT X,Y
40 NEXT X

```

The equivalent machine code program would be:

```

        ORG      0

PLOTXY EQU     22EEH
STACKA EQU     2D28H
UNSTACKA EQU    2DD5H

SINE:   XOR      A                ;SET X CO-ORD TO 0
        PUSH     AF                ;SAVE X CO-ORD ON STACK
        LD       A,100             ;STACK 100 DECIMAL
        CALL     STACKA
        LD       A,50
        CALL     STACKA            ;STACK 50 DECIMAL
        POP      AF                ;GET X CO-ORD
        PUSH     AF                ;AND SAVE
        CALL     STACKA            ;PLACE X ON FP STACK
        RST      28H               ;START CALCULATION
        DB       0A3H              ;STACK PI/2
        DB       0A4H              ;STACK 10 DECIMAL
        DB       05H               ;DIVIDE (PI/2 BY 10)
        DB       04H               ;MULTPLY (PI/20 BY X)
        DB       1FH               ;SINE (SINE(PI*X/20))
        DB       04H               ;MULTIPLY
                                   ;(50*SIN(PI*X/20))
        DB       0FH               ;ADD
                                   ;100+50*SIN(PI*X/20)
        DB       38H               ;END OF CALCULATION

        CALL     UNSTACKA          ;GET Y CO-ORD IN A REG
        LD       B,A               ;AND SAVE IN B

        POP      AF                ;GET X CO-ORD OFF STACK
        PUSH     AF                ;SAVE AGAIN

        LD       C,A               ;PLACE X CO-ORD IN C REG
        CALL     PLOTXY            ;PLOT A POINT AT X, Y
        POP      AF                ;GET X REGISTER
        INC      A                 ;INCREASE X CO-ORD
        RET      Z                 ;DONE 256 TIMES SO RETURN
        JR       SINE              ;KEEP PLOTTING

        END

```


10 Screen and attribute handling

The screen display is used in order to communicate information to the user. To PRINT information to the screen is a fairly simple and trivial matter in BASIC. However, when we delve into the realms of machine code accessing the screen becomes more complicated due to the complex screen lay out. The display file is split into two sections; the actual screen display file and the attribute file. The attribute file is easy to access because the data held in this file is organised sequentially. The attribute file consists of a 32 by 24 byte array. Each byte contains the information relating to a particular character position on the screen and tells the Spectrum what PAPER and INK colours to use for the display and whether the character is FLASH and/OR BRIGHT.

The Attribute File

The number to be stored in the attribute file can be calculated by using the following method:

- First set your total to zero.
- Add 128 to your total if you want it flashing.
- Add 64 to your total if you want it bright.
- Add 8 times the PAPER colour you want.
- Add the INK colour you want.

The value of the colours are:

- 0 Black
- 1 Blue
- 2 Red
- 3 Magenta
- 4 Green
- 5 Cyan
- 6 Yellow
- 7 White

The bit pattern of an attribute byte is set up like this:

```

Bit  7 6 5 4 3 2 1 0
     f b p p p i i i

```

Where f is the FLASH bit, b is the BRIGHT bit, p is PAPER number, i is INK number. So if, for example, we wanted the colour code for red INK on white PAPER with the BRIGHTness set on we would put the value $64+8*7+2=122$ in the appropriate location.

The address of the start of the attribute file is 22528 or 5800 hex. It can be represented by a grid of 32 column by 24 rows. The start of the file being in the top left-hand corner. If we wished to find the address of a given pair of co-ordinates (row 0, column 0 is in the top left-hand corner), then we could use the following piece of code:

```

;*****
; FIND ATTRIBUTE ADDRESS
; B CONTAINS THE ROW NUMBER
; C CONTAINS THE COLUMN NUMBER
; ON EXIT HL CONTAINS ADDRESS
CL-ATTR: LD    L, B
         LD    H, 0
         ADD  HL, HL
         ADD  HL, HL
         ADD  HL, HL
         ADD  HL, HL
         ADD  HL, HL ;FIND ROW TIMES 32
         LD   B, 0
         ADD  HL, BC ;ADD COLUMN OFFSET
         LD   BC, 5800H ;START OF ATTRIBUTES
         ADD  HL, BC ;ADD START OF ATTRIBUTES
         RET  ;ADDRESS NOW IN HL

```

As you can see we multiply the row number by 32 and by a series of ADD HL, HL instructions. Finally, we add the column offset. This isn't the quickest way of calculating the address but it is the easiest. If you read the chapter on shifting and rotating you may wish to calculate another way of finding the address using bit manipulation. This routine can be used for PEEKing or POKEing at the attribute file. If we wanted to look at the contents of a given row and column we could use the code:

```

PEEK:  CALL  CL-ATTR ;CALCULATE ADDRESS AT
        ;ROW B,COLUMN C
        LD   A,(HL) ;PUT CONTENTS IN A REGISTER.

```

If we wanted to POKE red INK, green PAPER at column 22, row 5 we could write the code:

```
POKE: LD    A,22H    ;CODE FOR RED INK, GREEN PAPER
      LD    BC,0516H ;SET BC TO ROW 5, COLUMN 22
      CALL CL-ATTR  ;FIND ADDRESS
      LD    (HL),A   ;POKE VALUE IN.
```

The Screen File

The screen file is a little more complicated than the attribute file! It consists of three sections of 8 character rows. Each character row is made up of 32 characters split into 8 pixel lines. Since we are working with eight bit bytes the resolution of the screen is 256 by 192. The resolution of a screen determines the number of little dots that make up the data on the screen. The higher the resolution, the more detailed the pictures on the screen. To determine an address on the screen for a given pair of co-ordinates is not an easy matter. If we look at the bit pattern for a screen address:

010ssllrrrcccc

ss is the section number, 0 being the top third, 1 being the middle and 2 the bottom third. 3 indicates an address in the attribute file. ll is the pixel line number (0-7) within a character. rrr is the line number within a section (0-7) and cccc is the column number (0-31).

Using this pattern we can determine an address anywhere on the screen, even down to one single bit. The section number is contained in the top two bits of the row. The pixel line number is also contained in the row, this time in the middle three bits. The pixel line number is the last three bits of the row. The column is represented by a number 0-255. The range 0 to 255 is used because the routine is designed to give the address and bit position of any pixel on the screen. The data for the routine is as follows:

row	ssrrlll
col	ccccbbb

where all the letters have their previous meanings and bbb is the bit position of the pixel.

Now let's look at the routine to calculate the screen address for any given row and column. The B register is LOADED with a row

number in the region 0-191, where row 0 is at the top of the screen. The c register is LOADED with the column number in the range of 0-255. After executing the routine the HL pair will contain the address on the screen, and the A register will contain the bit position (0-7) within that address.

```

; FIND PIXEL ADDRESS
; B CONTAINS Y COORD
; C CONTAINS X COORD
; ON EXIT HL CONTAINS ADDRESS
; A CONTAINS PIXEL NUMBER
PIXADD:
LD   A, B           ;GET Y REGISTER
RRA
SCF
RRA
RRA                ;X1x.. MOVE DOWN SECTION
AND  58H           ;MASK OFF SECTION. 010SS
LD   H, A          ;SAVE IN H REG
LD   A, B          ;GET Y AGAIN.
AND  7             ;WORK OUT PIXEL LINE WITHIN
                        ;CHARACTER
ADD  A, H          ;ADD PREVIOUS RESULT
LD   H, A          ;SAVE IN H REG
LD   A, C
RRCA
RRCA
RRCA
AND  1FH
LD   L, A          ;MOVE DOWN COLS 0-31
LD   A, B          ;GET ROW NUMBER
AND  38H           ;MASKING OFF ROW NUMBER
ADD  A, A
ADD  A, A
OR   L
LD   L, A
LD   A, C          ;GET BIT NUMBER
AND  7             ;A CONTAINS BIT NUMBER
RET

```

This routine could be used to PLOT points on the screen since the PLOT command in BASIC is limited to accessing only 256 by 176 points. To do this we call the PIXADD routine then rotate the pixel to the bit

position we want. So our program to PLOT a point at the co-ordinates in the BC register pair would be:

```

; PLOT A PIXEL AT THE COORDINATES
; X, Y
; B=Y COORD AND C=X COORD
PLOT:
CALL PIXADD      ;Find address for co-ords BC
LD  B, A         ;Put bit position in the B
                        ;register
INC  B           ;now in the range 1-8
XOR  A           ;set A to zero and clear
                        ;carry flag.
PIX:
RRA
DJNZ PIX         ;move pixel dot to position
                        ;required
XOR  (HL)        ;and place it at the address
LD  (HL), A
RET

```

As seen I have used an exclusive OR to place the dot on the screen. This has the effect of turning a pixel on if there wasn't already a dot at the calculated address or turning the pixel off, if it was already set.

When we PRINT a character onto the screen within a character boundary the offset for each byte of data which makes up a character is 256 so it is a simple matter to have a loop such as:

```

;DE POINTS TO DATA WHICH WE WANT TO PRINT
;HL POINTS TO SCREEN ADDRESS. TOP OF CHARACTER
;BOUNDARY
LD  B, 8         ;8 BYTES OF DATA
NXTPL: LD  A, (DE) ;GET DATA
LD  (HL), A     ;PLACE DATA ON SCREEN
INC  DE         ;POINT TO NEXT DATA
INC  H          ;POINT TO NEXT PIXEL LINE
DJNZ NXTPL     ;DO NEXT PIXEL LINE

```

The INCH instruction is the same as adding 256 to the HL register pair and has the effect of getting the next pixel line address below. The offset is always 256 only if we are within a character boundary. If this is not true we have to use the following routine below which I have called INCY:

INCY:

```

INC  H
LD   A, H
AND  7
RET  NZ           ; WITHIN CHAR BOUNDARY
LD   A, H
SUB  8
LD   H, A
LD   A, L
ADD  A, 32       ; NEXT CHAR LINE DOWN
                        ; (WITHIN SECTION)
                        ; ADD 32 DECIMAL

LD   L, A
RET  NC         ; DEF WITHIN SECTION
                        ; NEXT SECTION DOWN

LD   A, H
ADD  A, 8
LD   H, A
XOR  58H       ; 01011000 BINARY
RET  NZ       ; IS THERE A WRAPAROUND
                        ; NEEDED?

LD   H, 40H
RET

```

The routine could be written to run a little quicker. It seems a waste of time to first subtract eight from the H register if we have gone over a character boundary and then to add eight back. The reason we have done this is to incorporate a wrap around effect. This means anything which is printed over the bottom of the screen will appear on the top.

There is also a relationship between the address of the attribute file and the screen file. Study the bit patterns for the high byte of the start of each section on the screen file and the corresponding bit patterns on the attribute file. The table below shows how the high bytes of the addresses relate:

screen addr	screen bit pattern	attr addr	attr bit pattern
40H	01000000	58H	01011000
48H	01001000	59H	01011001
50H	01010000	5AH	01011010

To get the corresponding attribute address from a given screen address shift down the high byte to the left three times and then set bits three and four of the high byte.

```

; THIS PIECE OF CODE GIVES
; THE ADDRESS OF THE ATTRIBUTE FILE IN THE
; HL REGISTER PAIR
; FOR A GIVEN ADDRESS ON THE DISPLAY FILE
; IN THE HL REGISTER PAIR.
LD  A,H      ;01055000
SRA A        ;00105500
SRA A        ;00010550
SRA A        ;00001055
OR   50H     ;01011055
LD  H,A
RET

```

ANIMATION

Using the `INCY` routine we can write another routine which allows us to `PRINT` a character at any pixel position on the screen. Usually when we `PRINT` a character in `BASIC` the character is placed on the standard 32 by 24 grid. Therefore there are only 768 positions at which we can place that character. If we were to write a game using the `BASIC PRINT` statement movement of characters is limited to moving horizontally eight bits at a time and vertically eight pixel lines at a time. The following machine code routine will demonstrate how to move objects around the screen smoothly using pixel movement.

The routine in the ROM which deals with `PRINTING` characters in `BASIC` calculates the screen address for a given pair of co-ordinates. It is then a simple matter of lacing the eight bytes of data which make up a character onto the screen. The screen is constructed in such a way that each vertical line, where the character is to be placed, is 256 bytes below the last pixel line. However, this offset changes when we are `PRINTING` over a character or section boundary. If we wanted to draw a character on any of the 192 pixel lines we would need to keep using the `INCY` routine to find the addresses of successive pixel lines.

Therefore if our character stayed within a character boundary vertically then the following routine would print a character to the screen. The screen address is pointed to by the `HL` register pair and the character data is pointed to by the `DE` register pair.

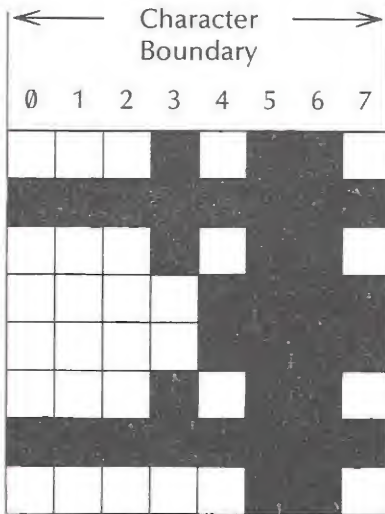
```

BOUNDH: LD      B, 8           ;GET DATA COUNT
NXC:    LD      A, (DE)       ;GET CHARACTER DATA
        LD      (HL), A      ;AND PLACE ON THE SCREEN
        CALL   INCY          ;NEXT PIXEL LINE DOWN
        INC    DE            ;POINT DE TO NEXT
                               ;CHARACTER DATA
        DJNZ   NXC          ;DO 8 TIMES

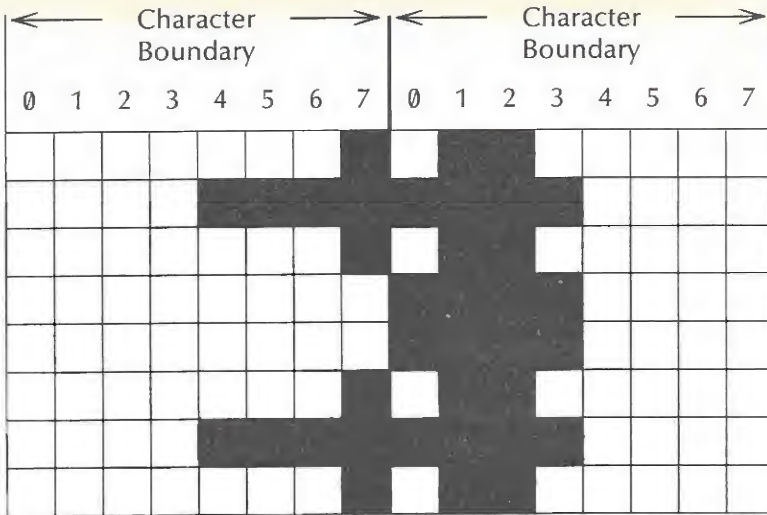
```

As you can see this portion of code is similar to the first routine we used to PRINT a character to the screen. The exception is that the INCH instruction is replaced by calling the routine INCY which calculates the address of the next pixel line down.

The next problem we have to overcome when PRINTING a character on the screen is to deal with its horizontal position. When we want to print an eight bit character, at any of the 256 bits, we may sometimes overlap between two character boundaries. This means that if we can calculate the bit position where the object is to be placed within one of the 32 horizontal positions on the screen we can scroll the eight bit number which makes up one line of the character through two bytes which we then PRINT onto the screen. Look at the following two diagrams. Diagram A shows a space ship being PRINTED within a character boundary. The data only occupies one byte for each horizontal line. When we wish to PRINT an eight by eight pixel object at any horizontal pixel position then we could get an overlap onto the adjacent character position as shown in diagram B. An overlap will occur seven in every eight horizontal bit positions. To find a character's bit position simply get it's x co-ordinate and mask off the bottom three bits by ANDING it with seven. Remember, that this 'bit position' is different from the one we use to describe bit instructions such as SET, RESET and BIT. This time the bit position starts from the left hand side of the byte.



8 by 8 pixel character
within boundary



8 by 8 pixel character over boundary

If given the bit position in which a character lies, then to obtain its two byte equivalent, get the object data to be printed and scroll it from left to right within two bytes. This 'scrolling' from left to right of a 16 bit number is identical to dividing the number by 2. This was explained in the chapter on rotating and shifting. The following piece of code divides a two byte number in the A register (the high part) and the C register (the low part) by 2.

```
LD    C,0 ;clear low byte first
SRL  A    ;scroll A reg from left to right into carry
RR   C    ;scroll c register from left to right through
      ;carry
```

Notice how we clear the low byte of the number by **LOADING** the C register with 0. We would of course do this scrolling until we reach the bit position which we require. Therefore, if the B register contained the bit position we would find the two byte number by using the code:

```
; B REGISTER CONTAINS THE BIT POSITION
LD    A, B ;PLACE B REGISTER INTO
          ;THE A REGISTER
AND   A    ;TEST FOR BIT POSITION=0
JR    Z, BOUND ;WITHIN ONE CHARACTER
          ;BOUNDARY SO DEAL WITH
          ;THIS CASE AT LABEL BOUND
```

```

        LD      C, 0           ;CLEAR RIGHT BYTE FIRST
GETB:   SRL    A              ;SCROLL A REGISTER
        RR     C              ;RIGHT THROUGH THE
                                ;C REGISTER
        DJNZ  GETB           ;UNTIL WE GET INTO
                                ;THE REQUIRED BIT
                                ;POSTION

```

Notice that before we scroll the character we test that the bit position is zero. If the bit position was zero then this means that our character is within a boundary so we deal with this at the label BOUND. If we did not do this and carried on through to scroll the data then we would find that we would end up scrolling the data 256 times.

After we have our two new characters which make up the object then it is simply a case of placing them on the screen. If, for example, the HL register pair was pointing to the screen address where we wanted to place the character then we would place the data at HL and HL+1

```

;DRAW OBJECT ONTO THE SCREEN
;AT THE ADDRESS IN THE HL PAIR

```

```

        LD      (HL), A       ;PLACE LEFT HAND SIDE
                                ;OF THE OBJECT
        INC     HL            ;POINT TO NEXT CHARACTER
                                ;BOUNDARY
        LD      (HL), C       ;PLACE RIGHT HAND SIDE
                                ;OF THE OBJECT

```

To animate, simply DRAW the object onto the screen and to move it, remove the object from its previous position. Then update its new position and DRAW it to the screen. The following machine code routine DRAWS and animates nine space ships on the screen. Each one follows a movement pattern. The object can move in any one of four directions. Direction one indicates that the ship is moving right, two left, four down and eight up. The movement pattern DIRTAB is a table of directions which the ship follows and ends with 255 or FF hex. The ships start at different locations in the table so that the movements are not synchronised.

Each ship has three bytes of data starting from SHIPTB, to represent its x and y co-ordinates and an offset position or vector count

pointing to a direction within the movement table. If a ship had a starting offset of 12 the first direction it would use is at the address `DIRTAB+12`. This contains a one and means that the ship would move right. When the ship comes to the end of the direction table (signified by the byte `FF` hex) it resets its vector count to zero thus pointing to the first byte of the direction table.

The `IX` register is used in this routine to point to each of the ships data. At the start of the program the ships are first drawn onto the screen. The main routine which deals with drawing characters uses the ROM routine `PIXADD` (`22AA` hex). When given the `X` and `Y` co-ordinates on the screen it will RETURN the screen address in the `HL` pair and the `A` register holds the bit position within a byte. The co-ordinates `X=0, Y=0` start the `A` register at the left hand side of the screen 22 lines from the top. Therefore, we only have 176 pixel lines vertically to which to DRAW the objects. Of course if you wanted you could use my `PIXADD` routine which makes full use of the 256 by 192 screen.

Notice in the routine `PRTCHR` how instead of LOADING in the data character bytes directly I first exclusive OR the data with the contents of the screen. This is extremely useful for I can if I wish MOVE over 'background objects' on the screen without corrupting the data. It is like using the `OVER 1` command in BASIC. As well as leaving the background it also serves a useful purpose for effacing the ship from the screen when MOVING it to a new position. Since we are using the `XOR` instruction this will turn off any bits that are all ready on and turn on bits already off.

Assembler Listing

```

;           MOVEMENT ROUTINE
;
;
;
;
;
;
;           ;EXAMPLE OF PIXEL MOVEMENT.

ORG      28000D
JP       START           ;START THE PROGRAM

```

```
SHIP:  DB      22, 255, 22, 15, 15, 22, 255, 22
        ;DATA FOR SPACE SHIP
```

```
PRINOBJ:                                ;SUBROUTINE PRINT OBJECT
                                           ;PRINTS AN OBJECT AT X, Y
                                           ;B REG=Y VALUE
                                           ;C REG=X VALUE
```

```
      PUSH  BC
      CALL  PRTCHR
      POP   BC
      RET
```

```
PRTCHR:                                ;PRINTS A SHIP ON ANY
                                           ;PIXEL POSTION
```

```
      LD    IX, SHIP
      CALL  22AAH      ;FIND PIXEL ADDRESS
```

```
;ROUTINE AT 22AAH FINDS THE ADDRESS ON SCREEN FOR A
;GIVEN X, Y CO-ORDS IN C REG AND B REG
;THE A REG IS RETURN WITH THE START OF PIXEL POSTION
;WITHIN THAT BYTE THIS ROUTINE IS ONLY LIMITED FOR Y
;BEING BETWEEN 0 AND 175 INSTEAD OF THE EXPECTED 0-191
;ALSO THE CO-ORD 0, 0 START IN THE BOTTOM LEFT HAND SIDE.
```

```
      LD    E, A      ;SAVE BIT POSTION IN E
                                           ;REGISTER
      LD    D, 8      ;PIXEL LINE COUNT
      AND   A         ;TEST FOR ZERO PIXEL
                                           ;POSTION
      JR    Z, WBOUND ;WITHIN BOUNDARY SO JUMP
```

```

LINE: LD B, E ;GET BIT POSITION IN
; B REGISTER

; SCROLL DATA B TIMES INTO REGS A AND C
LD A, (IX+0) ;GET DATA
LD C, 0 ;CLEAR RIGHT HAND SIDE

SCROLL:
SRL A ;SCROLL DATA DOWN TO BIT
; POSITION

RR C
DJNZ SCROLL

XOR (HL) ;MERGE LEFT HAND PART OF
; DATA IN!

LD (HL), A

INC HL ;TO THE RIGHT MARCH!
LD A, C ;GET SECOND CHAR
ONE: XOR (HL)
LD (HL), A ;MERGE SECOND CHAR IN.
DEC HL
INC IX ;BUMP NEXT DATA
CALL INCY ;FIND ADDRESS OF NEXT
; PIXEL LINE
DEC D ;HAVE WE DONE THE 8 BYTES?
JR NZ, LINE ;NO! SO DO NEXT LINE
RET ;YES THEN RETURN.

WBOUND: LD B, D ;LOAD B WITH 8
; (PIXEL LINE COUNT)

XBOUND:
LD A, (IX+0) ;GET DATA
XOR (HL) ;MERGE IN WITH DATA
LD (HL), A ;ALREADY ON THE SCREEN
CALL INCY ;NEXT PIXEL LINE DOWN
INC IX ;NEXT DATA BYTE
DJNZ XBOUND ;REPEAT 8 TIMES
RET

```

```

INCY:                                ;FINDS ADDRESS OF NEXT
                                       ;PIXEL LINE ON THE SCREEN

    INC     H
    LD      A, H
    AND     7
    RET     NZ                          ;WITHIN CHAR BOUNDARY
    LD      A, H
    SUB     8
    LD      H, A

    LD      A, L
    ADD     A, 32D                       ;NEXT CHAR LINE DOWN
                                       ;(<WITHIN SECTION)

    LD      L, A
    RET     NC                          ;CHAR WITHIN SECTION

                                       ;NEXT SECTION DOWN

    LD      A, H
    ADD     A, 8
    LD      H, A
    XOR     88                          ;01011000
    RET     NZ
    LD      H, 40H                       ;WRAP AROUND EFFECT
    RET

```

```

DIRTAB:
DB      8, 8, 8, 8, 8, 8, 8, 8, 8, 8, 8
DB      1, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1
DB      4, 4, 4, 4, 4, 4, 4, 4, 4, 4, 4
DB      2, 2, 2, 2, 2, 2, 2, 2, 2, 2, 2
DB      0FFH                          ;END OF DIRECTION TABLE

```

```

GETDIR:
LD     HL,DIRTAB      ;POINT TO DIRECTION TABLE
LD     E,(IY+VECTCN) ;GET SHIPS POINTER
LD     D,0
ADD    HL,DE          ;POINT TO DIRECTION
LD     A,(HL)         ;GET DIRECTION
CP     0FFH           ;IS THIS THE END OF
                        ;THE TABLE
JR     NZ,MOVEIT      ;NO THEN MOVE SHIP!
XOR    A              ;SET A TO ZERO
LD     (IY+VECTCN),A ;SET VECTOR COUNT TO ZERO
RET

```

```

MOVEIT:
INC    (IY+VECTCN)   ;INCREASE VECTOR COUNT
                        ;FOR NEXT GO
CP     8              ;GOING UP
JP     Z,UPD
CP     4
JP     Z,DOWND       ;GOING DOWN
CP     1
JP     Z,RIGHTD      ;GOING RIGHT
CP     2
JP     Z,LEFTD       ;GOING LEFT
RET

```

```

UPDAA:
CALL   PRINOBJ       ;REPRINT SHIP
                        ;AT NEW POSTION
LD     (IY+XPOS),C   ;SAVE NEW XPOSTION
LD     (IY+YPOS),B   ;SAVE NEW YPOSTION
RET

```

```

LEFTD:
CALL   PRINOBJ
DEC    C
JP     UPDAA

```

```

RIGHTD:
CALL   PRINOBJ
INC    C
JP     UPDAA

```

UPD:

```

CALL PRINOBJ
DEC B
JP UPDAA

```

DOWND:

```

CALL PRINOBJ
INC B
JP UPDAA

```

```

XPOS EQU 0 ; XPOS OFFSET
YPOS EQU 1 ; YPOS OFFSET
VECTCN EQU 2 ; VECTOR COUNT OFFSET

NUM EQU 9 ; NUMBER OF SHIPS
LEN EQU 3 ; LENGTH OF DATA FOR TABLE

```

```

; SHIP TABLE
; 3 BYTES PER SHIP
; 1ST BYTE =X CO-ORD
; 2ND BYTE =Y CO-ORD
; 3RD BYTE =VECTOR COUNT

```

SHIPTB:

```

DB 100, 100, 8
DB 120, 80, 7
DB 55, 45, 14
DB 30, 30, 20
DB 40, 30, 1
DB 130, 130, 24
DB 140, 140, 20
DB 140, 118, 2
DB 140, 150, 2

```



```

START:
      DI                      ;DISABLE INT
      LD      B, NUM          ;LOAD B REGISTER WITH
                               ;NUMBER OF SHIPS
      LD      IY, SHIPTB     ;IY POINTS TO START
                               ;OF SHIP TABLE
DRAW:  PUSH    BC            ;SAVE SHIP COUNTER
      LD      B, (IY+YPOS)    ;GET Y CO-ORD
      LD      C, (IY+XPOS)    ;GET X CO-ORD
      CALL   PRINOBJ         ;AND PRINT
      LD      DE, 3          ;DE CONTAINS OFFSET
      ADD    IY, DE          ;POINT TO NEXT SHIP'S DATA
      POP    BC             ;RESTORE COUNTER
      DJNZ   DRAW           ;DRAW NEXT SHIP

MOVE:
      LD      IY, SHIPTB     ;POINT TO SHIP TABLE
      LD      B, NUM          ;NUMBER OF SHIPS

NXT:
      PUSH   BC             ;SAVE COUNTER
      LD     B, (IY+YPOS)    ;GET Y CO-ORD
      LD     C, (IY+XPOS)    ;GET X CO-ORD
      CALL  GETDIR          ;GET DIRECTION AND MOVE
      LD     DE, 3          ;PLACE OFFSET IN DE
      ADD   IY, DE          ;AND POINT TO NEXT SHIPM DATA
      POP   BC             ;RESTORE SHIP COUNTER
      DJNZ  NXT            ;MOVE NEXT SHIP
      JR    MOVE           ;FOREVER AND SO ON.....
      END

```

Hexadecimal Listing

```

6D60  C3  5D  6E  16  FF  16  0F  0F
6D68  16  FF  16  C5  CD  71  6D  C1
6D70  C9  DD  21  63  6D  CD  AA  22
6D78  5F  16  08  A7  28  1C  43  DD
6D80  7E  00  0E  00  CB  3F  CB  19
6D88  10  FA  AE  77  23  79  AE  77
6D90  2B  DD  23  CD  A8  6D  15  20
6D98  E5  C9  42  DD  7E  00  AE  77

```

6DA0	CD	A8	6D	DD	23	10	F4	C9
6DAB	24	7C	E6	07	C0	7C	D6	08
6DB0	67	7D	C6	20	6F	D0	7C	C6
6DB8	08	67	EE	58	C0	26	40	C9
6DC0	08	08	08	08	08	08	08	08
6DC8	08	08	08	08	01	01	01	01
6DD0	01	01	01	01	01	01	01	01
6DD8	04	04	04	04	04	04	04	04
6DE0	04	04	04	04	02	02	02	02
6DE8	02	02	02	02	02	02	02	02
6DF0	FF	21	C0	6D	FD	5E	02	16
6DF8	00	19	7E	FE	FF	20	05	AF
6E00	FD	77	02	C9	FD	34	02	FE
6E08	08	CA	34	6E	FE	04	CA	3B
6E10	6E	FE	01	CA	2D	6E	FE	02
6E18	CA	26	6E	C9	CD	6B	6D	FD
6E20	71	00	FD	70	01	C9	CD	6B
6E28	6D	0D	C3	1C	6E	CD	6B	6D
6E30	0C	C3	1C	6E	CD	6B	6D	05
6E38	C3	1C	6E	CD	6B	6D	04	C3
6E40	1C	6E	64	64	08	78	50	07
6E48	37	2D	0E	1E	1E	14	28	1E
6E50	01	82	82	18	8C	8C	14	8C
6E58	76	02	8C	96	02	F3	06	09
6E60	FD	21	42	6E	C5	FD	46	01
6E68	FD	4E	00	CD	6B	6D	11	03
6E70	00	FD	19	C1	10	EE	FD	21
6E78	42	6E	06	09	C5	FD	46	01
6E80	FD	4E	00	CD	F1	6D	11	03
6E88	00	FD	19	C1	10	EE	18	E6
6E90	6E	FE	01	CA	2D	6E	FE	02
6E98	CA	26	6E	C9	CD	6B	6D	FD
6EA0	71	00	FD	70	01	C9	CD	6B
6EA8	6D	0D	C3	1C	6E	CD	6B	6D
6EB0	0C	C3	1C	6E	CD	6B	6D	05
6EB8	C3	1C	6E	CD	6B	6D	04	C3
6EC0	1C	6E	64	64	08	78	50	07
6EC8	37	2D	0E	1E	1E	14	28	1E
6ED0	01	82	82	18	8C	8C	14	8C
6ED8	76	02	8C	96	02	F3	06	09

```

6EE0  FD  21  42  6E  C5  FD  46  01
6EE8  FD  4E  00  CD  6B  6D  11  03
6EF0  00  FD  19  C1  10  EE  FD  21
6EF8  42  6E  06  09  C5  FD  46  01

```

There are many improvements that could be made to this routine. You could easily add other directions, other direction tables and other characters. In fact if you want to be more ambitious you could make the objects of variable width and height. In addition improve the method of placing data on the screen by DRAWing from the bottom upwards instead of from the top downwards. This method of DRAWing objects will reduce flicker from the raster by catching the object as it is being DRAWn. All you have to work out is a routine similar to INCY but find the next pixel line above for a given screen address. I do know of a couple of ways to achieve this but I am not going to spoil your fun by explaining it to you!

11 Interrupts on the Spectrum

Have you ever wished that your computer could execute more than one program at once? Well, this chapter will explain how, in effect, you can double the power of your Spectrum by seemingly *RUNNING* two programs at once!

Interrupts on the Z80 chip serve similar purposes to those on other processors. They tell the computer that an external device, such as a disk drive, printer, keyboard or modem requires some attention. Take, as an example, the case where we have linked up a printer printing out data to our computer.

There are two ways of checking whether the printer is ready to get a character from the microprocessor. The inefficient way is to use a loop which has a description like this:

WAIT:

 IS PRINTER READY?

 ANSWER=NO THEN GO TO WAIT

 ANSWER=YES THEN GET NEXT CHARACTER: SEND IT
 TO THE PRINTER: GO TO WAIT

As you can see the above method 'polls' the printer continually to see if it is ready for the next character. Most of its time is spent in this loop waiting for the printer, so a lot of CPU time is wasted! Wouldn't it be fine if we could continue with other parts of the program and only send characters when the printer is ready? Well we can by using interrupts! Your Spectrum uses interrupts to get characters from the keyboard and update the frames system variable.

What is happening on a Spectrum is that your computer is running Spectrum BASIC. Frequently, (1/50 of a second to be precise or 1/60 of a second in N. America) it remembers where it is and what line it is running. It also recalls what address it is executing in the ROM or RAM and executes a routine in ROM which scans the keyboard. After it has done this it will go back to the address it was executing prior to interruption.

On the Z80 processor there are four kinds of interrupts. These interrupts are split into two categories called non-maskable and maskable interrupts. We shall be looking at just two of the maskable interrupts called mode 1 and mode 2 interrupts. (maskable means we can switch the interrupts off if we wish.)

mode 1 interrupts

Every time an interrupt occurs the processor pushes the current program counter onto the stack and jumps to location 0038 hex.

To exit out of this interrupt we must use a 'RET' (return) or 'RETI' (return from interrupt) instruction.

This mode of interrupt is actually the one used by the Spectrum during the scan for a key routine as described above.

mode 2 interrupts

This is the most powerful of the interrupts on the Z80 processor and is sometimes known as vectored processing.

In a mode 2 interrupt the programmer can specify up to 128 interrupts for other external devices.

This mode of interrupt revolves round a table which can contain up to 128 addresses. We can also have more than one table to deal with other external devices.

The start of a table is always on a page boundary of a 256 byte section of memory, i.e. 000H, 100H, 200H, C200H, etc. To tell the processor where the vector table is we LOAD the I (Interrupt) register with the high byte of the page number. For example if our vector table was at location C000 hex then we would tell the processor by executing:

```
DI          ;DISABLE INTERRUPTS
IM 2       ;SET UP INTERRUPT MODE 2
LD I,C0H   ;LOAD I REGISTER WITH C0 HEX
EI          ;ENABLE INTERRUPTS
```

Note that we only need to specify the high byte as we are dealing with page boundaries. The second line of code tells the processor that we want to use mode 2 interrupts. The last instruction turns on the scanning of interrupts. If we wished to ignore any maskable interrupts at any time we would use the instruction:

```
DI          ;DISABLE INTERRUPTS
```

But wait! There is no instruction which allows us to LOAD the I register directly with a number. We can only LOAD the I register with the A register. We overcome this problem by using:

```
LD A,C0H ;LOAD THE A REGISTER WITH C0H
LD I,A   ;AND PUT IT INTO THE I REGISTER.
```

In interrupt mode 2, when a device causes an interrupt, it provides an offset data number which is the low byte of the table. The offset points to a two byte address within the table to which the processor jumps (after first, stacking its current Program Counter).

```
TABLE EQU      C000H

IM        2
LD        A,C0H
LD        I,A      ;I REGISTER IS LOADED WITH HIGH
                   BYTE OF TABLE

EI
.
.
.

TABLE: DEFW    KEYBROU ;ADDRESS OF KEYBOARD ROUTINE
DEFW      PRINROU    ;ADDRESS OF PRINTER ROUTINE
.
.
.

        PUT ANY OTHER VECTORS HERE FOR OTHER DEVICES

KEYBROU:
.
.
.
RET

PRINROU:
.
.
.
RET
```

Now if an interrupt occurs and the data supplied is 0 then the processor pushes the current program counter on the stack and jumps to the address at C000H. If the data supplied was 02 then it would jump to the address at C002H which is the printer routine.

What if the data supplied was 01? If that happens a crash is likely to occur! Do you know why? The programmer has to program the

device to RETURN a valid data vector with its lowest bit set to zero i.e. always even!

Now what device can we program to cause an interrupt on a Spectrum? How do we program its eight bit vector number? The answer is we don't have too!! The Spectrum is not in conventional interrupt programming. Every 1/50 of a second an interrupt is generated by the ULA, one of the chips inside the computer. And at the time of the interrupt the data 0ff hex is passed to the microprocessor. If in interrupt mode 2 this will cause a jump to the address of the vector table currently pointed at by the I register plus 256 bytes.

Example:

```

                ORG  0C000H

INTINT:  LD   A,0C0H           ;SET TABLE AT PAGE 0C0H
         LD   I,A             ;AND PLACE IN THE I REG
         IM   2               ;SET UP FOR INTERRUPT
                                   ;MODE 2
         EI                   ;AND ENABLE
         RET

                ORG  0C0FFH           ;PLACE VECTOR AT PAGE+0FFH
         DEFW INTROU          ;ADDRESS OF
                                   ;INTERRUPT ROUTINE

INTROU:
         DI                   ;STOP ANY MORE INTERRUPTS
         PUSH HL              ;SAVE HL
         LD   HL,5B00H        ;POINT TO THE ATTRIBUTE
                                   ;FILE
         LD   (HL),255        ;AND SHOW SOME COLOUR
         POP  HL              ;RESTORE HL
         JP   0038H           ;JUMP BACK TO BASIC.
    
```

In our first example we notice that in our interrupt routine we disable the interrupt. This is only necessary when an interrupt is longer than 1/50 of a second so that we don't interrupt an interrupt! The JP 0038H jump to 0038 hex is the jump to BASIC keyboard scan routine. This scans the keyboard, updates the frame count and then

enables the interrupts. If we didn't wish to RETURN to BASIC then we would end an interrupt routine with:

```

        EI      ;ENABLE INTERRUPT
        RET     ;RETURN FROM INT
OR      EI
        RETI

```

Due to a hardware quirk in the Spectrum the value of the I register is limited to certain values 0-16 and 32-64. This means that for the 16k models we can only have our vector table in the page address 0-16 which is ROM!

There is, however, an end of page value which has a two byte value jumping out to RAM. This is page 28 hex.

LINKING THE INTERRUPT WITH THE RASTER

Re-set your Spectrum and type in the following program:

```

        ORG     30000

RASTER: HALT
        LD      A, 1
        OUT    (&FEH), A
        LD      HL, 500H          ;**EXPERIMENT**
                                   ;WITH THIS VALUE!

LOOP:   DEC     HL
        LD      A, L
        OR     H
        JR     NZ, LOOP

        LD      A, 2
        OUT    (&FEH), A
        JR     RASTER

        END

```

Here's a hexadecimal listing of the same program:

```

7D00  76 3E 01 D3 FE 21 00 05
7D08  2B 7D B4 20 FB 3E 02 D3
7D10  FE 18 ED 00

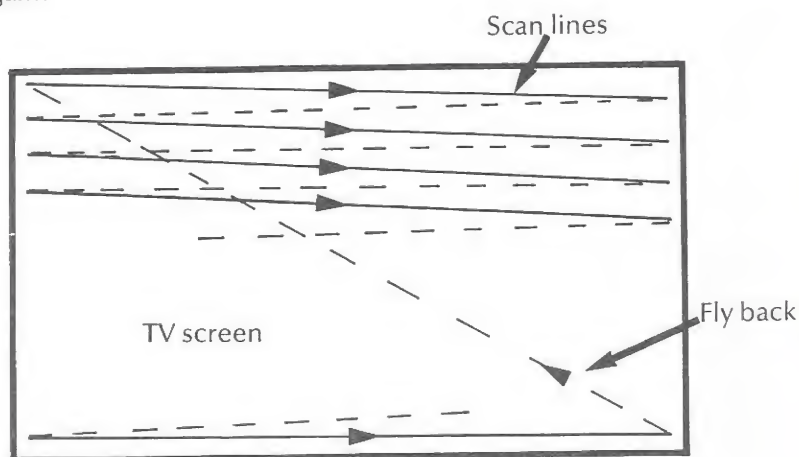
```


When you run the program you should get a BORDER split into two colours blue and red

The HALT instruction on the Z80 is used to wait for an interrupt. The computer will wait at a HALT instruction until some external device causes an interrupt. In the case of the Spectrum the ULA causes the interrupt. Therefore, the effect of the HALT instruction is to wait 1/50th of a second. Of course if we disabled all interrupts by using the instruction DI then the computer would wait for ever unless a Non-maskable interrupt (one that can not be disabled) was activated. In our program we use the HALT instruction to link with the raster beam to cause a split in the BORDER colour. Objects can be drawn when the raster is at the top or flying back thus reducing screen flicker. A lot of game programmers use this technique when writing fast arcade games.

Try pressing the keys when RUNNING this program. Notice how the BORDERS go up and down. Do you know why? It is due to the keyboard routine (which is called by the interrupt routine) taking different lengths of time to execute depending on which keys it finds pressed.

Every 1/50th of a second the computer REDRAWs the screen. The screen is updated by an election beam which scans across the pixels turning them on or off if they are set or re-set. The beam starts from the top left hand side of the screen and scans left to right across each line. After reaching the bottom the beam (or raster) flies diagonally back to the top left where it starts to update the screen again.



Interrupts are a powerful feature of the Z80 processor and must be used with care. It is not always true that the data RETURN for the low byte of the vector is 0FFH if some other device is on the back of the

Spectrum. When an interrupt occurs with a Kempston joystick on the Spectrum the data on the databus is the actual data returned from `READING` the joystick. In order to overcome this problem we could fill our vector table with an address in which the low and high bytes are the same: `8080h 7777h 1616h`.

Remember that on the interrupt the PC register will jump to the location in the table. We could put our interrupt routine at that address or have another jump instruction. Unfortunately, because of hardware problems limiting the value of the I register, this method is impossible on the 16k model Spectrums.

We'll now give you listings for two routines that use interrupts. The first is a `TRACE` program that can be used to help you debug `BASIC` programs. Every 1/50th of a second an interrupt occurs causing the transfer of the program counter to address 38 hex. This is where the keyboard is scanned and other 'housekeeping' tasks performed. We can cause the interrupt transfer to point to a routine which looks at the system variables `PPC` and `SUBPPC`. These contain the line number and statement number which `BASIC` is currently executing. `PRINTING` the values of these variables to the screen tells the `BASIC` programmer which statements the program is interpreting and the sequence of execution. This provides an extremely valuable aid for mapping the flow of the program, which in turn can greatly assist debugging.

This is not like the true `TRACE` functions found on some other computers since we can only see what line we are executing every 1/50th of a second. Some of the faster `BASIC` statements could be missed. The trace function is enabled by typing the instruction `RAND USR 32330` and disabled by typing `RAND USR 32338`. When you are running `BASIC` and the trace function is enabled then the line you are executing at the time of the interrupt is displayed on the top left hand side of the screen. To slow `BASIC` down as it is executing with the `TRACE` function press the 'Q' key. This is most useful as it sometimes gets difficult to see the line numbers being `PRINTED`.

Assembler Listing

```

ORG      32330D

;TRACE ROUTINE
;FOR 16K SPECTRUM

ZEROADD EQU      15360      ;ADDRESS FOR NUMERIC DATA
PCC      EQU      23621      ;LINE NUMBER EXECUTED

```

```

SUBPCC EQU 23623 ; STATEMENT WITHIN LINE

TRON:
    LD A, 28H ; LOAD A WITH 28H
    LD I, A ; AND PLACE IN THE
                ; I REGISTER
    IM 2 ; SET UP INTERRUPT MODE 2
    EI ; AND ENABLE
    RET

TROFF: IM 1 ; INTERRUPT MODE 1
        RET

        ORG 7E5CH ; INTERRUPT ROUTINE
                ; STARTS HERE

TRACE: DI
        PUSH AF ; SAVE REGS
        PUSH BC
        PUSH DE
        PUSH HL
        PUSH IX

        LD HL, (PCC) ; LOAD PROGRAM POINTER
        LD A, H
        INC A
        JR Z, SKIP

        LD DE, 16384
        CALL CONV ; PRINT NUMBER ON SCREEN.
        INC DE

        LD A, (SUBPCC) ; GET SUB-LINE NUMBER.
                ; ONE BYTE NUMBER
        LD H, 0 ; AND TRANSFER TO
                ; HL REGISTER PAIR.

        LD L, A
        CALL CONV2
    
```

```

        LD      A, 0FBH

        IN      A, (0FEH)
        RRA
        JR      C, SKIP

        CALL    WAIT
        CALL    WAIT
SKIP:
        POP     IX
        POP     HL
        POP     DE
        POP     BC
        POP     AF
        JP      0038H

WAIT:   LD      HL, 00
        LD      DE, 00
        LD      BC, 00
        LDIR
        RET

PRDIGIT:                                ;HL POINTS TO SCREEN
                                           ;ADDRESS
                                           ;A CONTAINS DIGIT
                                           ;NUMBER 0-9

        PUSH   DE                        ;SAVE REGISTERS
        PUSH   BC
        PUSH   IX
        PUSH   HL

        LD     H, 0                      ;PUT CHARACTER OFFSET
        LD     L, A                      ;IN HL REGISTER

        ADD    HL, HL                    ;MULTIPLY BY 8
        ADD    HL, HL
        ADD    HL, HL

```

```

EX      DE, HL
LD      IX, ZEROADD      ;POINT TO START OF
                          ;NUMERIC DATA
ADD     IX, DE           ;ADD OFFSET TO START OF
                          ;NUMBER
EX      DE, HL
LD      B, B             ;SET COUNTER

NXDAT:  LD      A, (IX)   ;GET NUMERICAL DATA
        LD      (DE), A  ;PLACE ON SCREEN
        INC     D         ;ADJUST SCREEN ADDRESS
        INC     IX       ;ADJUST DATA POINTER.
        DJNZ   NXDAT     ;DO NEXT DATA

        POP     HL
        POP     IX
        POP     BC       ;RESTORE REGISTERS
        POP     DE
        RET

DECT2:  DEFW    1000
DECT3:  DEFW    100
        DEFW    10
        DEFW    1
CONV2:  LD      IX, DECT3
        JR      NDIG2
CONV:   LD      IX, DECT2
NDIG2:  LD      B, (IX+1)
        LD      C, (IX+0)
        LD      A, '0'-1
        AND     A
CAR:    INC     A
        SBC     HL, BC
        JR      NC, CAR
        ADD     HL, BC
        CALL   PRDIGIT
        INC     DE
        INC     IX

```

```

INC      IX
DEC      C
JR       NZ,NDIG2
RET

```

```

END

```

Hexadecimal Listing

```

7E4A    3E  28  ED  47  ED  5E  FB  C9
7E52    ED  56  C9  00  00  00  00  00
7E5A    00  00  F3  F5  C5  D5  E5  DD
7E62    E5  2A  45  5C  7C  3C  28  1D
7E6A    11  00  40  CD  CE  7E  13  3A
7E72    47  5C  26  00  6F  CD  C8  7E
7E7A    3E  FB  DB  FE  1F  38  06  CD
7E82    90  7E  CD  90  7E  DD  E1  E1

7E8A    D1  C1  F1  C3  38  00  21  00
7E92    00  11  00  00  01  00  00  ED
7E9A    B0  C9  D5  C5  DD  E5  E5  26
7EA2    00  6F  29  29  29  EB  DD  21
7EAA    00  3C  DD  19  EB  06  08  DD
7EB2    7E  00  12  14  DD  23  10  F7
7EBA    E1  DD  E1  C1  D1  C9  E8  03
7EC2    64  00  0A  00  01  00  DD  21

7ECA    C2  7E  18  04  DD  21  C0  7E
7ED2    DD  46  01  DD  4E  00  3E  2F
7EDA    A7  3C  ED  42  30  FB  09  CD
7EE2    9C  7E  13  DD  23  DD  23  00
7EEA    20  E6  C9  CD  CE  7E  13  3A
7EF2    47  5C  26  00  6F  CD  C8  7E
7EFA    3E  FB  DB  FE  1F  38

```

The second of our two interrupt driven routines allows us to have an on-screen clock constantly telling us the time, even when we are RUNNING a BASIC program. After placing the machine code routine in memory key in the BASIC listing below. This serves to set the time on the clock. After RUNNING, the BASIC program clock should be constantly updated on the top right hand side of the screen. You can stop the clock at any time by entering NEW. This disables the clock by re-setting the interrupt mode to 1, thereby causing the Z80 to

branch off to 38 hex on every interrupt. To start the clock off again simply type RANDOMIZEUSR 32330

BASIC Listing

```

10 CLEAR 32325:LET T=32438
20 INPUT "HOURS";H:LET H=INT(H
): IF H<0 OR H>12 THEN GO TO 20
30 INPUT "MINS ";M:LET M=INT(M
): IF M<0 OR M>59 THEN GO TO 30
40 IF H>9 THEN LET H=H+6
50 IF M>9 THEN LET M=M+6*INT(M
/10)
60 POKE T,H:POKE T+1,M:POKE T+
2,0:RANDOMIZEUSR 32330

```

Assembler Listing

```

ORG 32330D

;CLOCK ROUTINE
;FOR 16K SPECTRUM

TRON:
LD A,28H ;SET UP I REGISTER
LD I,A ;TO PAGE 28 HEX
IM 2 ;SET INTERRUPT MODE 2
EI ;AND ENABLE
RET

FRAMES:
DB 0

ORG 7E5CH ;START OF INTERRUPT ROUTINE

CLOCK:
DI ;DISABLE INTERRUPTS
PUSH AF ;SAVE REGISTERS ON
PUSH BC ;THE STACK.
PUSH DE
PUSH HL
PUSH IX

LD A,(FRAMES) ;UPDATE 1/50 SECOND
;COUNTER.

```

```

    INC    A
    LD     (FRAMES), A
    CP     50                ; HAVE WE COUNTED THROUGH
                                ; 1 SEC?
    JR     NZ, PRCLOCK      ; NO, SO PRINT TIME ANYWAY.

    XOR    A                ; SET FRAMES
    LD     (FRAMES), A     ; TO ZERO.

    LD     DE, TIMLIM      ; GET BCD TIME LIMITS
    LD     HL, SECS       ; POINT TO TIMER COUNTERS
    LD     B, 3            ; NUMBER OF COUNTERS TO
                                ; UPDATE.

NXBCD:  LD     A, (HL)     ; GET TIME COUNTER
        ADD    A, 1       ; INCREASE BY ONE
        DAA                ; BCD
        LD     (HL), A
        LD     A, (DE)    ; GET LIMIT
        CP     (HL)      ; HAVE WE REACHED LIMIT FOR
                                ; THAT TIMER?

        JR     NZ, PRCLOCK

; NO SO GO AND PRINT TIME.
        LD     (HL), 0    ; RESET TIME COUNTER
        INC    DE         ; POINT TO NEXT TIME LIMIT
        DEC    HL         ; POINT TO NEXT TIME COUNTER
        DJNZ  NXBCD      ; DO NEXT DIGIT.

; AT THIS POINT THE HOURS ARE RESET TO ZERO.

        INC    HL         ; POINT TO HOURS DIGIT
        INC    (HL)      ; SET TO 1 O'CLOCK.

PRCLOCK:
        ; ROUTINE TO PRINT THE CLOCK ON THE SCREEN

        LD     HL, 16384+31-8 ; TOP RIGHT HAND CORNER
        LD     DE, HRS      ; DE POINTS TO BCD DIGITS
        LD     B, 3        ; COUNTER

```



```

NXT:  LD    A,(DE)      ;GET DIGIT
      LD    C,A        ;SAVE IN C REG
      AND   0F0H      ;GET FIRST DIGIT
      RRCA                ;MOVE DOWN TO BITS 0-3
      RRCA
      RRCA
      RRCA
      CALL  PRDIGIT    ;PRINT DIGIT.
      INC  HL          ;POINT TO NEXT PART
                        ;OF SCREEN
      LD    A,C        ;GET DIGIT
      AND   0FH       ;MASK OFF BOTTOM 4 BITS.
      CALL  PRDIGIT    ;PRINT DIGIT
      INC  DE          ;POINT TO NEXT
                        ;TWO BCD DIGITS.
      INC  HL          ;ONE SPACE BETWEEN DIGITS.
      INC  HL
      DJNZ NXT        ;DO NEXT DIGITS

      POP  IX          ;RESTORE REGISTERS
      POP  HL
      POP  DE
      POP  BC
      POP  AF
      JP   0038H
    
```

```

TIMLIM: DB    60H,60H,13H
HRS:    DB    0
MINS:   DB    0
SECS:   DB    0
    
```

```

ZEROADD EQU    15744 ;ADDRESS OF START OF
                    ;NUMERIC DATA
    
```

```

PRDIGIT:
                    ;HL POINTS TO SCREEN
                    ;ADDRESS
                    ;A CONTAINS DIGIT
                    ;NUMBER 0-9

      PUSH HL        ;SAVE REGISTERS
      PUSH BC
    
```

```

SLA    A           ;MULTIPLY DIGIT BY 8
SLA    A
SLA    A
LD     B, 0       ;GET OFFSET
LD     C, A       ;IN BC REGISTER
LD     IX, ZEROADD ;POINT TO START OF
                        ;NUMERIC DATA
ADD    IX, BC     ;ADD OFFSET TO START OF
                        ;NUMBER

LD     B, 8       ;SET COUNTER

NXDAT: LD     A, (IX) ;GET NUMERICAL DATA
LD     (HL), A    ;PLACE ON SCREEN
INC    H          ;ADJUST SCREEN ADDRESS
INC    IX         ;ADJUST DATA POINTER.
DJNZ   NXDAT     ;DO NEXT DATA

POP    BC         ;RESTORE REGISTERS
POP    HL
RET

END

```

Hexadecimal Listing

```

7E4A  3E  28  ED  47  ED  5E  FB  C9
7E52  00  00  00  00  00  00  00  00
7E5A  00  00  F3  F5  C5  D5  E5  DD
7E62  E5  3A  52  7E  3C  32  52  7E
7E6A  FE  32  20  1D  AF  32  52  7E
7E72  11  B3  7E  21  B8  7E  06  03
7E7A  7E  C6  01  27  77  1A  BE  20
7E82  08  36  00  13  2B  10  F1  23

7E8A  34  21  17  40  11  B6  7E  06
7E92  03  1A  4F  E6  F0  0F  0F  0F
7E9A  0F  CD  B9  7E  23  79  E6  0F
7EA2  CD  B9  7E  13  23  23  10  E9
7EAA  DD  E1  E1  D1  C1  F1  C3  38
7EB2  00  60  60  13  00  00  00  E5
7EBA  C5  CB  27  CB  27  CB  27  06
7EC2  00  4F  DD  21  80  3D  DD  09

```

7ECA	06	08	DD	7E	00	77	24	DD
7ED2	23	10	F7	C1	E1	C9	00	00
7EDA	00	00	F3	F5	C5	D5	E5	DD
7EE2	E5	3A	52	7E	3C	32	52	7E
7EEA	FE	32	20	1D	AF	32	52	7E
7EF2	11	B3	7E	21	B8	7E	06	03
7EFA	7E	C6	01	27	77	1A		

12 Machine code miscellany

In this final chapter, I present a complete machine code game and a variety of routines. You can use these to enhance your BASIC programs or incorporate into your own machine code programs. The techniques we've seen in the course of the book are all represented here. Study of the programs should help you in writing your own and will increase your repertoire of routines and your program library.

BRICKOUT

This version of a venerable arcade game is one of my favourite programs. There is no better way to become proficient in machine code than by writing games. The object of this game is to knock three layers of 32 bricks away from the top part of the screen. The player controls a bat and directs a ball moving along the screen to knock down the coloured bricks. If the ball passes the bat then the player loses a life. The game ends when the player loses all his lives. The game listing has been broken into sections and heavily annotated to help you see the structure of the program.

The bat is controlled by the <z> and <x> keys which make it move left or right respectively. The <CAPS SHIFT> key can be used to make the bat go twice as fast, (a 'Cheat Key' if you like!). The game program comes in two parts. One is in BASIC and the other in machine code. To start the game after entering the code you RUN the BASIC program which calls the machine code routine. When the player has lost the score is printed on the RETURN to BASIC as the variable SC and the user is asked if he wants another game. Here's the BASIC program:

```
10 PRINT #0;"press a key when  
ready"  
20 IF INKEY$="" THEN GO TO 20  
30 CLS  
40 RANDOMIZE USR 30000  
45 PLOT 0,175: DRAW 255,0: PLO  
T 0,0: DRAW 0,175  
50 GO TO 10
```

The machine code routine comes in three main sections. The first initializes the score, the number of bricks left, the number of balls left and draws the screen. The second routine, MOVBAT, moves the user's bat, controlled by the <Z>, <X> and <CAPS SHIFT> keys. The last routine, MUBALL, deals with moving the ball around the screen, knocking out bricks, rebounding off the bat and walls, and updating the score. We'll break down the assembler listing, and give the whole hexadecimal listing at the end.

To start off the game we jump into the portion of code labelled START. This follows the initialisation:

```

                ORG      32000
                JP       START
BATYX:  DEFW      160FH      ;BATS POSITION
BALLX:  DB        10H       ;BALLS X POSITION
BALLY:  DB        01H       ;BALLS Y POSITION
TPBLYX:
                DEFW      0           ;TEMP AREA
XINC:   DB        1           ;X MOVEMENT
YINC:   DB        1           ;Y MOVEMENT
LEVEL:  DB        4           ;LEVEL
                ;NUMBER OF HALTS FOR DELAY
SCORE:  DEFW      0           ;SCORE
BALLS:  DB        0           ;NUMBER OF BALLS
HITS:   DB        0           ;NUMBER OF BRICKS HIT

PATTBL:

SPACE:  DB        0, 0, 0, 0, 0, 0, 0, 0

;DATA FOR BALL
BALLCH: DB        3CH, 7EH, 0FFH, 0FFH, 0FFH, 0FFH, 7EH, 3CH

```

; DATA FOR BAT

```

BATCHS: DB      3FH, 7FH, 0FFH, 0FFH, 0FFH, 0FFH, 7FH, 3FH
           DB      0FFH, 0FFH, 0FFH, 0FFH, 0FFH, 0FFH, 0FFH, 0FFH
           DB      0FCH, 0FEH, 0FFH, 0FFH, 0FFH, 0FFH, 0FEH, 0FCH

```

; DATA FOR BRICK

```

BRICK1: DB      0FFH, 81H, 81H, 81H, 81H, 81H, 81H, 0FFH

```

START:

```

LD      A, 2
CALL   1601H      ; OPEN CHANNEL TWO
XOR    A          ; SET A REGISTER TO ZERO
OUT    (0FEH), A ; SET BORDER TO BLACK

LD      HL, 0      ; RE-SET SCORE TO 0
LD      (SCORE), HL
LD      A, 5
LD      (BALLS), A ; SET NUMBER OF BALLS TO 5
LD      A, 4
LD      (LEVEL), A
LD      A, 96      ; NUMBER OF BRICKS
LD      (HITS), A  ; PLACE IN HITS
CALL   SETUP      ; SET UP BRICKS ON SCREEN
CALL   RNOBAL      ; PLACE THE BALL ON THE
                       ; SOMEWHERE ON THE SCREEN
LD      HL, 1610H  ; INITIALIZE BATS X, Y CO-ORDS
LD      (BATYX), HL
CALL   PRTBAT      ; AND PRINT THE BAT

```

BATAGN:

```

LD      A, (BALLS) ; LOOK AT THE NUMBER
                       ; OF BALLS LEFT
AND    A          ; IS IT ZERO?
JR     Z, GMOVE   ; IF IT IS GO TO
                       ; DEAL WITH END OF GAME

```

; WE ARE STILL PLAYING

```

CALL   MOVBAT      ; MOVE BAT
CALL   MUBALL      ; MOVE BALL

```

```

EI
HALT                ;WAIT FOR 1/50 OF A SECOND
DI
JR      BATAGN      ;KEEP PLAYING

```

```

GMOVER:
LD      BC, ( SCORE ) ;PASS SCORE TO
                        ;BC REGISTER PAIR
EI
RET     ;ENABLE INTERRUPTS
        ;AND RETURN TO BASIC

```

The game ends when there are no balls left, which causes the routine to jump to the label `GMOVER`. This gets the current score and places it into the `BC` register pair to be passed back to `BASIC`.

`MOVBAT` is used to control the movement of the bat according to the keys `<Z>` or `<X>` being pressed. If the user presses the key `<CAPS SHIFT>` then the routine goes back to the label `FIVE90` to move the bat again.

A call is made to either `RIGHTTB` or `LEFTTB` to move the bat right or left. After this the routine `RETURNS` to move the ball.

```

MOVBAT:
                        ;MOVE PLAYERS BAT

CALL    PRTBAT        ;PRINT BAT

LD      A, 0FEH       ;SET UP TO SCAN BOTTOM
IN      A, ( 0FEH )   ;LEFT HAND SIDE OF KEYBOARD
AND     1FH           ;MASK OFF LOWER FOUR BITS
CP      1FH           ;SEE IF ALL BITS ARE SET
RET     Z             ;NO KEY PRESSED SO RETURN

CALL    CLRBAT        ;CLEAR BAT OFF SCREEN
FIVE90:
BIT     1, A          ;IF PRESSED 'Z'
CALL    Z, LEFTB      ;THEN MOVE LEFT
BIT     2, A          ;IF PRESSED 'X'
CALL    Z, RIGHTB     ;THEN MOVE RIGHT
BIT     0, A          ;HAVE WE PRESSED SHIFT
JR      NZ, BATPRT    ;KEY?IF NOT JUST PRINT BAT
SET     0, A
JR      FIVE90        ;TURN OFF SHIFT KEY
                        ;HAVE ONE MORE GO

```

BATPRT:

```

CALL    PRTBAT      ;PRINT BAT ON SCREEN
RET     ;AND RETURN

```

When moving the ball left or right a check must be made to make sure that the bat does not go off the screen. The variable **BATXY** holds the X,Y co-ordinate of the left hand side of the bat. The bat is made up of three characters.

RIGHTB:

```

PUSH    AF          ;GOING RIGHT, SAVE AF PAIR
LD      HL,(BATYX)  ;GET X, Y CO-ORD OF BAT
                          ;IN HL PAIR
LD      A,10H      ;LOAD A REGISTER WITH 29
CP      L           ;TEST TO SEE IF WE HAVE
                          ;HIT THE RIGHT SIDE
JR      Z,REDGE     ;HIT, SO DON'T UPDATE
INC     L           ;INCREASE X CO-ORD
LD      (BATYX),HL ;AND SAVE

```

REDGE:

```

POP     AF          ;RESTORE KEY STATUS
RET     ;AND RETURN

```

LEFTB:

```

PUSH    AF          ;GOING LEFT, SAVE KEY STATUS
LD      HL,(BATYX)  ;GET X, Y CO-ORDS
LD      A,L         ;TEST IF HIT LEFT HAND SIDE
AND     A           ;IE IF EQUAL TO 0
JR      Z,LEDGE     ;HIT SO DON'T UPDATE
DEC     L           ;DECREASE ONE OFF X CO-ORD
LD      (BATYX),HL ;AND SAVE

```

LEDGE:

```

POP     AF          ;RESTORE KEY STATUS
RET     ;AND RETURN

```

The routine **CLRBAT** is used to remove the bat from the screen. To do this we **PRINT** the character **SPACE** which consists of zeros. While the routine **PRTBAT** is used to **PRINT** the bat to the screen. Both these routines call the routine **PRTCH** which **PRINTS** the character held in the **A** register. In this routine a **CALL** is made to two **ROM** routines. The routine at **0E9E** hex calculates the screen address for a given **Y** co-ordinate. The routine at the address **0E88** hex calculates the attribute in the **DE** register pair for a given screen address.


```

CLRBAT:
    PUSH    AF                ;SAVE AF REGISTER
    LD      HL,(BATYX)       ;GET X, Y CO-ORD
    LD      BC,338H         ;SET INK AND PAPER
                                ;WHITE PAPER BLACK INK
CLRIT:
                                ;B REGISTER IS
                                ;LOADED WITH 3
    PUSH    BC                ;SAVE CHAR CODE
    PUSH    HL                ;SAVE X, Y CO-ORD AND
                                ;COUNTER
    XOR     A                 ;SET A TO ZERO
    CALL    PRTCH            ;PRINT SPACE
    POP     HL                ;RESTORE X, Y
    INC     L                 ;POINT TO NEXT CHAR OF BAT
    POP     BC                ;RESTORE X, Y CO-ORD
                                ;AND COUNTER
    DJNZ   CLRIT             ;RUB OFF 3 CHARACTERS
    POP     AF                ;RESTORE AF REGISTER
    RET

PRTBAT:
    LD      HL,(BATYX)       ;GET X, Y CO-ORD
    LD      BC,339H         ;SET B=3 AND COLOUR TO
                                ;WHITE PAPER AND RED INK
    LD      A,2              ;INITIALIZE A REG TO FIRST
                                ;CHARACTER OF BAT

NEXBAT:
    PUSH    BC                ;SAVE COLOUR AND COUNTER
    PUSH    HL                ;SAVE X, Y CO-ORD
    CALL    PRTCH            ;PRINT PART OF BAT
    INC     A                 ;NEXT PART OF BAT
    POP     HL                ;RESTORE X, Y
    INC     L                 ;NEXT X POSITION OF BAT
    POP     BC                ;RESTORE COUNTER AND COLOUR
    DJNZ   NEXBAT           ;DO 3 TIMES
    RET

PRINTCHAR:
    ;H=Y L=X A=CHAR NUMBER C=COLOUR

```

PRTCH:

```

      PUSH    AF
      PUSH    BC           ;SAVE CHARACTER
      PUSH    HL           ;SAVE COLOUR
                          ;SAVE X, Y CO-ORDS

      PUSH    BC
      PUSH    AF           ;SAVE COLOUR
      PUSH    HL           ;SAVE CHARACTER
                          ;SAVE X, Y CO-ORDS

      LD      A, H
      CALL    0E9EH        ;LOAD A WITH Y CO-ORD
      POP     DE           ;CALCULATE SCREEN ADDRESS
      LD      D, 0         ;PLACE X CO-ORD IN E REG
      ADD     HL, DE       ;PLACE 0 IN D
      EX      DE, HL      ;FIND SCREEN ADDRESS
                          ;AND PLACE IN DE

      POP     AF           ;GET CHARACTER CODE
      LD      BC, PATTBL   ;BC POINTS TO CHARACTER SET
      LD      H, 0         ;LOAD H WITH 0
      LD      L, A         ;LOAD A WITH CHARACTER
                          ;NUMBER

      ADD     HL, HL       ;TIMES BY 2
      ADD     HL, HL       ;TIMES BY 4
      ADD     HL, HL       ;TIMES BY 8
      ADD     HL, BC       ;ADD CHARACTER
                          ;TABLE ADDRESS
                          ;HL NOW POINTS
                          ;TO CHARACTER DATA

      LD      B, 8         ;LOAD B WITH DATA COUNT

NXTROW: LD      A, (HL)    ;GET CHARACTER DATA
      LD      (DE), A     ;AND PLACE ON SCREEN
      INC     HL           ;POINT TO NEXT CHARACTER
                          ;DATA

      INC     D           ;POINT TO NEXT PIXEL
                          ;LINE ON THE SCREEN

      DJNZ   NXTROW      ;DO THIS UNTILL
                          ;WE HAVE FINISHED
                          ;PRINTING THE CHARACTER

      EX     DE, HL       ;LET HL NOW BE
                          ;THE SCREEN ADDRESS

      CALL   0E88H        ;CALCULATE THE
                          ;ATTRIBUTE ADDRESS

      POP    BC           ;RESTORE COLOUR CODE

```

```

LD      A, C           ;PLACE IN A REGISTER
LD      (DE), A       ;SET ATTRIBUTE
POP     HL             ;RESTORE X, Y CO-ORD
POP     BC             ;RESTORE COLOUR CODE
POP     AF             ;RESTORE CHARACTER
RET                                           ;RETURN FROM PRINTING

```

The routine SETUP is called only once: at the start of each new game. It is used to draw the bricks on the screen.

```
;SET START SCREEN
```

```
SETUP:
```

```

CALL    0D6BH         ;CLEAR SCREEN
LD      BC, 2020H     ;32 GREEN BRICKS
LD      A, 5          ;PLACE BRICK CHAR IN A REG
LD      HL, 300H      ;START X, Y CO-ORD OF BRICKS
CALL    NXCOL         ;DRAW BRICKS

LD      BC, 2018H     ;COLOUR =18H MAGENTA
LD      HL, 400H      ;Y=4  X=0
CALL    NXCOL         ;DRAW BRICKS

LD      BC, 2030H     ;COLOUR =30H YELLOW
LD      HL, 500H      ;Y=5, X=0

```

```
NXCOL:
```

```

CALL    PRCH         ;PRINT BRICK
INC     L             ;POINT TO NEXT X CO-ORD
DJNZ   NXCOL        ;REPEAT 32 TIMES

RET                                           ;RETURN

```

PEEK is the routine which is used to detect any collision between the ball and any bricks or the bat. The x and y co-ordinates are placed in the HL pair and after calling this routine the attribute or colour code is returned in the A register.

PEEK:

```

; RETURNS ATTRIBUTE
; OF GIVEN X, Y ( IN HL PAIR )
; IN A REGISTER
LD      A, L      ; PLACE X CO-ORD
; IN A REGISTER
LD      L, H      ; PLACE Y CO-ORD
; IN L REGISTER
LD      H, 0      ; 32 BIT NUMBER 50
; PLACE 0 IN H
ADD     HL, HL    ; TIMES BY 2
ADD     HL, HL    ; TIMES BY 4
ADD     HL, HL    ; TIMES BY 8
ADD     HL, HL    ; TIMES BY 16
ADD     HL, HL    ; TIMES BY 32
LD      B, 0      ; LOAD B REG WITH 0
LD      C, A      ; PLACE X CO-ORD IN C REGISTER
ADD     HL, BC    ; FIND OFFSET
LD      BC, 5800H
ADD     HL, BC    ; CALCULATE ATTRIBUTE
; ADDRESS
LD      A, ( HL ) ; GET CONTENTS OF
; THAT ADDRESS AND PLACE
; IN A REGISTER
RET     ; RETURN

```

BRKOUT deals with the ball colliding with an object. It determines which coloured brick (if any) it has hit and gives an appropriate score. When the ball hits a brick the variable HITS is deducted by one to keep a count of the number of bricks still left standing. It branches off to NOEND if any are still left.

If all the bricks are knocked down then the routine will reset the number of bricks by LOADING the variable HITS with 96 (i.e. three rows of 32 bricks). The player is rewarded by a bonus of two balls and the variable LEVEL is decreased. This controls the delay when moving the ball, thus increasing its speed for the next game.

BRKOUT:

```

; BRICK HAS HIT SOMETHING DO A TEST

LD      BC, 0
CP      30H      ; HAVE WE HIT A YELLOW BRICK
JR      NZ, NTYLW ; NOT YELLOW

```

```

LD      A, -1          ;SEND BALL IN OTHER DIRECTION
LD      (YINC), A
LD      BC, 2         ;ADD TO SCORE
JR      BEEP         ;AND MAKE A NOISE ABOUT IT!

NTYLW:
CP      18H          ;HAVE WE HIT A
                ;MAGENTA BRICK?
JR      NZ, NTMAGN   ;NOT MAGENTA
LD      A, -1        ;SEND BALL IN
                ;OTHER DIRECTION

LD      (YINC), A
LD      BC, 5        ;SCORE
JR      BEEP         ;AND MAKE A SOUND

NTMAGN:
CP      20H          ;HAVE WE HIT A GREEN BRICK?
JR      NZ, ERROR    ;GOD KNOWS WHAT WE HIT!
LD      BC, 10       ;GIVE HIM A BIG SCORE
JR      BEEP         ;AND MAKE A NOISE!

ERROR:
LD      DE, 40H
LD      HL, 666H
CALL   3B5H         ;MAKE A LONGER BEEP!

BEEP:
LD      HL, (SCORE)  ;GET SCORE
ADD     HL, BC       ;AND ADD 0, 5 OR 16
LD      (SCORE), HL ;SAVE UPDATED SCORE
LD      HL, HITS     ;POINT TO NUMBER OF HITS
DEC     (HL)        ;SUBTRACT ONE
JR      NZ, NOEND    ;ALL BRICKS HIT?

LD      (HL), 96     ;RESET NUMBER OF BRICKS
LD      A, (LEVEL)  ;GET LEVEL
AND     A           ;TEST FOR ZERO LEVEL
JR      Z, MAXLEV    ;DO NOT BOTHER MAKING
                ;ANY MORE DIFFICULT

DEC     A           ;ONE OFF THE LEVEL
LD      (LEVEL), A  ;AND SAVE
LD      A, (BALLS)  ;GET NUMBER OF BALLS
ADD     A, 2        ;AND GIVE HIM TWO MORE
LD      (BALLS), A  ;AND SAVE

```

```

MAXLEV:
        CALL    RNDBAL        ;GET A RANDOM BALL POSTION
        CALL    SETUP        ;SET UP THE WALL
NOEND:
        LD      DE, 8
        LD      HL, 666H
        CALL    3B5H         ;BEEP
        LD      A, 0         ;MAKE SURE
        OUT    (0FEH), A     ;WE HAVE A BLACK BORDER
        RET

```

The routine MUBALL is one of the main routines which deals with the movement of the ball. The ball has a y direction (YINC) and x direction (XINC). These two variables are offsets which are added to the ball's x and y co-ordinates. These are either 1 or -1. If the ball passes the bottom of the screen, one is deducted off the number of remaining balls. If there is still any left then a branch is made to RNDBALL which sets up another ball at a random x position. If the ball collides with an object than its x direction and/or y direction is reversed.

```

MUBALL:
        LD      HL, (BALLX)  ;GET BALLS X, Y CO-ORD
        ;BALL DOESNT GO THROUGH
        ;THE BAT...

        LD      A, (YINC)    ;GET Y DIRECTION
        ADD    A, H         ;ADD TO Y
        LD      H, A        ;AND SAVE NEW Y CO-ORD
        LD      A, (XINC)    ;GET X DIRECTION
        ADD    A, L         ;ADD TO CURRENT X CO-ORD
        LD      L, A        ;AND SAVE NEW X CO-ORD
        PUSH   HL          ;SAVE THIS
        CALL   PEEK        ;LOOK AT THE COLOUR
                        ;OF NEW X, Y
        POP    HL          ;RESTORE SCREEN ADDRESS
        CP     39H         ;IS IT THE BAT
        ;HIT?
        JR     NZ, NTBAT    ;NO ITS NOT!

        LD      A, (YINC)    ;REVERSE Y DIRECTION
        NEG
        LD      (YINC), A   ;AND SAVE
        ;SEMI REBOUND DIRN.

```

```

LD      A, (XINC)      ; REVERSE X DIRECTION
NEG
ADD     A, L           ; ADD X CO-ORD
LD      L, A           ; AND SAVE IN L
CALL   PEEK           ; LOOK AT COLOURS HERE
CP      38H           ; IS IT NORMAL BACKGROUND
JR      NZ, MUBALL    ; NO, THEN MOVE BALL

; SEND BALL BACK THE WAY IT CAME ...

LD      A, (BALLX)    ; GET X CO-ORD
AND     A              ; TEST FOR LEFT HAND SIDE
JR      Z, MUBALL     ; NO THEN MOVE BALL
LD      A, (XINC)
NEG
LD      (XINC), A     ; AND SAVE
JR      MUBALL        ; MOVE THE BALL

NTBAT:
LD      HL, (BALLX)   ; GET X, Y CO-ORD
LD      A, (XINC)     ; GET X DIRECTION
ADD     A, L           ; GET NEW X CO-ORD
LD      L, A           ; AND SAVE IN L REG
LD      (TPBLYX), A   ; PLACE NEW X CO-ORD IN TEMP
AND     A              ; IS IT AT THE
                          ; LEFT HAND SIDE?
JR      Z, NGXINC     ; YES THEN GO TO
                          ; CHANGE X DIRECTION
CP      1FH           ; IS IT ON THE
                          ; RIGHT HAND SIDE?
JR      C, YCHECK     ; NO SO CHECK Y MOVEMENT

NGXINC:
LD      A, (XINC)     ; GET X DIRECTION
NEG
LD      (XINC), A     ; AND SAVE

YCHECK:
LD      A, (YINC)     ; GET Y DIRECTION
ADD     A, H           ; GET NEW Y CO-ORD
LD      H, A           ; AND SAVE IN H REG
LD      (TPBLYX+1), A ; AS WELL AS TEMP+1
AND     A              ; HAVE WE HIT THE TOP?
JR      Z, NGYINC     ; YES THEN CHANGE
                          ; Y DIRECTION

```

```

CP      23                ; HAVE WE HIT THE BOTTOM?
JR      NC, BALOUT      ; OUT OF BOUNDS
CALL    PEEK            ; NO, LOOK AT WHERE
                        ; WE ARE GOING TO.
CP      38H             ; IS IT BLANK?
JR      Z, GOED         ; YES, CARRY ON
CALL    BRKOUT          ; HIT SOMETHING SO CHECK!

NGYINC:
LD      A, (YINC)       ; GET Y DIRECTION
NEG                      ; CHANGE DIRETCION
LD      (YINC), A       ; AND SAVE

GOED:
LD      A, 0            ; SET UP CHARACTER AS SPACE
; ERASE OLD
LD      C, 38H          ; SET UP BACKGROUND COLOUR
LD      HL, (BALLX)     ; GET BALLS X, Y CO-ORDS
CALL    PRITCH          ; AND BLANK OUT
LD      A, 1            ; GET BALL CHARACTER
LD      HL, (TPBLYX)    ; GET TEMP X, Y
LD      (BALLX), HL     ; AND SAVE IN BALLS
                        ; X, Y CO-ORDS
CALL    PRITCH          ; PRINT THE BALL!
LD      A, (LEVEL)     ; GET LEVEL

LEVL:
EI
HALT                      ; 1/50 OF A SECOND DELAY
DI
DEC A
JR      NZ, LEVLP      ; DELAY DEPENDENT ON LEVEL
RET

BALOUT:
LD      A, 0            ; SET UP CHARACTER AS SPACE
; ERASE BALL            ; BACKGROUND COLOUR
LD      C, 38H
LD      HL, (BALLX)    ; GET X, Y CO-ORDS
CALL    PRITCH          ; AND PRINT
CALL    RNDBAL          ; GET RANDOM X, Y CO-ORD
LD      HL, BALLS      ; ONE OFF NUMBER OF BALLS
DEC     (HL)
RET

; RANDOM X-POSN FOR BALL

```



```

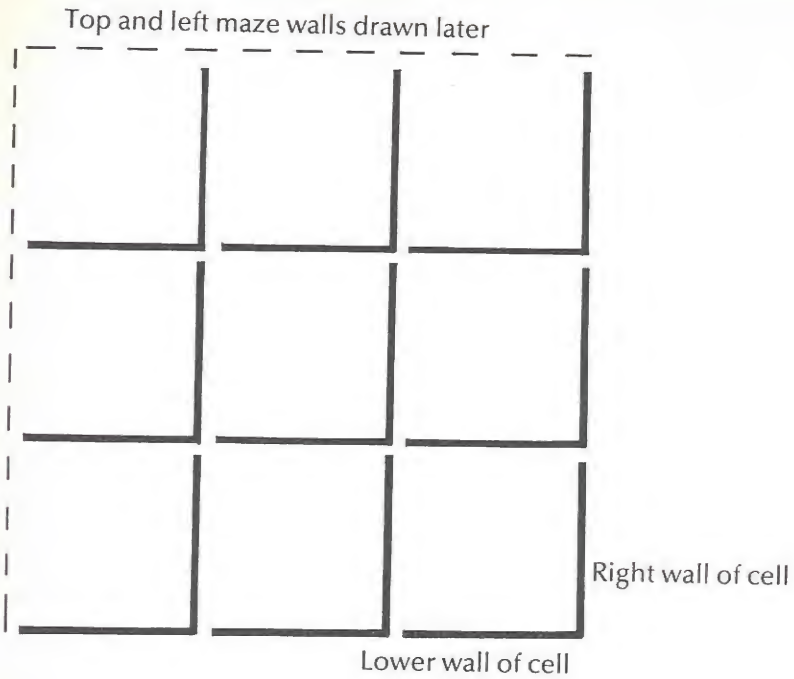
RDNBAL:                                ;GET FRAMES
      LD      A, (23672)
      ;LSB OF FRAMES
      SRL    A
      AND    0FH                        ;0-15
      ADD    A, 5                        ;5-20
      LD     (BALLX), A                  ;SAVE RANDOM X CO-ORD
                                          ;FOR BALL
      LD     A, 6                        ;INITIALIZE Y CO-ORD
      ;SET Y -POS
      LD     (BALLX+1), A                ;AND SAVE
      LD     HL, (BALLX)                 ;GET X, Y CO-ORD
      LD     (TPBLYX), HL                ;AND SAVE IN TEMP X, Y
      LD     B, 50
DLOOP:
      EI
      HALT
      DI
      DJNZ   DLOOP                       ;WAIT FOR A WHILE
      RET
      END

```

MAZE GENERATOR

This program generates random mazes consisting of a 32 by 22 grid of cells. It can, however, be easily adapted to produce mazes of any desired width and height. The algorithm used generates mazes where there is only one route from one cell to another. This routine could be used in games to produce a maze for an adventure game program such as the famous 'Hall of the Things' by Crystal Software.

A maze is constructed of cells and each of these cells are surrounded by up to four walls. They can be regarded as having just two walls as the other surrounding cells provide the other two. Diagram 1 shows how the wall would be made up on a 3 by 3 maze. Notice that there are no walls on the top and left sides of the maze. These we can draw later, after we have constructed the rest of the maze.



Hexadecimal Listing

7D00	C3	40	7D	0F	16	10	01	00
7D08	00	01	01	04	00	00	00	00
7D10	00	00	00	00	00	00	00	00
7D18	3C	7E	FF	FF	FF	FF	7E	3C
7D20	3F	7F	FF	FF	FF	FF	7F	3F
7D28	FF	FF	FF	FF	FF	FF	FF	FF
7D30	FC	FE	FF	FF	FF	FF	FE	FC
7D38	FF	81	81	81	81	81	81	FF
7D40	3E	02	CD	01	16	AF	D3	FE
7D48	21	00	00	22	0C	7D	3E	05
7D50	32	0E	7D	3E	04	32	0B	7D
7D58	3E	60	32	0F	7D	CD	1A	7E
7D60	CD	5F	7F	21	10	16	22	03
7D68	7D	CD	D9	7D	3A	0E	7D	A7
7D70	28	0B	CD	83	7D	CD	B7	7E
7D78	FB	76	F3	18	EF	ED	4B	0C

7D80	7D	FB	C9	CD	D9	7D	3E	FE
7D88	0B	FE	E6	1F	FE	1F	C8	CD
7D90	C5	7D	CB	4F	CC	B7	7D	CB
7D98	57	CC	A8	7D	CB	47	20	04
7DA0	CB	C7	18	EE	CD	D9	7D	C9
7DAB	F5	2A	03	7D	3E	1D	BD	28
7DB0	04	2C	22	03	7D	F1	C9	F5
7DB8	2A	03	7D	7D	A7	28	04	2D
7DC0	22	03	7D	F1	C9	F5	2A	03
7DC8	7D	01	38	03	C5	E5	AF	CD
7DD0	ED	7D	E1	2C	C1	10	F5	F1
7DD8	C9	2A	03	7D	01	39	03	3E
7DE0	02	C5	E5	CD	ED	7D	3C	E1
7DE8	2C	C1	10	F5	C9	F5	C5	E5
7DF0	C5	F5	E5	7C	CD	9E	0E	D1
7DF8	16	00	19	EE	F1	01	10	7D
7E00	26	00	6F	29	29	29	09	06
7E08	08	7E	12	23	14	10	FA	EB
7E10	CD	88	0E	C1	79	12	E1	C1
7E18	F1	C9	CD	6B	0D	01	20	20
7E20	3E	05	21	00	03	CD	37	7E
7E28	01	18	20	21	00	04	CD	37
7E30	7E	01	30	20	21	00	05	CD
7E38	ED	7D	2C	10	FA	C9	7D	6C
7E40	26	00	29	29	29	29	29	06
7E48	00	4F	09	01	00	58	09	7E
7E50	C9	01	00	00	FE	30	20	0A
7E58	3E	FF	32	0A	7D	01	02	00
7E60	18	20	FE	18	20	0A	3E	FF
7E68	32	0A	7D	01	05	00	18	12
7E70	FE	20	20	05	01	0A	00	18
7E78	09	11	40	00	21	66	06	CD
7E80	B5	03	2A	0C	7D	09	22	0C
7E88	7D	21	0F	7D	35	20	1A	36
7E90	60	3A	0B	7D	A7	28	0C	3D
7E98	32	0B	7D	3A	0E	7D	C6	02
7EA0	32	0E	7D	CD	5F	7F	CD	1A
7EA8	7E	11	08	00	21	66	06	CD
7EB0	B5	03	3E	00	D3	FE	C9	2A
7EB8	05	7D	3A	0A	7D	84	67	3A

7EC0	09	7D	85	6F	E5	CD	3E	7E
7EC8	E1	FE	39	20	26	3A	0A	7D
7ED0	ED	44	32	0A	7D	3A	09	7D
7ED8	ED	44	85	6F	CD	3E	7E	FE
7EE0	38	20	D4	3A	05	7D	A7	28
7EE8	CE	3A	09	7D	ED	44	32	09
7EF0	7D	18	C4	2A	05	7D	3A	09
7EF8	7D	85	6F	32	07	7D	A7	28
7F00	04	FE	1F	38	08	3A	09	7D
7F08	ED	44	32	09	7D	3A	0A	7D
7F10	84	67	32	08	7D	A7	28	0E
7F18	FE	17	30	31	CD	3E	7E	FE
7F20	38	28	0B	CD	51	7E	3A	0A
7F28	7D	ED	44	32	0A	7D	3E	00
7F30	0E	38	2A	05	7D	CD	ED	7D
7F38	3E	01	2A	07	7D	22	05	7D
7F40	CD	ED	7D	3A	0B	7D	FB	76
7F48	F3	3D	20	FA	C9	3E	00	0E
7F50	38	2A	05	7D	CD	ED	7D	CD
7F58	5F	7F	21	0E	7D	35	C9	3A
7F60	78	5C	CB	3F	E6	0F	C6	05
7F68	32	05	7D	3E	06	32	06	7D
7F70	2A	05	7D	22	07	7D	06	32
7F78	FB	76	F3	10	FB	C9	A7	28

The theory behind the maze generator is to walk randomly round the maze knocking down walls as we proceed. To begin we walk a set number of steps around the maze knocking down the walls if we meet any obstructions. On our second walk we start off in a cell that we have not previously entered. We again walk randomly around the maze knocking walls down. We keep on walking till we arrive at a cell which we had visited before on 'other walks'. When we arrive at such a cell we have then completed a path from one random walk to another on a different random walk. This process is repeated on all the untouched cells until we have proceeded through all the maze.

To implement this algorithm in machine code we represent our maze by having two arrays, the size of which are the size of the number of cells in the maze. One is called *BUILD*, the other *MAZE*. The array *BUILD* holds the path numbers and route which we 'walk' along while the array *MAZE* holds the 'wall' patterns. A wall pattern shows the structure of the two walls in a cell. The array *MAZE* is initialized

with the two walls intact. This is represented by the two first bits of its number being set high (i.e. the number three). Knocking down the walls is represented by re-setting a particular bit. If bit 0 of the number represents the bottom wall and bit 1 represents the right hand side wall then we can see the process if we knock down a wall. Going downwards we reset bit 0 of the cell we are in. If we knock down a wall going up we reset bit 0 of the cell above, the cell we are entering. Going right we re-set bit 1 of the cell we are in, going left we re-set bit 1 of the adjacent cell.

One point we have to look out for is that we do not 'back track' on a particular walk we are doing. We do this by giving each walk a path number and if we do happen to back track on our original path then we do not bother to knock down any walls. Using this method we guarantee our maze does not have any gaping holes and that it is singular in nature.

The program comes in two parts, one BASIC and one machine code. The machine code routine generates a random maze. The BASIC program draws the top and left hand side of the wall to complete the maze. When you use the generator in a game the unused bits in the array MAZE can be used to represent up to 63 objects such as axes, torches, wands or nasties! The second array is unused once the maze is generated so it could be used to store other variables or data in the game. The maze takes about two seconds to generate, very slow by machine code standards. Perhaps you could set yourself the task to make it faster. One way of improving the speed for 48K Spectrum owners would be to place the routine higher up in the memory map above the address 32768. Moving the code here would stop the Z80 CPU 'waiting' for the Spectrum's ULA to update the screen.

BASIC Listing

```

1 CLEAR 29000
10 CLS:LET SC=USR 32000
20 CLS:PRINT AT 10,10;"SCORE =
";SC
30 FOR X=1 TO 200:NEXT X
40 PRINT #0;"PRESS A KEY TO ST
ART"
50 PAUSE 0
60 GO TO 10

```

Assembler Listing

```

                ORG     30000
                JP     START
XPOS:          DB     0
YPOS:          DB     0
PATH:          DB     0

OPENCH EQU     1601H
UDG     EQU     23675                ;UDG ADDRESS

INTMAZE:
    LD     A, 2                ;OPEN SCREEN CHANNEL
    CALL  OPENCH
    LD     HL, NOUGHT        ;SET USER DEF GRAPHICS
                                ;TO OURS

    LD     (UDG), HL
    CALL  RANDI                ;INITIALIZE RANDOM
                                ;NUMBER GENERATOR

    LD     HL, MAZE          ;RE-BUILD THE MAZE
    LD     DE, MAZE+1
    LD     BC, 22*32        ;OF 22 BY 32
    LD     (HL), 3          ;WITH WALLS
    LDIR

; HL POINTS TO BUILD
; DE POINTS TO BUILD+1

    LD     BC, 22*32        ;CLEAR THE ARRAY
    LD     (HL), 0          ;BUILD WITH 0
    LDIR
    XOR    A                ;SET THE START X CO-ORD
    LD     (XPOS), A
    LD     (YPOS), A        ;SET THE START Y CO-ORD
    INC   A
    LD     (PATH), A        ;SET THE STARTING
                                ;PATH NUMBER

    RET

                                ;FINISH INITIALIZING

RAND0:         DB     0                ;RANDOM VAR 0
RAND1:         DB     0                ;RANDOM VAR 1

```

```

RAND2: DB      0      ;RANDOM VAR 2
RAND3: DB      0      ;RANDOM VAR 3

RAND:                                     ;GENERATE RANDOM NUMBER
                                           ;BETWEEN
                                           ;0 AND 255

LD      A,(RAND1)      ;GET RANDOM SEED
RRCA                                     ;A MOD 8 * 32
RRCA
RRCA
PUSH    BC              ;SAVE BC PAIR
PUSH    AF              ;SAVE AF PAIR
LD      A,(RAND2)      ;GET SECOND RANDOM VARIABLE
LD      B,A             ;AND PLACE IN B REGISTER
LD      A,(RAND3)      ;GET THIRD RANDOM VARIABLE
LD      C,A            ;AND PLACE IN C REGISTER
POP     AF              ;RESTORE AF
ADD     A,B             ;(RAND1)MOD 8 *32 + (RAND2)
ADD     A,C             ;+(RAND3)
RLCA                                       ;ALIGN BITS
RLCA                                       ;AND SAVE NEW RANDOM
                                           ;VARIABLES

LD      (RAND0),A
LD      A,B
LD      (RAND1),A
LD      A,C
LD      (RAND2),A
LD      A,(RAND0)
LD      (RAND3),A
POP     BC              ;RESTORE BC PAIR
RET

RANDI:
LD      A,0             ;SET UP RANDOM VARIABLES
LD      (RAND0),A
LD      A,173
LD      (RAND1),A
LD      A,206
LD      (RAND2),A
LD      A,R             ;ENSURE SOME RANDOMNESS
LD      (RAND3),A
RET

```

```

LENW    EQU    255D           ;LARGEST WALK

WALK:
LD      B, LENW

KEW:
                                ;KEEP WALKING

CALL    RANDW
LD      A, (HL)               ;GET CONTENTS OF
                                ;NEW POSTION
AND     A
JR      Z, PUTIN              ;TEST FOR NEW LOCATION
CP      C                     ;ZERO SO MARK PATH!
JR      Z, PUTIN              ;GOING BACK ON PATH?
CP      C                     ;YES MARK IT!
JR      Z, PUTIN              ;HAVE REACHED A VALUE LOWER
RET

PUTIN:
LD      A, (PATH)             ;GET PATH NUMBER
LD      (HL), A               ;AND PLACE IN BUILD
DJNZ    KEW                   ;DO THIS FOR LENW MAXIMUM
LD      A, (PATH)             ;ONLY DO LENW FOR PATH 1
CP      1
JR      NZ, WALK

RET

RANDW:
CALL    RAND                   ;GET RANDOM NUMBER
AND     3                     ;MASK OFF FOR
                                ;NUMBERS 0 TO 3

AND     A

JR      Z, NORTH              ;IF ZERO GO NORTH
CP      1
JR      Z, SOUTH              ;IF 1 GO SOUTH
CP      2
JR      Z, WEST               ;IF 2 GO WEST

;GO EAST

```


EAST:

```

LD      A, (XPOS)      ;GET X CO-ORD
CP      31              ;TEST TO SEE IF WE ARE ON
JR      Z, RANDW       ;THE RIGHT HAND SIDE
                          ;IF SO GO AGAIN
INC     A              ;ELSE INCREASE
                          ;X CO-ORD BY 1
LD      (XPOS), A      ;AND SAVE
LD      C, (HL)        ;OLD VALUE IN C REGISTER
INC     HL             ;NEW POSTION
LD      A, (HL)        ;NEW VALUE
CP      C              ;ARE THEY EQUAL?
RET     Z              ;DON'T BACKTRACK!

DEC     HL             ;GET OLD POSTION
CALL    RES1           ;RESET BIT 1 OF OLD CELL
INC     HL             ;POINT TO NEW CELL
RET

```

WEST:

```

LD      A, (XPOS)      ;GET X CO-ORD
AND     A              ;TEST FOR LEFT HAND SIDE
JR      Z, RANDW       ;PICK ANOTHER DIRECTION

DEC     A              ;GO LEFT

LD      (XPOS), A      ;SAVE X CO-ORD
LD      C, (HL)        ;OLD VALUE

DEC     HL             ;GO LEFT
LD      A, (HL)        ;GET NEW VALUE
CP      C              ;COMPARE WITH OLD VALUE
RET     Z              ;NO BACKTRACKING

```



```

INC      A                ; GOING DOWN
LD       (YPOS), A       ; SAVE Y CO-ORD
LD       C, (HL)        ; OLD VALUE.

LD       DE, 32          ; OFFSET FOR GOING DOWN
ADD      HL, DE
LD       A, (HL)        ; GET NEW PATH NUMBER
CP       C              ; COMPARE WITH OLD PATH
RET      Z              ; DON'T BACKTRACK
AND      A              ; CLEAR CARRY
SBC     HL, DE          ; NORMAL SUBTRACTION
CALL    RES0           ; KNOCK DOWN BOTTOM WALL
LD       DE, 32        ; GET HL BACK
ADD      HL, DE
RET

PRINTC:
PUSH    HL              ; SAVE REGISTERS
PUSH    BC

RST     010H           ; PRINT A REGISTER
                          ; TO CURRENT CHANNEL
                          ; RESTORE REGISTERS

POP     BC
POP     HL
RET

DISPLAY:
LD      HL, MAZE       ; DISPLAY MAZE TO
LD      A, 22         ; SCREEN. LOAD A WITH

LINE:   PUSH    AF     ; NUMBER OF LINES DOWN
                          ; SAVE LINE COUNT

DRAW:   LD      B, 32  ; GET CHARACTER COUNT
                          ; GET MAZE VALUE
LD      A, (HL)
ADD     A, 144        ; ADD BASE OF UDG
CALL    PRINTC       ; PRINT CHARACTER
INC     HL           ; NEXT MAZE CELL
DJNZ   DRAWS        ; REPEAT 32 TIMES

```

```

POP      AF          ;RESTORE LINE NUMBER
DEC      A          ;ONE OFF LINE NUMBER
JR       NZ,LINE    ;REPEAT 22 TIME
RET

```

BUILD:

```

LD       HL, BUILD  ;POINT TO ARRAY BUILD
XOR      A          ;INITIALIZE THE X CO-ORD
LD       (XPOS), A  ;AND Y CO-ORD
LD       (YPOS), A

```

FIND:

```

LD       A, (HL)    ;GET PATH NUMBER
AND      A          ;TEST FOR ZERO
JR       Z, SKIP    ;START WALKING ON ZERO
INC      HL         ;NEXT ONE ACROSS
LD       A, (XPOS)  ;UPDATE X CO-ORD
INC      A
LD       (XPOS), A
CP       32         ;HAVE WE GONE RIGHT ACROSS?
JR       NZ, FIND   ;NO, SO CARRY ON LOOKING
XOR      A          ;YES... RESET X CO-ORD
LD       (XPOS), A
LD       A, (YPOS)  ;GO ONE DOWN
INC      A
LD       (YPOS), A
CP       22         ;HAVE WE GONE ALL
JR       NZ, FIND   ;THE WAY DOWN?. NO THEN KEEP
RET      ;LOOKING.. ELSE RETURN

```

SKIP:

```

LD       A, (PATH)  ;GET PATH NUMBER
LD       (HL), A    ;PLACE AT NEW CELL
CALL     WALK       ;DO A RANDOM WALK
LD       A, (PATH)  ;UPDATE....
INC      A
LD       (PATH), A  ;NEW PATH NUMBER
JR       BUILD      ;CARRY ON BUILDING

```

```

START:
      CALL    INTMAZE      ;CLEAR MAZE
      CALL    BUILDM      ;BUILD MAZE
      CALL    DISPLAY     ;DISPLAY MAZE
      RET

```

```

;USER DEFINE CHARACTER SET FOR MAZE

```

```

NOUGHT:
      DB      0,0,0,0,0,0,0,0
      DB      0,0,0,0,0,0,0,255
      DB      1,1,1,1,1,1,1,1
      DB      1,1,1,1,1,1,1,255

```

```

MAZE:  DS      32*22
BUILD:  DS      32*22

```

```

      END

```

Hexadecimal Listing

7530	C3	8D	76	00	00	00	3E	02
7538	CD	01	16	21	97	76	22	7B
7540	5C	CD	90	75	21	E7	76	11
7548	B8	76	01	C0	02	36	03	ED
7550	E0	01	C0	02	36	00	ED	E0
7558	AF	32	33	75	32	34	75	3C
7560	32	35	75	C9	00	00	00	00
7568	3A	65	75	0F	0F	0F	C5	F5
7570	3A	66	75	47	3A	67	75	4F
7578	F1	80	81	07	07	32	64	75
7580	78	32	65	75	79	32	66	75
7588	3A	64	75	32	67	75	C1	C9
7590	3E	00	32	64	75	3E	AD	32
7598	65	75	3E	CE	32	66	75	ED
75A0	5F	32	67	75	C9	06	FF	CD
75A8	C0	75	7E	A7	28	04	B9	28


```

76F0  00  00  00  00  00  00  00  00
76F8  00  00  00  00  00  00  00  00
7700  00  00  00  00  00  00  00  00
7708  00  00  00  00  00  00  00  00
7710  00  00  00  00  00  00  00  00
7718  00  00  00  00  00  00  00  00
7720  00  00  00  00  00  00  00  00
7728  00  00  00  00  00  00  00  00

```

LARGE PRINT

I wrote this routine to enhance my own BASIC programs. The routine PRINTS characters on the screen twice the width of normal characters. I have 'patched' part of the BASIC operating system so that the large characters can be PRINTED from BASIC and will accept all the control characters, such as INK, PAPER, AT, TAB, etc. To enable the large PRINT facility we first call the routine at address 30000. This gives the Spectrum an additional channel, channel number 5. Then, to PRINT large characters to the screen we simply use the BASIC syntax:

```
PRINT#5;"STRING"
```

Here's a sample BASIC program which demonstrates how the routine can be used:

```

10 CLS : RANDOMIZE USR 30000
20 PRINT "This program demonstr
rates"
30 PRINT "How to get ";: PRINT
#5;"Large";: PRINT " letters"
40 PRINT #5;"    on the screen
"
50 PRINT
60 PRINT "It can cope with con
trol codes"
70 PRINT #5;AT 5,5;"such as AT
"
75 PRINT
80 PRINT #5; INK 5; PAPER 2;"a
nd colours"
90 PRINT #5; INVERSE 1;TAB 7;"
inverse"
100 PRINT #5; FLASH 1;" as well
as flashing"

```

Assembler Listing

```

                ORG      30000

CURCHL EQU      23633
REPORTJ EQU     15C4H      ;INVALID I/O DEVICE
STRMS EQU       23568D    ;STREAMS
STREAM5 EQU     STRMS+6+5*2 ;OPEN CHANNEL 5

CHARS EQU       23606D    ;CHARS CHARACTER ADDRESS
UDG EQU         23675D    ;UDG ADDRESS
CHANS EQU       23631D    ;CHANNEL ADDRESS
CHANINF EQU     STREAM5+2 ;CHANNEL 5

POCHANGE EQU    0A80H
TVDATA EQU     23566
TVDATL EQU     TVDATA
TVDATAH EQU    TVDATA+1
POCONT EQU     0A87H

INITP:
                ;SET UP PRINT#4 COMMAND
                ;MOVE CHANNEL INFORMATION
                LD      HL, CHANIND
                LD      DE, CHANINF
                LD      BC, 5
                LDIR
                LD      HL, CHANINF      ;FIND DISTANCE BETWEEN CHAN
                LD      DE, (CHANS)
                AND     A
                SBC    HL, DE
                INC    HL
                LD      (STREAM5), HL    ;SET UP STREAMS
                RET

CHANIND:
                DEFW   PRINTD      ;PRINT OUT ROUTINE
                DEFW   REPORTJ    ;INPUT ROUTINE.
                DEFB   'D'

```


PRINTD:

; WHEN BASIC CALLS THIS ROUTINE THE
; A REG CONTAINS CHAR NUMBER.

```

CP      20H          ;TEST TO SEE IF PRINTABLE
JP      NC, CAR      ;PRINT THE CHARACTER
CALL    0B03H        ;GET CURRENT PRINT POSITION
CP      06           ;PRINT '?'
                          ;FOR CODES 00- 05 HEX

JP      C, 0A69H
CP      18H          ;AND 18H TO 1FH

JP      NC, 0A69H
CP      16
JP      C, 09F4H+16D ;GO TO ROMS TABLE.

LD      HL, ATTAB    ;ASSUME CONTROL CHAR
                          ;IS AT OR TAB

CP      22
JR      NC, RIGHT    ;YOU WERE RIGHT!
LD      HL, INKOVER

```

RIGHT:

```

PUSH    HL           ;RETURN ADDRESS IS PUSHED
JP      0B03H        ;FETCH CURRENT CHARACTER

```

POTV2D:

```

LD      DE, POCONTD  ;SAVE FIRST OPERAND
                          ;IN TVDATH

LD      (TVDATH), A
JP      POCHANGE     ;CHANGE ADDRESS OF
                          ;CURRENT CHANNEL

```

ATTAB:

```

LD      DE, POTV2D   ;NEXT TIME ROUND GOTO POTVD
JR      POTV1D       ;SAVE CHARACTER CODE
                          ;IN TVDATAL

```

INKOVER:

```

LD      DE, POCONTD  ;NEXT TIME POCONTD
POTV1D: LD (TVDATL), A ;SAVE CONTROL CODE
JP      POCHANGE     ;CHANGE OUTPUT ADDRESS

```



```

CALL    09F4H
LD      A, 144          ;NOW PRINT USER DEFINED
                        ;GRAPHICS

CALL    09F4H
POP     HL              ;RESTORE UDG
LD      (UDG), HL
RET

FETD:
CALL    NYBBLE         ;DO ONE NYBBLE
LD      L, H
;NOW GET NEXT NYBBLE
NYBBLE:
LD      B, 4           ;NUMBER OF BITS
NBIT:   RRCA

;TRY CHANGING ABOVE OPCODE TO RLCA!!! FOR A BIT OF FUN.

RR      C              ;MAKE TWICE AS FAT
SRA    C
DJNZ   NBIT           ;DO 8 TIMES
LD     H, C
RET

DUGD:   DS      8*2

END

```

Hexadecimal Listing

```

7530    21  4A  75  11  22  5C  01  05
7538    00  ED  B0  21  22  5C  ED  5B
7540    4F  5C  A7  ED  52  23  22  20
7548    5C  C9  4F  75  C4  15  44  FE
7550    20  D2  91  75  CD  03  0B  FE
7558    06  DA  69  0A  FE  18  D2  69
7560    0A  FE  10  DA  04  0A  21  7D
7568    75  FE  16  30  03  21  82  75

```

7570	E5	C3	03	0B	11	8B	75	32
7578	0F	5C	C3	80	0A	11	74	75
7580	18	03	11	8B	75	32	0E	5C
7588	C3	80	0A	11	4F	75	C3	8A
7590	0A	ED	5B	36	5C	26	00	6F
7598	29	29	29	19	EB	21	DE	75
75A0	E5	DD	E1	06	08	1A	C5	CD
75A8	CF	75	C1	DD	75	00	DD	74
75B0	08	DD	23	13	10	EF	2A	7B
75B8	5C	E5	21	DE	75	22	7B	5C
75C0	3E	91	CD	F4	09	3E	90	CD
75C8	F4	09	E1	22	7B	5C	C9	CD
75D0	D3	75	6C	06	04	0F	CB	19
75D8	CB	29	10	F9	61	C9	00	00
75E0	00	00	00	00	00	00	00	00
75E8	00	00	00	00	00	00	82	75
75F0	E5	C3	03	0B	11	8B	75	32
75F8	0F	5C	C3	80	0A	11	74	75

PIXEL SCROLL

This routine allows the user to scroll any portion of the screen to either left or right. It has a 'wrap-around' effect, and so could be most useful when writing arcade games with scrolling background scenery of mountains, high rise flats or the like. When calling the routine the HL register pair must point to the screen address of the position from which you wish to scroll. The program below is a demonstration program showing how the routine can be used from BASIC:

```

5 FOR X=1 TO 32*22:PRINT CHR#
143;:NEXT X
10 PRINT AT 0,0;"THIS SCROLL W
ILL GO LEFT"
20 PRINT "WITH THIS LINE!!"
30 PRINT AT 8,0;"THIS SCROLL W
ILL GO RIGHT"
40 PRINT "ALONG WITH THIS LINE
!!"
50 RANDOMIZE USR 32000
60 GO TO 50

```

Here are the listings for the scroll routine:

Assembler Listing

```

        ORG      32000D

NLINE$ EQU      16          ;NUMBER OF LINES TO SCROLL
NBYTES EQU      32          ;NUMBER OF BYTES TO SCROLL
ADD     EQU      16384       ;SCREEN ADDRESS
ADD2    EQU      16384+256*8 ;SCREEN ADDRESS2

        JP      TEST        ;TEST THE SCROLL

;*****
;These two routines SLEFT SCROLL LEFT
; and SRIGHT SCROLL RIGHT scroll the screen
; left and right respectively.They use the rout-
; -ine INCY which finds address of corresponding
; pixel line addresses.
; On entry to the routine HL points to the top
; left hand side of the portion of the screen
; to be scrolled.
; The other values which the program will give are
; NBYTES number of bytes to scroll ie width
; NLINE$ number of lines to scroll
; Both routines have a wrap-around effect.

SLEFT:
        LD      HL,ADD+NBYTES-1
                                ;POINT TO RIGHT HAND SIDE
        LD      C,NLINE$       ;NUMBER OF LINES TO SCROLL

LINE:   PUSH    HL              ;SAVE SCREEN ADDRESS
        LD      B,NBYTES       ;NUMBER OF BYTES TO SCROLL
                                ;ACROSS

CHARX:
        RL      (HL)           ;SCROLL LEFT THROUGH
        DEC    HL              ;CARRY
        DJNZ   CHARX          ;REPEAT NBYTES TIMES
        POP    HL              ;RESTORE RIGHT HAND
                                ;SIDE ADDRESS
        LD     A,0             ;SET A TO ZERO

```

```

ADC    A, A           ;PLACE CARRY IN BIT 0 OF
OR     (HL)          ;RIGHT HAND SIDE
LD     (HL), A
CALL  INCY           ;GET ADDRESS OF NEXT PIXEL
                           ;LINE
DEC    C             ;DOWN, ONE LESS LINE
JR     NZ, LINE      ;REPEAT UNTIL DONE ALL LINES
RET

SRIGHT:
LD     HL, ADD2      ;POINT TO LEFT HAND
                           ;OF SCREEN
LD     C, NLINES     ;NUMBER OF LINES TO SCROLL
LINER: PUSH HL       ;SAVE LEFT HAND
                           ;SIDE ADDRESS
LD     B, NBYTES     ;NUMBER OF BYTES TO SCROLL

CHARR:
RR     (HL)          ;SCROLL RIGHT THROUGH CARRY
INC    HL            ;TO THE RIGHT
DJNZ  CHARR          ;REPEAT UNTIL DONE
                           ;NBYTES TIMES
POP   HL             ;RESTORE LEFT HAND SIDE
LD     A, 0          ;ROTATE CARRY INTO LEFT
                           ;HAND SIDE

RRA
OR     (HL)
LD     (HL), A
CALL  INCY           ;NEXT PIXEL LINE DOWN
DEC    C             ;ONE LESS PIXEL LINE
JR     NZ, LINER    ;REPEAT UNTIL NO MORE LINES
RET

INCY:
                           ;NEXT PIXEL LINE DOWN
                           ;HL POINTS TO SCREEN
                           ;ADDRESS
INC    H             ;NEXT LINE DOWN
LD     A, H          ;TEST IF WITHIN CHARACTER
AND   7
RET   NZ             ;WITHIN CHAR SO RETURN

LD     A, L          ;NEXT CHARACTER DOWN

```

```

ADD    A, 20H
LD     L, A
RET    C                ; DO NOT ADJUST SECTOR SINCE
                        ; WE HAVE GONE OVER

```

```

LD     A, H            ; WITHIN SECTOR SO
SUB    B
LD     H, A            ; RE-ADJUST
RET

```

TEST:

```

CALL   SLEFT          ; SCROLL LEFT
CALL   SRIGHT         ; SCROLL RIGHT
RET

```

END

Hexadecimal Listing

7D00	C3	46	7D	21	1F	40	0E	10
7D08	E5	06	20	CB	16	2B	10	FB
7D10	E1	3E	00	8F	B6	77	CD	37
7D18	7D	0D	20	EC	C9	21	00	48
7D20	0E	10	E5	06	20	CB	1E	23
7D28	10	FB	E1	3E	00	1F	B6	77
7D30	CD	37	7D	0D	20	EC	C9	24
7D38	7C	E6	07	C0	7D	C6	20	6F
7D40	D8	7C	D6	08	67	C9	CD	03
7D48	7D	CD	1D	7D	C9	00	00	00
7D50	00	00	00	00	00	00	00	00
7D58	00	00	00	00	00	00	00	00
7D60	00	00	00	00	00	00	00	00
7D68	00	00	00	00	00	00	00	00
7D70	00	00	00	00	00	00	00	00
7D78	00	00	00	00	00	00	00	00

AUTO LINE NUMBER

This routine produces line numbers automatically when the user is typing in a BASIC program. Like the Clock and Trace programs given

in Chapter 11, this uses interrupts. The routine LOADS the ASCII line number into the system variable LAST-K every 1/50th of a second. This causes the line number to be placed in the edit area and lower screen of the Spectrum. To enable the auto line facility, key in the instruction RAND USR 32333. To turn off the auto line first DELETE the line number currently being edited and then enter RAND USR 32330.

Don't forget to CLEAR memory to keep the machine code safe. CLEAR 32329 is suitable for this.

Assembler Listing

```

                ORG      32330D

ECHOE EQU      23682      ;COL NUMBER AND ROW NUMBER
LASTK EQU      23560      ;LAST KEY TO BE PRESSED
FLAGS EQU      23611      ;KEYBOARD FLAGS
EPCC EQU      23625      ;CURRENT LINE NUMBER
PPC EQU      23621      ;CURRENT LINE EXECUTING

DISINT:
    IM      1      ;ENABLE DEFAULT INTERRUPTS
    RET

ENABLE:
    XOR     A      ;SET A TO ZERO
    LD     (STATE),A ;NOT OUTPUTTING ASCII CHARS
    LD     A,28H   ;SET I REG TO PAGE 28H
    LD     I,A
    IM     2      ;AND ENABLE INTERRUPT MODE 2
    EI
    RET

                ORG      7E5CH      ;START OF INTERRUPT ROUTINE

                CALL     38H      ;SCAN KEYBOARD FIRST
                DI      ;DISABLE INTERRUPTS FIRST
                PUSH AF      ;SAVE REGISTERS
                PUSH BC
                PUSH DE
                PUSH HL
                PUSH IX

```



```

LD      A, (STATE)      ;GET STATE
AND     A                ;TEST FOR ZERO
JR      NZ, DML         ;ALL READY DOING A LINE

LD      A, (PPC+1)      ;ARE WE EXECUTING A LINE?
CP      255
JR      NZ, BYE         ;YES SO EXIT FROM ROUTINE

LD      A, (ECHOE)
CP      20H
JR      NZ, BYE

LD      A, (ECHOE+1)
CP      17H
JR      NZ, BYE         ;ARE WE AT BOTTOM OF SCREEN
;ARRIVE HERE IF WE ARE AT THE BOTTOM OF THE SCREEN

LD      A, (LASTK)      ;GET LASTK
CP      0DH             ;IF NOT RETURN
JR      NZ, BYE         ;THEN EXIT FROM ROUTINE

FIRST: LD      A, 4      ;SET UP NO OF CHARS
                ;TO PRINT
LD      (STATE), A
LD      HL, DECTL      ;START OF TABLE
LD      (CDATA), HL

LD      HL, (EPCC)     ;GET CURRENT LINE NO
LD      DE, 000AH      ;GET STEP NUMBER
ADD     HL, DE         ;GET NEXT LINE NUMBER
LD      (LINE), HL     ;AND SAVE

DML:   LD      A, (STATE) ;GET STATE
DEC     A              ;ONE LESS CHAR
LD      (STATE), A     ;TO PRINT
LD      HL, (LINE)     ;GET LINE NUMBER
CALL   CONV            ;OUTPUT ASCII TO LASTK

```

```

BYE:                                ;RESTORE REGS
      POP IX
      POP HL
      POP DE
      POP BC
      POP AF
      EI
      RET

DECTL: DEFW    1000D                ;START OF TABLE
DECTX: DEFW    100D
      DEFW    10D
      DEFW    1D

CONV:  LD      IX,(CDATA)           ;GET CURRENT TABLE POINTER
NDIGIT: LD     C,(IX+0)             ;GET LOW BYTE OF MULTIPLES
                                           ;OF TENS
      LD      B,(IX+1)             ;GET HIGH BYTE OF MULTIPLE
      LD      A,'0'-1              ;SET UP A REG WITH 30 HEX
      AND     A                     ;RESET CARRY
FIDIG:  INC     A                    ;ADD 1 TO A REGISTER
      SBC     HL,BC                 ;UNTIL WE GET A CARRY
      JR      NC,FIDIG
      ADD     HL,BC                 ;CORRECT NUMBER IN HL
      LD      (LINE),HL            ;AND SAVE IT!
      CALL   OUTP2                 ;OUTPUT ASCII CHAR IN A
                                           ;REGISTER
      INC     IX                    ;POINT TO NEXT MULTIPLE
      INC     IX
      PUSH    IX                    ;TRANSFER IX TO HL
                                           ;REGISTER PAIR
      POP     HL
      LD      (CDATA),HL           ;AND SAVE
      RET

OUTP2: LD      (LASTK),A            ;PLACE CHAR IN LASTK
      LD      HL,FLAGS             ;SIGNIFY WE ....
      SET     5,(HL)               ;PRESSED A KEY
      RET

```

```

CDATA:  DEFW  0           ;CURRENT LINE DATA
STATE:  DB    0           ;NO OF CHARS TO PRINT
LINE:   DEFW  0           ;LINE NUMBER

```

END

Hexadecimal Listing

```

7E4A  ED  56  C9  AF  32  E9  7E  3E
7E52  28  ED  47  ED  5E  FB  C9  00
7E5A  00  00  CD  38  00  F3  F5  C5
7E62  D5  E5  DD  E5  3A  E9  7E  A7
7E6A  20  31  3A  46  5C  FE  FF  20
7E72  37  3A  82  5C  FE  20  20  30
7E7A  3A  83  5C  FE  17  20  29  3A
7E82  08  5C  FE  0D  20  22  3E  04

7E8A  32  E9  7E  21  B2  7E  22  E7
7E92  7E  2A  49  5C  11  0A  00  19
7E9A  22  EA  7E  3A  E9  7E  3D  32
7EA2  E9  7E  2A  EA  7E  CD  BA  7E
7EAA  DD  E1  E1  D1  C1  F1  FB  C9
7EB2  E8  03  64  00  0A  00  01  00
7EBA  DD  2A  E7  7E  DD  4E  00  DD
7EC2  46  01  3E  2F  A7  3C  ED  42

7ECA  30  FB  09  22  EA  7E  CD  DE
7ED2  7E  DD  23  DD  23  DD  E5  E1
7EDA  22  E7  7E  C9  32  08  5C  21
7EE2  3B  5C  CB  EE  C9  00  00  00
7EEA  00  00  3A  46  5C  FE  FF  20
7EF2  37  3A  82  5C  FE  20  20  30
7EFA  3A  83  5C  FE  17  20

```

SORT

Another program which can be used with BASIC, this sort routine which allows you to sort strings into alphabetical order. The routine, when called in BASIC searches for the dimensional array `AS`. It should be first set up with the number of objects to sort and the length of each string. If the string is not found or the length is too large then it will exit from the sort routine with an appropriate error message. When you wish to sort the string you simply call the machine code from BASIC by using the instruction `RAND USR 32000`. This will then sort

out the string in ascending order. The method used to sort out the strings is known as a 'Bubble Sort'. This method of sorting is not the most efficient. However, under one second to sort out 100 strings of 25 characters in length is not slow!

The BASIC listing below demonstrates how the machine code program is used:

```

5 LET sort=32000
10 DIM a$(100,25)
20 FOR p=1 to 100
30 FOR c=1 to 25
40 LET a$(p,c)=CHR$( (RND*26)+
65)
50 NEXT c
60 NEXT p
70 PRINT #0;"Press L to list,S
to sort"
80 LET k$=INKEY$: IF k$="" THE
N GO TO 80
90 IF k$="L" OR k$="1" THEN GO
SUB 120: GO TO 70
100 IF k$<>"s" AND k$<>"S" THEN
GO TO 80
110 CLS: PRINT "sorting":RANDOM
IZE USR sort: BEEP 1,1: GO SUB 1
20: STOP
120 FOR p=1 TO 100
130 PRINT a$(p)
140 NEXT p
150 RETURN

```

Here are the listings for the sort routine:

Assembler Listing

```

          ORG      32000
VARS      EQU      236270

START:
          LD       HL,(VARS)      ;SET HL TO POINT TO
                                   ;VARIABLE AREA
TEST:    LD       A,(HL)         ;GET 1ST BYTE OF VARIABLE
          CP       128           ;END OF VARS MARKER?
          JR       Z,NOTFOUND    ;FINISHED LOOKING AT VARS

```

```

CP      193          ; IS IT A$?
JP      Z, FOUND    ; YES FOUND IT!

AND     11100000B   ; MASK OFF TOP THREE BITS
CP      01100000B   ; SINGLE DATA
JR      Z, ADISIX   ; YES ADD 6
CP      11100000B   ; 'FOR NEXT' VARIABLE?
JR      Z, ADI19    ; YES ADD 19
CP      10100000B   ; VARIABLE NAME LARGER
                          ; THEN ONE LETTER?
JR      Z, SKIPC    ; YES THEN SKIP PASS
                          ; VARIABLE NAME

; 010 OR 110
INC     HL
LD      E, (HL)     ; GET LENGTH LOW
INC     HL
LD      D, (HL)     ; GET LENGTH HIGH
INC     HL
ADD     HL, DE      ; SKIP PASS VARIABLE
JP      TEST        ; TEST FOR NEXT VARIABLE

NOTFOUND:
RST     08
DB      01          ; VARIABLE NOT FOUND ERROR!

ERROR:
RST     08          ; SUBSCRIPT WRONG ERROR!
DB      02

ADI19:
LD      DE, 19     ; GO PASS VARIABLE
ADD     HL, DE
JP      TEST        ; TEST NEXT VARIABLE

SKIPC:
INC     HL          ; SKIP PASS VARIABLE.
BIT     7, (HL)    ; NAME
JR      Z, SKIPC   ; TILL BIT 7 IS SET

ADISIX:
LD      DE, 6      ; GO PASS VARIABLE

```

```

ADD    HL, DE
JP     TEST          ;TEST NEXT VARIABLE

```

FOUND:

```

INC    HL
INC    HL
INC    HL          ;POINT TO NUMBER OF DIMS
                ;MUST BE TWO OR LESS

```

```

LD     A, (HL)
CP     2
JR     NZ, ERROR   ;SHOULD BE TWO DIMENSIONS.

```

```

INC    HL          ;POINT TO NUMBER OF
                ;ELEMENTS
LD     B, (HL)     ;NUMBER OF ELEMENTS
INC    HL          ;GET HIGH BYTE!
LD     A, (HL)
AND    A
JR     NZ, ERROR   ;LARGER THEN 255 ELEMENTS
INC    HL
LD     C, (HL)     ;LENGTH OF STRINGS
INC    HL
LD     A, (HL)
AND    A
JR     NZ, ERROR   ;LARGER THEN 255 CHARACTERS
INC    HL
;HL NOW POINTS TO START OF STRING

LD     A, C        ;SAVE SIZE
LD     (SIZE), A

```

SORT:

```

;HL POINTS TO START OF STRING
;B CONTAINS NUMBER OF STRINGS
;C CONTAINS LENGTH OF STRING

```

```

NEXTS:  PUSH    BC           ; SAVE NUMBER AND LENGTH
        XOR     A           ; RESET SWAP FLAG
        LD      (FLAG), A
        PUSH   HL           ; SAVE ADDRESS OF FIRST
                               ; STRING

NEXTEL:  PUSH   HL           ; SAVE ADDRESS OF STRING
        LD     E, C         ; GET LENGTH OF STRING
        LD     D, 0         ; AND PLACE IN DE REGISTER
        ADD   HL, DE        ; POINT TO SECOND STRING
        EX    DE, HL        ; AND PLACE IN
                               ; THE DE REGISTER
        POP   HL           ; RESTORE ADDRESS OF STRING
        CALL  COMPARE        ; COMPARE THE TWO STRINGS
        CALL  C, SWAP        ; IN ASCENDING ORDER

        EX    DE, HL        ; HL NOW POINTS TO
                               ; NEXT STRING

        DJNZ  NEXTEL        ; REPEAT COMPARISION UNTIL
                               ; DONE UP TO CURRENT NUMBER
                               ; OF STRINGS
        POP   HL           ; GET ADDRESS OF
                               ; FIRST STRING
        POP   BC           ; RESTORE COUNTERS
        LD   A, (FLAG)     ; GET SWAP FLAG
        AND  A             ; TEST FOR ZERO
        RET  Z             ; NO SWAPS MADE SO SORTED
        DEC  B             ; ONE LESS TO SORT
        JR   NZ, NEXTS

COMPARE:  PUSH   HL           ; SAVE REGISTERS
        PUSH  DE
        PUSH  BC

COMPARS:  ; COMPARE STRINGS
          ; ONE POINTED BY THE HL PAIR
          ; AND ONE POINTED BY
          ; THE DE PAIR

```

```

        LD      A, (DE)      ;GET CHARACTER
        SUB    (HL)        ;COMPARE AGAINST THE
                           ;SAME ONE IN
                           ;THE SECOND STRING
        JR     NZ, BYEFC    ;NOT EQUAL EXIT FROM
                           ;COMPARISON
        INC    HL          ;POINT TO NEXT CHARACTER
        INC    DE          ;POINT TO NEXT CHARACTER
        DEC    C           ;REPEAT UNTIL COMPARED
                           ;ALL CHARACTERS
        JR     NZ, COMPARS

BYEFC:  POP    BC          ;RESTORE REGISTERS
        POP    DE
        POP    HL
        RET

SWAP:
        ;SWAP THE TWO STRINGS POINTED
        ;BY THE HL PAIR AND THE DE PAIR

        PUSH   BC          ;SAVE REGISTERS
        PUSH   DE
        PUSH   HL

        LD     A, (SIZE)   ;GET SIZE
        LD     C, A        ;PLACE IN THE LOW BYTE
                           ;OF THE COUNTER
        LD     B, 0        ;NOT LARGER THEN 255

        LD     DE, BUFF    ;DE POINTS TO THE BUFFER
        LDIR                    ;MOVE THE STRING FROM HL
                           ;TO THE BUFFER

        POP    DE          ;PUT ORIGINAL HL IN DE
        PUSH   DE

        LD     C, A        ;GET COUNTER
        LDIR                    ;MOVE TO SECOND STRING

        LD     HL, BUFF    ;POINT TO BUFFER
        LD     C, A        ;GET COUNT
        LDIR                    ;AND SWAP

```



```

LD      (FLAG), A      ; SIGNIFY A SWAP WAS MADE

POP     HL              ; RESTORE REGISTERS
POP     DE
POP     BC
RET

BUFF:   DS      255     ; BUFFER
FLAG:   DB      0      ; SWAP FLAG
SIZE:   DB      0      ; SIZE OF STRING
END

```

Hexadecimal Listing

```

7D00    2A  4B  5C  7E  FE  80  28  1C
7D08    FE  C1  CA  3B  7D  E6  E0  FE
7D10    60  28  21  FE  E0  28  11  FE
7D18    A0  28  14  23  5E  23  56  23
7D20    19  C3  03  7D  CF  01  CF  02
7D28    11  13  00  19  C3  03  7D  23
7D30    CB  7E  28  FB  11  06  00  19
7D38    C3  03  7D  23  23  23  7E  FE

7D40    02  20  E3  23  46  23  7E  A7
7D48    20  DC  23  4E  23  7E  A7  20
7D50    D5  23  79  32  A7  7E  C5  AF
7D58    32  A6  7E  E5  E5  59  16  00
7D60    19  EB  E1  CD  77  7D  DC  87
7D68    7D  EB  10  F0  E1  C1  3A  A6
7D70    7E  A7  C8  05  20  E0  C9  E5
7D78    D5  C5  1A  96  20  05  23  13

7D80    0D  20  F7  C1  D1  E1  C9  C5
7D88    D5  E5  3A  A7  7E  4F  06  00
7D90    11  A7  7D  ED  B0  D1  D5  4F
7D98    ED  B0  21  A7  7D  4F  ED  B0
7DA0    32  A6  7E  E1  D1  C1  C9  00
7DA8    00  00  00  00  00  00  00  00
7DB0    00  00  00  00  00  00  00  00
7DB8    00  00  00  00  00  00  00  00

```

RECURSION

This program is similar to the music routine given in Chapter 9 but is slightly more elaborate and complex. The tune I have given is the one I translated (from the Spectrum manual) from the section on the BEEP command. You can however write your own music. See the table given in Chapter nine. The routine is called by setting the IX register to point to the music data. The data represents the notes to be played and the duration. Each note and duration is represented by two bytes making a total of four. The first two bytes make up the frequency of the note and the second two the duration. The nice thing about this music routine is that it has the ability to play substrings of music. The routine scans first of all for the frequency in the table. If the low byte of the frequency is a one then this indicates that the following two bytes are the address of a substring to be played. The end of a string of music is indicated by having the byte 0. Substrings can be nested to many levels dependent on the RAM you have left. The whole principle behind this routine is that of recursion. It's a routine which calls itself, in the same way as BASIC subroutines can.

Assembler Listing

```

                ORG      32000D

BEEPER EQU      03B5H          ;ADDRESS OF BEEPER ROUTINE
LD      IX, FRERE          ;POINT TO MUSIC
CALL    PLAY              ;AND PLAY IT SAM!
RET

PLAY:
    PUSH    IX              ;SAVE STRING POSTION
    LD      L, (IX+0)        ;LOW PITCH
    LD      H, (IX+1)        ;HIGH PITCH
    LD      E, (IX+2)        ;LOW DURATION
    LD      D, (IX+3)        ;HIGH DURATION
    LD      A, L              ;LOOK AT LOW PITCH
    CP      01
    JR      Z, PLS          ;PLAY SUBSTRING
    JR      C, BYE          ;ZERO SO BYE
    CALL    BEEPER          ;PLAY NOTE
    POP     IX              ;GET STRING POSTION
    LD      DE, 4            ;NEXT NOTE AND DURATION
    ADD    IX, DE
    JR      PLAY            ;KEEP PLAYING SAM!

```

```

PLS:      INC      IX          ;POINT
          INC      IX          ;TO RETURN POSTION
          INC      IX
          LD       L, H
          ;ADJUST SUBSTRING ADDRESS
          LD       H, E
          POP      AF          ;GET RID OF OLD STRING
          ;ADDRESS
          PUSH     IX          ;PUT IN NEW STRING ADDRESS

          PUSH     HL          ;TRANFER SUBSTRING
          POP      IX          ;ADDRESS TO IX REGISTER
          CALL     PLAY        ;PLAY SUBSTRING
          POP      IX          ;RETURNED FROM PLAYING
          ;SUBSTRING
          JR       PLAY        ;KEEP PLAYING

BYE:      POP      IX
          RET

```

```

FRERE:   DB       01
          DEFW    FRERE1
          DB       01
          DEFW    FRERE1

```

```

FRERE1:  DB       01
          DEFW    TUNE1
          DB       01
          DEFW    TUNE1

```

```

          DB       01
          DEFW    TUNE2
          DB       01
          DEFW    TUNE2

```

```
DB      01
DEFW    TUNE3
DB      01
DEFW    TUNE3
```

```
DB      01
DEFW    TUNE4
DB      01
DEFW    TUNE4
```

TUNE1:

```
DEFW    66AH
DEFW    105H
DEFW    5B3H
DEFW    125H
DEFW    560H
DEFW    98H
DEFW    5B3H
DEFW    92H
DEFW    66AH
DEFW    105H
DEFW    00
```

TUNE2:

```
DEFW    560H
DEFW    137H
DEFW    4C6H
DEFW    15DH
DEFW    43DH
DEFW    188H
DEFW    00
```

TUNE3:

```

DEFW 43DH
DEFW 126H
DEFW 3FFH
DEFW 67H
DEFW 43DH
DEFW 0C4H
DEFW 4C6H
DEFW 0AEH
DEFW 560H
DEFW 98H
DEFW 5B3H
DEFW 92H
DEFW 66AH
DEFW 105H
DEFW 00

```

TUNE4:

```

DEFW 66AH
DEFW 105H
DEFW 89AH
DEFW 0C4H
DEFW 66AH
DEFW 20AH
DEFW 00
END

```

Hexadecimal Listing

```

7D00 DD 21 41 7D CD 08 7D C9
7D08 DD E5 DD 6E 00 DD 66 01
7D10 DD 5E 02 DD 56 03 7D FE
7D18 01 28 0E 38 21 CD B5 03
7D20 DD E1 11 04 00 DD 19 18
7D28 DF DD 23 DD 23 DD 23 6C
7D30 63 F1 DD E5 E5 DD E1 CD
7D38 08 7D DD E1 18 CA DD E1

```

7D40	C9	01	47	7D	01	47	7D	01
7D48	5F	7D	01	5F	7D	01	75	7D
7D50	01	75	7D	01	83	7D	01	83
7D58	7D	01	A1	7D	01	A1	7D	6A
7D60	06	05	01	B3	05	25	01	60
7D68	05	9B	00	B3	05	92	00	6A
7D70	06	05	01	00	00	60	05	37
7D78	01	C6	04	5D	01	3D	04	88
7D80	01	00	00	3D	04	26	01	FF
7D88	03	67	00	3D	04	C4	00	C6
7D90	04	AE	00	60	05	9B	00	B3
7D98	05	92	00	6A	06	05	01	00
7DA0	00	6A	06	05	01	9A	08	C4
7DA8	00	6A	06	0A	02	00	00	6C
7DE0	63	F1	DD	E5	E5	DD	E1	CD
7DE8	08	7D	DD	E1	18	CA	DD	E1
7DC0	C9	01	47	7D	01	47	7D	01
7DC8	5F	7D	01	5F	7D	01	75	7D
7DD0	01	75	7D	01	83	7D	01	83
7DD8	7D	01	A1	7D	01	A1	7D	6A
7DE0	06	05	01	B3	05	25	01	60
7DE8	05	9B	00	B3	05	92	00	6A
7DF0	06	05	01	00	00	60	05	37
7DF8	01	C6	04	5D	01	3D	04	88

Appendix 1

Z80 instructions listed by mnemonic

8E		142		ADC A,(HL)
DD 8E dd		221 142 dd		ADC A,(IX'd)
FD 8E dd		253 142 dd		ADC A,(IY'd)
8F		143		ADC A,A
88		136		ADC A,B
89		137		ADC A,C
8A		138		ADC A,D
8B		139		ADC A,E
8C		140		ADC A,H
8D		141		ADC A,L
CE XX		206 XX		ADC A,N
ED 4A		237 74		ADC HL,BC
ED 5A		237 90		ADC HL,DE
ED 6A		237 106		ADC HL,HL
ED 7A		237 122		ADC HL,SP
86		134		ADD A,(HL)
DD 86 dd		221 134 dd		ADD A,(IX'd)
FD 86 dd		253 134 dd		ADD A,(IY'd)
87		135		ADD A,A
80		128		ADD A,B
81		129		ADD A,C
82		130		ADD A,D
83		131		ADD A,E
84		132		ADD A,H
85		133		ADD A,L
C6 XX		198 XX		ADD A,N
09		9		ADD HL,BC
19		25		ADD HL,DE
29		41		ADD HL,HL
39		57		ADD HL,SP
DD 09		221 9		ADD IX,BC
DD 19		221 25		ADD IX,DE
DD 29		221 41		ADD IX,IX
DD 39		221 57		ADD IX,SP
FD 09		253 9		ADD IY,BC

FD 19		253 25		ADD IY,DE
FD 29		253 41		ADD IY,IY
FD 39		253 57		ADD IY,SP
A6		166		AND (HL)
DD A6 dd		221 166 dd		AND (IX'd)
FD A6 dd		253 166 dd		AND (IY'd)
A7		167		AND A
A0		160		AND B
A1		161		AND C
A2		162		AND D
A3		163		AND E
A4		164		AND H
A5		165		AND L
E6 XX		230 XX		AND N
CB 46		203 70		BIT 0,(HL)
DD CB dd 46		221 203 dd 70		BIT 0,(IX'd)
FD CB dd 46		253 203 dd 70		BIT 0,(IY'd)
CB 47		203 71		BIT 0,A
CB 40		203 64		BIT 0,B
CB 41		203 65		BIT 0,C
CB 42		203 66		BIT 0,D
CB 43		203 67		BIT 0,E
CB 44		203 68		BIT 0,H
CB 45		203 69		BIT 0,L
CB 4E		203 78		BIT 1,(HL)
DD CB dd 4E		221 203 dd 78		BIT 1,(IX'd)
FD CB dd 4E		253 203 dd 78		BIT 1,(IY'd)
CB 4F		203 79		BIT 1,A
CB 48		203 72		BIT 1,B
CB 49		203 73		BIT 1,C
CB 4A		203 74		BIT 1,D
CB 4B		203 75		BIT 1,E
CB 4C		203 76		BIT 1,H
CB 4D		203 77		BIT 1,L
CB 56		203 86		BIT 2,(HL)
DD CB dd 56		221 203 dd 86		BIT 2,(IX'd)
FD CB dd 56		253 203 dd 86		BIT 2,(IY'd)
CB 57		203 87		BIT 2,A
CB 50		203 80		BIT 2,B
CB 51		203 81		BIT 2,C
CB 52		203 82		BIT 2,D
CB 53		203 83		BIT 2,E
CB 54		203 84		BIT 2,H
CB 55		203 85		BIT 2,L

CB 5E	203 94	BIT 3,(HL)
DD CB dd 5E	221 203 dd 94	BIT 3,(IX'd)
FD CB dd 5E	253 203 dd 94	BIT 3,(IY'd)
CB 5F	203 95	BIT 3,A
CB 58	203 88	BIT 3,B
CB 59	203 89	BIT 3,C
CB 5A	203 90	BIT 3,D
CB 5B	203 91	BIT 3,E
CB 5C	203 92	BIT 3,H
CB 5D	203 93	BIT 3,L
CB 66	203 102	BIT 4,(HL)
DD CB dd 66	221 203 dd 102	BIT 4,(IX'd)
FD CB dd 66	253 203 dd 102	BIT 4,(IY'd)
CB 67	203 103	BIT 4,A
CB 60	203 96	BIT 4,B
CB 61	203 97	BIT 4,C
CB 62	203 98	BIT 4,D
CB 63	203 99	BIT 4,E
CB 64	203 100	BIT 4,H
CB 65	203 101	BIT 4,L
CB 6E	203 110	BIT 5,(HL)
DD CB dd 6E	221 203 dd 110	BIT 5,(IX'd)
FD CB dd 6E	253 203 dd 110	BIT 5,(IY'd)
CB 6F	203 111	BIT 5,A
CB 68	203 104	BIT 5,B
CB 69	203 105	BIT 5,C
CB 6A	203 106	BIT 5,D
CB 6B	203 107	BIT 5,E
CB 6C	203 108	BIT 5,H
CB 6D	203 109	BIT 5,L
CB 76	203 118	BIT 6,(HL)
DD CB dd 76	221 203 dd 118	BIT 6,(IX'd)
FD CB dd 76	253 203 dd 118	BIT 6,(IY'd)
CB 77	203 119	BIT 6,A
CB 70	203 112	BIT 6,B
CB 71	203 113	BIT 6,C
CB 72	203 114	BIT 6,D
CB 73	203 115	BIT 6,E
CB 74	203 116	BIT 6,H
CB 75	203 117	BIT 6,L
CB 7E	203 126	BIT 7,(HL)
DD CB dd 7E	221 203 dd 126	BIT 7,(IX'd)
FD CB dd 7E	253 203 dd 126	BIT 7,(IY'd)
CB 7F	203 127	BIT 7,A

CB 78	203 120	BIT 7,B
CB 79	203 121	BIT 7,C
CB 7A	203 122	BIT 7,D
CB 7B	203 123	BIT 7,E
CB 7C	203 124	BIT 7,H
CB 7D	203 125	BIT 7,L
DC XXXX	220 XXXX	CALL C,NN
FC XXXX	252 XXXX	CALL M,NN
D4 XXXX	212 XXXX	CALL NC,NN
CD XXXX	205 XXXX	CALL NN
C4 XXXX	196 XXXX	CALL NZ,NN
F4 XXXX	244 XXXX	CALL P,NN
EC XXXX	236 XXXX	CALL PE,NN
E4 XXXX	228 XXXX	CALL PO,NN
CC XXXX	204 XXXX	CALL Z,NN
3F	63	CCF
BE	190	CP (HL)
DD BE dd	221 190 dd	CP (IX'd)
FD BE dd	253 190 dd	CP (IY'd)
BF	191	CP A
B8	184	CP B
B9	185	CP C
BA	186	CP D
BB	187	CP E
BC	188	CP H
BD	189	CP L
FE XX	254 XX	CP N
ED A9	237 169	CPD
ED B9	237 185	CPDR
ED A1	237 161	CPI
ED B1	237 177	CPIR
2F	47	CPL
27	39	DAA
35	53	DEC (HL)
DD 35 dd	221 53 dd	DEC (IX'd)
FD 35 dd	253 53 dd	DEC (IY'd)
3D	61	DEC A
05	5	DEC B
0B	11	DEC BC
0D	13	DEC C
15	21	DEC D
1B	27	DEC DE
1D	29	DEC E
25	37	DEC H

2B		43		DEC HL
DD	2B	221	43	DEC IX
FD	2B	253	43	DEC IY
2D		45		DEC L
3B		59		DEC SP
F3		243		DI
10		16	XX	DJNZ N
FB		251		EI
E3		227		EX (SP),HL
DD	E3	221	227	EX (SP),IX
FD	E3	253	227	EX (SP),IY
08		8		EX AF,AF'
EB		235		EX DE,HL
D9		217		EXX
76		118		HALT
ED	46	237	70	IM 0
ED	56	237	86	IM 1
ED	5E	237	94	IM 2
ED	78	237	120	IN A,(C)
DB	XX	219	XX	IN A,(N)
ED	40	237	64	IN B,(C)
ED	48	237	72	IN C,(C)
ED	50	237	80	IN D,(C)
ED	58	237	88	IN E,(C)
ED	60	237	96	IN H,(C)
ED	68	237	104	IN L,(C)
34		52		INC (HL)
DD	34 dd	221	52 dd	INC (IX'd)
FD	34 dd	253	52 dd	INC (IY'd)
3C		60		INC A
04		4		INC B
03		3		INC BC
0C		12		INC C
14		20		INC D
13		19		INC DE
1C		28		INC E
24		36		INC H
23		35		INC HL
DD	23	221	35	INC IX
FD	23	253	35	INC IY
2C		44		INC L
33		51		INC SP
ED	AA	237	170	IND
ED	BA	237	186	INDR

ED A2	237 162	INI
ED B2	237 178	INIR
E9	233	JP (HL)
DD E9	221 233	JP (IX)
FD E9	253 233	JP (IY)
DA XXXX	218 XXXX	JP C,NN
FA XXXX	250 XXXX	JP M,NN
D2 XXXX	210	JP NC,NN
C3 XXXX	195 XXXX	JP NN
C2 XXXX	194 XXXX	JP NZ,NN
F2 XXXX	242 XXXX	JP P,NN
EA XXXX	234 XXXX	JP PE,NN
E2 XXXX	226 XXXX	JP PO,NN
CA XXXX	202 XXXX	JP Z,NN
38 XX	56 XX	JR C,N
18 XX	24 XX	JR N
30 XX	48 XX	JR NC,N
20 XX	32 XX	JR NZ,N
28 XX	40 XX	JR Z,N
0E XX	14 XX	LD C,N
02	2	LD (BC),A
12	18	LD (DE),A
77	119	LD (HL),A
70	112	LD (HL),B
71	113	LD (HL),C
72	114	LD (HL),D
73	115	LD (HL),E
74	116	LD (HL),H
75	117	LD (HL),L
36 XX	54 XX	LD (HL),N
DD 77 dd	221 119 dd	LD (IX'd),A
DD 70 dd	221 112 dd	LD (IX'd),B
DD 71 dd	221 113 dd	LD (IX'd),C
DD 72 dd	221 114 dd	LD (IX'd),D
DD 73 dd	221 115 dd	LD (IX'd),E
DD 74 dd	221 116 dd	LD (IX'd),H
DD 75 dd	221 117 dd	LD (IX'd),L
DD 36 dd XX	221 54 dd XX	LD (IX'd),N
FD 77 dd	253 119 dd	LD (IY'd),A
FD 70 dd	253 112 dd	LD (IY'd),B
FD 71 dd	253 113 dd	LD (IY'd),C
FD 72 dd	253 114 dd	LD (IY'd),D
FD 73 dd	253 115 dd	LD (IY'd),E
FD 74 dd	253 116 dd	LD (IY'd),H

FD 75 dd	253 117 dd	LD (IY'd),L
FD 36 dd XX	253 54 dd XX	LD (IY'd),N
32 XXXX	50 XXXX	LD (NN),A
ED 43 XXXX	237 67 XXXX	LD (NN),BC
ED 53 XXXX	237 83 XXXX	LD (NN),DE
22 XXXX	34 XXXX	LD (NN),HL
ED 63 XXXX	237 99 XXXX	LD (NN),HL
DD 22 XXXX	221 34 XXXX	LD (NN),IX
FD 22 XXXX	253 34 XXXX	LD (NN),IY
ED 73 XXXX	237 115 XXXX	LD (NN),SP
0A	10	LD A,(BC)
1A	26	LD A,(DE)
7E	126	LD A,(HL)
DD 7E dd	221 126 dd	LD A,(IX'd)
FD 7E dd	253 126 dd	LD A,(IY'd)
3A XXXX	58 XXXX	LD A,(NN)
7F	127	LD A,A
78	120	LD A,B
79	121	LD A,C
7A	122	LD A,D
7B	123	LD A,E
7C	124	LD A,H
ED 57	237 87	LD A,I
7D	125	LD A,L
3E XX	62 XX	LD A,N
ED 5F	237 95	LD A,R
46	70	LD B,(HL)
DD 46 dd	221 70 dd	LD B,(IX'd)
FD 46 dd	253 70 dd	LD B,(IY'd)
47	71	LD B,A
40	64	LD B,B
41	65	LD B,C
42	66	LD B,D
43	67	LD B,E
44	68	LD B,H
45	69	LD B,L
06 XX	6 XX	LD B,N
ED 4B XXXX	237 75 XXXX	LD BC,(NN)
01 XXXX	1 XXXX	LD BC,NN
4E	78	LD C,(HL)
DD 4E dd	221 78 dd	LD C,(IX'd)
FD 4E dd	253 78 dd	LD C,(IY'd)
4F	79	LD C,A
48	72	LD C,B

49		73		LD C,C
4A		74		LD C,D
4B		75		LD C,E
4C		76		LD C,H
4D		77		LD C,L
56		86		LD D,(HL)
DD	56 dd	221	86 dd	LD D,(IX'd)
FD	56 dd	253	86 dd	LD D,(IY'd)
57		87		LD D,A
50		80		LD D,B
51		81		LD D,C
52		82		LD D,D
53		83		LD D,E
54		84		LD D,H
55		85		LD D,L
16	XX	22	XX	LD D,N
ED	5B XXXX	237	91 XXXX	LD DE,(NN)
11		17	XXXX	LD DE,NN
5E		94		LD E,(HL)
DD	5E dd	221	94 dd	LD E,(IX'd)
FD	5E dd	253	94 dd	LD E,(IY'd)
5F		95		LD E,A
58		88		LD E,B
59		89		LD E,C
5A		90		LD E,D
5B		91		LD E,E
5C		92		LD E,H
5D		93		LD E,L
1E	XX	30	XX	LD E,N
66		102		LD H,(HL)
DD	66 dd	221	102 dd	LD H,(IX'd)
FD	66 dd	253	102 dd	LD H,(IY'd)
67		103		LD H,A
60		96		LD H,B
61		97		LD H,C
62		98		LD H,D
63		99		LD H,E
64		100		LD H,H
65		101		LD H,L
26	XX	38	XX	LD H,N
2A	XXXX	42	XXXX	LD HL,(NN)
ED	6B XXXX	237	107 XXXX	LD HL,(NN)
21	XXXX	33	XXXX	LD HL,NN
ED	47	237	71	LD I,A

DD 2A XXXX	221 42	LD IX,(NN)
DD 21 XXXX	221 33 XXXX	LD IX,NN
FD 2A XXXX	253 42	LD IY,(NN)
FD 21 XXXX	253 33 XXXX	LD IY,NN
6E	110	LD L,(HL)
DD 6E dd	221 110 dd	LD L,(IX'd)
FD 6E dd	253 110 dd	LD L,(IY'd)
6F	111	LD L,A
68	104	LD L,B
69	105	LD L,C
6A	106	LD L,D
6B	107	LD L,E
6C	108	LD L,H
6D	109	LD L,L
2E XX	46	LD L,N
ED 4F	237 79	LD R,A
ED 7B XXXX	237 123 XXXX	LD SP,(NN)
F9	249	LD SP,HL
DD F9	221 249	LD SP,IX
FD F9	253 249	LD SP,IY
31 XXXX	49 XXXX	LD SP,NN
ED A8	237 168	LDD
ED B8	237 184	LDDR
ED A0	237 160	LDI
ED B0	237 176	LDIR
ED 44	237 68	NEG
00	0	NOP
B6	182	OR (HL)
DD B6 dd	221 182 dd	OR (IX'd)
FD B6 dd	253 182 dd	OR (IY'd)
B7	183	OR A
B0	176	OR B
B1	177	OR C
B2	178	OR D
B3	179	OR E
B4	180	OR H
B5	181	OR L
F6 XX	246 XX	OR N
ED BB	237 187	OTDR
ED B3	237 179	OTIR
ED 79	237 121	OUT (C),A
ED 41	237 65	OUT (C),B
ED 49	237 73	OUT (C),C
ED 51	237 81	OUT (C),D

ED 59	237 89	OUT (C),E
ED 61	237 97	OUT (C),H
ED 69	237 105	OUT (C),L
D3 XX	211 XX	OUT (N),A
ED AB	237 171	OUTD
ED A3	237 163	OUTI
F1	241	POP AF
C1	193	POP BC
D1	209	POP DE
E1	225	POP HL
DD E1	221 225	POP IX
FD E1	253 225	POP IY
F5	245	PUSH AF
C5	197	PUSH BC
D5	213	PUSH DE
E5	229	PUSH HL
DD E5	221 229	PUSH IX
FD E5	253 229	PUSH IY
CB 86	203 134	RES 0,(HL)
DD CB dd 86	221 203 dd 134	RES 0,(IX'd)
FD CB dd 86	253 203 dd 134	RES 0,(IY'd)
CB 87	203 135	RES 0,A
CB 80	203 128	RES 0,B
CB 81	203 129	RES 0,C
CB 82	203 130	RES 0,D
CB 83	203 131	RES 0,E
CB 84	203 132	RES 0,H
CB 85	203 133	RES 0,L
CB 8E	203 142	RES 1,(HL)
DD CB dd 8E	221 203 dd 142	RES 1,(IX'd)
FD CB dd 8E	253 203 dd 142	RES 1,(IY'd)
CB 8F	203 143	RES 1,A
CB 88	203 136	RES 1,B
CB 89	203 137	RES 1,C
CB 8A	203 138	RES 1,D
CB 8B	203 139	RES 1,E
CB 8C	203 140	RES 1,H
CB 8D	203 141	RES 1,L
CB 96	203 150	RES 2,(HL)
DD CB dd 96	221 203 dd 150	RES 2,(IX'd)
FD CB dd 96	253 203 dd 150	RES 2,(IY'd)
CB 97	203 151	RES 2,A
CB 90	203 144	RES 2,B
CB 91	203 145	RES 2,C

CB 92	203 146	RES 2,D
CB 93	203 147	RES 2,E
CB 94	203 148	RES 2,H
CB 95	203 149	RES 2,L
CB 9E	203 158	RES 3,(HL)
DD CB dd 9E	221 203 dd 158	RES 3,(IX'd)
FD CB dd 9E	253 203 dd 158	RES 3,(IY'd)
CB 9F	203 159	RES 3,A
CB 98	203 152	RES 3,B
CB 99	203 153	RES 3,C
CB 9A	203 154	RES 3,D
CB 9B	203 155	RES 3,E
CB 9C	203 156	RES 3,H
CB 9D	203 157	RES 3,L
CB A6	203 166	RES 4,(HL)
DD CB dd A6	221 203 dd 166	RES 4,(IX'd)
FD CB dd A6	253 203 dd 166	RES 4,(IY'd)
CB A7	203 167	RES 4,A
CB A0	203 160	RES 4,B
CB A1	203 161	RES 4,C
CB A2	203 162	RES 4,D
CB A3	203 163	RES 4,E
CB A4	203 164	RES 4,H
CB A5	203 165	RES 4,L
CB AE	203 174	RES 5,(HL)
DD CB dd AE	221 203 dd 174	RES 5,(IX'd)
FD CB dd AE	253 203 dd 174	RES 5,(IY'd)
CB AF	203 175	RES 5,A
CB A8	203 168	RES 5,B
CB A9	203 169	RES 5,C
CB AA	203 170	RES 5,D
CB AB	203 171	RES 5,E
CB AC	203 172	RES 5,H
CB AD	203 173	RES 5,L
CB B6	203 182	RES 6,(HL)
DD CB dd B6	221 203 dd 182	RES 6,(IX'd)
FD CB dd B6	253 203 dd 182	RES 6,(IY'd)
CB B7	203 183	RES 6,A
CB B0	203 176	RES 6,B
CB B1	203 177	RES 6,C
CB B2	203 178	RES 6,D
CB B3	203 179	RES 6,E
CB B4	203 180	RES 6,H
CB B5	203 181	RES 6,L

CB BE	203 190	RES 7,(HL)
DD CB dd BE	221 203 dd 190	RES 7,(IX'd)
FD CB dd BE	253 203 dd 190	RES 7,(IY'd)
CB BF	203 191	RES 7,A
CB B8	203 184	RES 7,B
CB B9	203 185	RES 7,C
CB BA	203 186	RES 7,D
CB BB	203 187	RES 7,E
CB BC	203 188	RES 7,H
CB BD	203 189	RES 7,L
C9	201	RET
D8	216	RET C
F8	248	RET M
D0	208	RET NC
C0	192	RET NZ
F0	240	RET P
E8	232	RET PE
E0	224	RET PO
C8	200	RET Z
ED 4D	237 77	RETI
ED 45	237 69	RETN
CB 16	203 22	RL (HL)
DD CB dd 16	221 203 dd 20	RL (IX'd)
FD CB dd 16	253 203 dd 20	RL (IY'd)
CB 17	203 23	RL A
CB 10	203 16	RL B
CB 11	203 17	RL C
CB 12	203 18	RL D
CB 13	203 19	RL E
CB 14	203 20	RL H
CB 15	203 21	RL L
17	23	RLA
CB 06	203 6	RLC (HL)
DD CB dd 06	221 203 dd 6	RLC (IX'd)
FD CB dd 06	253 203 dd 6	RLC (IY'd)
CB 07	203 7	RLC A
CB 00	203 0	RLC B
CB 01	203 1	RLC C
CB 02	203 2	RLC D
CB 03	203 3	RLC E
CB 04	203 4	RLC H
CB 05	203 5	RLC L
07	7	RLCA
ED 6F	237 111	RLD

CB 1E	203 30	RR (HL)
DD CB dd 1E	221 203 dd 30	RR (IX'd)
FD CB dd 1E	253 203 dd 30	RR (IY'd)
CB 1F	203 31	RR A
CB 18	203 24	RR B
CB 19	203 25	RR C
CB 1A	203 26	RR D
CB 1B	203 27	RR E
CB 1C	203 28	RR H
CB 1D	203 29	RR L
1F	31	RRA
CB 0E	203 14	RRC (HL)
DD CB dd 0E	221 203 dd 14	RRC (IX'd)
FD CB dd 0E	253 203 dd 14	RRC (IY'd)
CB 0F	203 15	RRC A
CB 08	203 8	RRC B
CB 09	203 9	RRC C
CB 0A	203 10	RRC D
CB 0B	203 11	RRC E
CB 0C	203 12	RRC H
CB 0D	203 13	RRC L
0F	15	RRCA
ED 67	237 103	RRD
C7	199	RST 0
D7	215	RST 10
DF	223	RST 18
E7	231	RST 20
EF	239	RST 28
F7	247	RST 30
FF	255	RST 38
CF	207	RST 8
9E	158	SBC A,(HL)
DD 9E dd	221 158 dd	SBC A,(IX'd)
FD 9E dd	253 158 dd	SBC A,(IY'd)
9F	159	SBC A,A
98	152	SBC A,B
99	153	SBC A,C
9A	154	SBC A,D
9B	155	SBC A,E
9C	156	SBC A,H
9D	157	SBC A,L
DE XX	222 XX	SBC A,N
ED 42	237 66	SBC HL,BC
ED 52	237 82	SBC HL,DE

ED 62	237 98	SBC HL,HL
ED 72	237 114	SBC HL,SP
37	55	SCF
CB C6	203 198	SET 0,(HL)
DD CB dd C6	221 203 dd 198	SET 0,(IX'd)
FD CB dd C6	253 203 dd 198	SET 0,(IY'd)
CB C7	203 199	SET 0,A
CB C0	203 192	SET 0,B
CB C1	203 193	SET 0,C
CB C2	203 194	SET 0,D
CB C3	203 195	SET 0,E
CB C4	203 196	SET 0,H
CB C5	203 197	SET 0,L
CB CE	203 206	SET 1,(HL)
DD CB dd CE	221 203 dd 206	SET 1,(IX'd)
FD CB dd CE	253 203 dd 206	SET 1,(IY'd)
CB CF	203 207	SET 1,A
CB C8	203 200	SET 1,B
CB C9	203 201	SET 1,C
CB CA	203 202	SET 1,D
CB CB	203 203	SET 1,E
CB CC	203 204	SET 1,H
CB CD	203 205	SET 1,L
CB D6	203 214	SET 2,(HL)
DD CB dd D6	221 203 dd 214	SET 2,(IX'd)
FD CB dd D6	253 203 dd 214	SET 2,(IY'd)
CB D7	203 215	SET 2,A
CB D0	203 208	SET 2,B
CB D1	203 209	SET 2,C
CB D2	203 210	SET 2,D
CB D3	203 211	SET 2,E
CB D4	203 212	SET 2,H
CB D5	203 213	SET 2,L
CB DE	203 222	SET 3,(HL)
DD CB dd DE	221 203 dd 222	SET 3,(IX'd)
FD CB dd DE	253 203 dd 222	SET 3,(IY'd)
CB DF	203 223	SET 3,A
CB D8	203 216	SET 3,B
CB D9	203 217	SET 3,C
CB DA	203 218	SET 3,D
CB DB	203 219	SET 3,E
CB DC	203 220	SET 3,H
CB DD	203 221	SET 3,L
CB E6	203 230	SET 4,(HL)

DD CB dd E6	221 203 dd 230	SET 4,(IX'd)
FD CB dd E6	253 203 dd 230	SET 4,(IY'd)
CB E7	203 231	SET 4,A
CB E0	203 224	SET 4,B
CB E1	203 225	SET 4,C
CB E2	203 226	SET 4,D
CB E3	203 227	SET 4,E
CB E4	203 228	SET 4,H
CB E5	203 229	SET 4,L
CB EE	203 238	SET 5,(HL)
DD CB dd EE	221 203 dd 238	SET 5,(IX'd)
FD CB dd EE	253 203 dd 238	SET 5,(IY'd)
CB EF	203 239	SET 5,A
CB E8	203 232	SET 5,B
CB E9	203 233	SET 5,C
CB EA	203 234	SET 5,D
CB EB	203 235	SET 5,E
CB EC	203 236	SET 5,H
CB ED	203 237	SET 5,L
CB F6	203 246	SET 6,(HL)
DD CB dd F6	221 203 dd 246	SET 6,(IX'd)
FD CB dd F6	253 203 dd 246	SET 6,(IY'd)
CB F7	203 247	SET 6,A
CB F0	203 240	SET 6,B
CB F1	203 241	SET 6,C
CB F2	203 242	SET 6,D
CB F3	203 243	SET 6,E
CB F4	203 244	SET 6,H
CB F5	203 245	SET 6,L
CB FE	203 254	SET 7,(HL)
DD CB dd FE	221 203 dd 254	SET 7,(IX'd)
FD CB dd FE	253 203 dd 254	SET 7,(IY'd)
CB FF	203 255	SET 7,A
CB F8	203 248	SET 7,B
CB F9	203 249	SET 7,C
CB FA	203 250	SET 7,D
CB FB	203 251	SET 7,E
CB FC	203 252	SET 7,H
CB FD	203 253	SET 7,L
CB 26	203 38	SLA (HL)
DD CB dd 26	221 203 dd 38	SLA (IX'd)
FD CB dd 26	253 203 dd 38	SLA (IY'd)
CB 27	203 39	SLA A
CB 20	203 32	SLA B

CB 21	203 33	SLA C
CB 22	203 34	SLA D
CB 23	203 35	SLA E
CB 24	203 36	SLA H
CB 25	203 37	SLA L
CB 2E	203 46	SRA (HL)
DD CB dd 2E	221 203 dd 46	SRA (IX'd)
FD CB dd 2E	253 203 dd 46	SRA (IY'd)
CB 2F	203 47	SRA A
CB 28	203 40	SRA B
CB 29	203 41	SRA C
CB 2A	203 42	SRA D
CB 2B	203 43	SRA E
CB 2C	203 44	SRA H
CB 2D	203 45	SRA L
CB 3E	203 62	SRL (HL)
DD CB dd 3E	221 203 dd 62	SRL (IX'd)
FD CB dd 3E	253 203 dd 62	SRL (IY'd)
CB 3F	203 63	SRL A
CB 38	203 56	SRL B
CB 39	203 57	SRL C
CB 3A	203 58	SRL D
CB 3B	203 59	SRL E
CB 3C	203 60	SRL H
CB 3D	203 61	SRL L
96	150	SUB (HL)
DD 96 dd	221 150 dd	SUB (IX'd)
FD 96 dd	253 150 dd	SUB (IY'd)
97	151	SUB A
90	144	SUB B
91	145	SUB C
92	146	SUB D
93	147	SUB E
94	148	SUB H
95	149	SUB L
D6 XX	214 XX	SUB N
EE XX	238 XX	XOR N
AE	174	XOR (HL)
DD AE dd	221 174 dd	XOR (IX'd)
FD AE dd	253 174 dd	XOR (IY'd)
AF	175	XOR A
A8	168	XOR B
A9	169	XOR C

AA	170	XOR D
AB	171	XOR E
AC	172	XOR H
AD	173	XOR L

Appendix 2

Z80 instructions listed by opcode

00		0		NOP
01	XXXX	1	XXXX	LD BC,NN
02		2		LD (BC),A
03		3		INC BC
04		4		INC B
05		5		DEC B
06	XX	6	XX	LD B,N
07		7		RLCA
08		8		EX AF,AF'
09		9		ADD HL,BC
0A		10		LD A,(BC)
0B		11		DEC BC
0C		12		INC C
0D		13		DEC C
0E	XX	14	XX	LD C,N
0F		15		RRCA
10		16	XX	DJNZ N
11		17	XXXX	LD DE,NN
12		18		LD (DE),A
13		19		INC DE
14		20		INC D
15		21		DEC D
16	XX	22	XX	LD D,N
17		23		RLA
18	XX	24	XX	JR N
19		25		ADD HL,DE
1A		26		LD A,(DE)
1B		27		DEC DE
1C		28		INC E
1D		29		DEC E
1E	XX	30	XX	LD E,N
1F		31		RRA
20	XX	32	XX	JR NZ,N
21	XXXX	33	XXXX	LD HL,NN
22	XXXX	34	XXXX	LD (NN),HL

23		35	INC HL
24		36	INC H
25		37	DEC H
26	XX	38	LD H,N
27		39	DAA
28	XX	40	JR Z,N
29		41	ADD HL,HL
2A	XXXX	42	LD HL,(NN)
2B		43	DEC HL
2C		44	INC L
2D		45	DEC L
2E	XX	46	LD L,N
2F		47	CPL
30	XX	48	JR NC,N
31	XXXX	49	LD SP,NN
32	XXXX	50	LD (NN),A
33		51	INC SP
34		52	INC (HL)
35		53	DEC (HL)
36	XX	54	LD (HL),N
37		55	SCF
38	XX	56	JR C,N
39		57	ADD HL,SP
3A	XXXX	58	LD A,(NN)
3B		59	DEC SP
3C		60	INC A
3D		61	DEC A
3E	XX	62	LD A,N
3F		63	CCF
40		64	LD B,B
41		65	LD B,C
42		66	LD B,D
43		67	LD B,E
44		68	LD B,H
45		69	LD B,L
46		70	LD B,(HL)
47		71	LD B,A
48		72	LD C,B
49		73	LD C,C
4A		74	LD C,D
4B		75	LD C,E
4C		76	LD C,H
4D		77	LD C,L
4E		78	LD C,(HL)

4F	79	LD C,A
50	80	LD D,B
51	81	LD D,C
52	82	LD D,D
53	83	LD D,E
54	84	LD D,H
55	85	LD D,L
56	86	LD D,(HL)
57	87	LD D,A
58	88	LD E,B
59	89	LD E,C
5A	90	LD E,D
5B	91	LD E,E
5C	92	LD E,H
5D	93	LD E,L
5E	94	LD E,(HL)
5F	95	LD E,A
60	96	LD H,B
61	97	LD H,C
62	98	LD H,D
63	99	LD H,E
64	100	LD H,H
65	101	LD H,L
66	102	LD H,(HL)
67	103	LD H,A
68	104	LD L,B
69	105	LD L,C
6A	106	LD L,D
6B	107	LD L,E
6C	108	LD L,H
6D	109	LD L,L
6E	110	LD L,(HL)
6F	111	LD L,A
70	112	LD (HL),B
71	113	LD (HL),C
72	114	LD (HL),D
73	115	LD (HL),E
74	116	LD (HL),H
75	117	LD (HL),L
76	118	HALT
77	119	LD (HL),A
78	120	LD A,B
79	121	LD A,C
7A	122	LD A,D

7B	123	LD A,E
7C	124	LD A,H
7D	125	LD A,L
7E	126	LD A,(HL)
7F	127	LD A,A
80	128	ADD A,B
81	129	ADD A,C
82	130	ADD A,D
83	131	ADD A,E
84	132	ADD A,H
85	133	ADD A,L
86	134	ADD A,(HL)
87	135	ADD A,A
88	136	ADC A,B
89	137	ADC A,C
8A	138	ADC A,D
8B	139	ADC A,E
8C	140	ADC A,H
8D	141	ADC A,L
8E	142	ADC A,(HL)
8F	143	ADC A,A
90	144	SUB B
91	145	SUB C
92	146	SUB D
93	147	SUB E
94	148	SUB H
95	149	SUB L
96	150	SUB (HL)
97	151	SUB A
98	152	SBC A,B
99	153	SBC A,C
9A	154	SBC A,D
9B	155	SBC A,E
9C	156	SBC A,H
9D	157	SBC A,L
9E	158	SBC A,(HL)
9F	159	SBC A,A
A0	160	AND B
A1	161	AND C
A2	162	AND D
A3	163	AND E
A4	164	AND H
A5	165	AND L
A6	166	AND (HL)

A7	167	AND A
A8	168	XOR B
A9	169	XOR C
AA	170	XOR D
AB	171	XOR E
AC	172	XOR H
AD	173	XOR L
AE	174	XOR (HL)
AF	175	XOR A
B0	176	OR B
B1	177	OR C
B2	178	OR D
B3	179	OR E
B4	180	OR H
B5	181	OR L
B6	182	OR (HL)
B7	183	OR A
B8	184	CP B
B9	185	CP C
BA	186	CP D
BB	187	CP E
BC	188	CP H
BD	189	CP L
BE	190	CP (HL)
BF	191	CP A
C0	192	RET NZ
C1	193	POP BC
C2 XXXX	194 XXXX	JP NZ,NN
C3 XXXX	195 XXXX	JP NN
C4 XXXX	196 XXXX	CALL NZ,NN
C5	197	PUSH BC
C6 XX	198 XX	ADD A,N
C7	199	RST 0
C8	200	RET Z
C9	201	RET
CA XXXX	202 XXXX	JP Z,NN
CC XXXX	204 XXXX	CALL Z,NN
CD XXXX	205 XXXX	CALL NN
CE XX	206 XX	ADC A,N
CF	207	RST 8
D0	208	RET NC
D1	209	POP DE
D2 XXXX	210	JP NC,NN
D3 XX	211 XX	OUT (N),A

D4	XXXX	212	XXXX	CALL NC,NN
D5		213		PUSH DE
D6	XX	214	XX	SUB N
D7		215		RST 10
D8		216		RET C
D9		217		EXX
DA	XXXX	218	XXXX	JP C,NN
DB	XX	219	XX	IN A,(N)
DC	XXXX	220	XXXX	CALL C,NN
DE	XX	222	XX	SBC A,N
DF		223		RST 18
E0		224		RET PO
E1		225		POP HL
E2	XXXX	226	XXXX	JP PO,NN
E3		227		EX (SP),HL
E4	XXXX	228	XXXX	CALL PO,NN
E5		229		PUSH HL
E6	XX	230	XX	AND N
E7		231		RST 20
E8		232		RET PE
E9		233		JP (HL)
EA	XXXX	234	XXXX	JP PE,NN
EB		235		EX DE,HL
EC	XXXX	236	XXXX	CALL PE,NN
EE	XX	238	XX	XOR N
EF		239		RST 28
F0		240		RET P
F1		241		POP AF
F2	XXXX	242	XXXX	JP P,NN
F3		243		DI
F4	XXXX	244	XXXX	CALL P,NN
F5		245		PUSH AF
F6	XX	246	XX	OR N
F7		247		RST 30
F8		248		RET M
F9		249		LD SP,HL
FA	XXXX	250	XXXX	JP M,NN
FB		251		EI
FC	XXXX	252	XXXX	CALL M,NN
FE	XX	254	XX	CP N
FF		255		RST 38
CB 00		203	0	RLC B
CB 01		203	1	RLC C
CB 02		203	2	RLC D

CB 03	203 3	RLC E
CB 04	203 4	RLC H
CB 05	203 5	RLC L
CB 06	203 6	RLC (HL)
CB 07	203 7	RLC A
CB 08	203 8	RRC B
CB 09	203 9	RRC C
CB 0A	203 10	RRC D
CB 0B	203 11	RRC E
CB 0C	203 12	RRC H
CB 0D	203 13	RRC L
CB 0E	203 14	RRC (HL)
CB 0F	203 15	RRC A
CB 10	203 16	RL B
CB 11	203 17	RL C
CB 12	203 18	RL D
CB 13	203 19	RL E
CB 14	203 20	RL H
CB 15	203 21	RL L
CB 16	203 22	RL (HL)
CB 17	203 23	RL A
CB 18	203 24	RR B
CB 19	203 25	RR C
CB 1A	203 26	RR D
CB 1B	203 27	RR E
CB 1C	203 28	RR H
CB 1D	203 29	RR L
CB 1E	203 30	RR (HL)
CB 1F	203 31	RR A
CB 20	203 32	SLA B
CB 21	203 33	SLA C
CB 22	203 34	SLA D
CB 23	203 35	SLA E
CB 24	203 36	SLA H
CB 25	203 37	SLA L
CB 26	203 38	SLA (HL)
CB 27	203 39	SLA A
CB 28	203 40	SRA B
CB 29	203 41	SRA C
CB 2A	203 42	SRA D
CB 2B	203 43	SRA E
CB 2C	203 44	SRA H
CB 2D	203 45	SRA L
CB 2E	203 46	SRA (HL)

CB 2F	203 47	SRA A
CB 38	203 56	SRL B
CB 39	203 57	SRL C
CB 3A	203 58	SRL D
CB 3B	203 59	SRL E
CB 3C	203 60	SRL H
CB 3D	203 61	SRL L
CB 3E	203 62	SRL (HL)
CB 3F	203 63	SRL A
CB 40	203 64	BIT 0,B
CB 41	203 65	BIT 0,C
CB 42	203 66	BIT 0,D
CB 43	203 67	BIT 0,E
CB 44	203 68	BIT 0,H
CB 45	203 69	BIT 0,L
CB 46	203 70	BIT 0,(HL)
CB 47	203 71	BIT 0,A
CB 48	203 72	BIT 1,B
CB 49	203 73	BIT 1,C
CB 4A	203 74	BIT 1,D
CB 4B	203 75	BIT 1,E
CB 4C	203 76	BIT 1,H
CB 4D	203 77	BIT 1,L
CB 4E	203 78	BIT 1,(HL)
CB 4F	203 79	BIT 1,A
CB 50	203 80	BIT 2,B
CB 51	203 81	BIT 2,C
CB 52	203 82	BIT 2,D
CB 53	203 83	BIT 2,E
CB 54	203 84	BIT 2,H
CB 55	203 85	BIT 2,L
CB 56	203 86	BIT 2,(HL)
CB 57	203 87	BIT 2,A
CB 58	203 88	BIT 3,B
CB 59	203 89	BIT 3,C
CB 5A	203 90	BIT 3,D
CB 5B	203 91	BIT 3,E
CB 5C	203 92	BIT 3,H
CB 5D	203 93	BIT 3,L
CB 5E	203 94	BIT 3,(HL)
CB 5F	203 95	BIT 3,A
CB 60	203 96	BIT 4,B
CB 61	203 97	BIT 4,C
CB 62	203 98	BIT 4,D

CB 63	203 99	BIT 4,E
CB 64	203 100	BIT 4,H
CB 65	203 101	BIT 4,L
CB 66	203 102	BIT 4,(HL)
CB 67	203 103	BIT 4,A
CB 68	203 104	BIT 5,B
CB 69	203 105	BIT 5,C
CB 6A	203 106	BIT 5,D
CB 6B	203 107	BIT 5,E
CB 6C	203 108	BIT 5,H
CB 6D	203 109	BIT 5,L
CB 6E	203 110	BIT 5,(HL)
CB 6F	203 111	BIT 5,A
CB 70	203 112	BIT 6,B
CB 71	203 113	BIT 6,C
CB 72	203 114	BIT 6,D
CB 73	203 115	BIT 6,E
CB 74	203 116	BIT 6,H
CB 75	203 117	BIT 6,L
CB 76	203 118	BIT 6,(HL)
CB 77	203 119	BIT 6,A
CB 78	203 120	BIT 7,B
CB 79	203 121	BIT 7,C
CB 7A	203 122	BIT 7,D
CB 7B	203 123	BIT 7,E
CB 7C	203 124	BIT 7,H
CB 7D	203 125	BIT 7,L
CB 7E	203 126	BIT 7,(HL)
CB 7F	203 127	BIT 7,A
CB 80	203 128	RES 0,B
CB 81	203 129	RES 0,C
CB 82	203 130	RES 0,D
CB 83	203 131	RES 0,E
CB 84	203 132	RES 0,H
CB 85	203 133	RES 0,L
CB 86	203 134	RES 0,(HL)
CB 87	203 135	RES 0,A
CB 88	203 136	RES 1,B
CB 89	203 137	RES 1,C
CB 8A	203 138	RES 1,D
CB 8B	203 139	RES 1,E
CB 8C	203 140	RES 1,H
CB 8D	203 141	RES 1,L
CB 8E	203 142	RES 1,(HL)

CB 8F	203 143	RES 1, A
CB 90	203 144	RES 2, B
CB 91	203 145	RES 2, C
CB 92	203 146	RES 2, D
CB 93	203 147	RES 2, E
CB 94	203 148	RES 2, H
CB 95	203 149	RES 2, L
CB 96	203 150	RES 2, (HL)
CB 97	203 151	RES 2, A
CB 98	203 152	RES 3, B
CB 99	203 153	RES 3, C
CB 9A	203 154	RES 3, D
CB 9B	203 155	RES 3, E
CB 9C	203 156	RES 3, H
CB 9D	203 157	RES 3, L
CB 9E	203 158	RES 3, (HL)
CB 9F	203 159	RES 3, A
CB A0	203 160	RES 4, B
CB A1	203 161	RES 4, C
CB A2	203 162	RES 4, D
CB A3	203 163	RES 4, E
CB A4	203 164	RES 4, H
CB A5	203 165	RES 4, L
CB A6	203 166	RES 4, (HL)
CB A7	203 167	RES 4, A
CB A8	203 168	RES 5, B
CB A9	203 169	RES 5, C
CB AA	203 170	RES 5, D
CB AB	203 171	RES 5, E
CB AC	203 172	RES 5, H
CB AD	203 173	RES 5, L
CB AE	203 174	RES 5, (HL)
CB AF	203 175	RES 5, A
CB B0	203 176	RES 6, B
CB B1	203 177	RES 6, C
CB B2	203 178	RES 6, D
CB B3	203 179	RES 6, E
CB B4	203 180	RES 6, H
CB B5	203 181	RES 6, L
CB B6	203 182	RES 6, (HL)
CB B7	203 183	RES 6, A
CB B8	203 184	RES 7, B
CB B9	203 185	RES 7, C
CB BA	203 186	RES 7, D

CB BB	203 187	RES 7,E
CB BC	203 188	RES 7,H
CB BD	203 189	RES 7,L
CB BE	203 190	RES 7,(HL)
CB BF	203 191	RES 7,A
CB C0	203 192	SET 0,B
CB C1	203 193	SET 0,C
CB C2	203 194	SET 0,D
CB C3	203 195	SET 0,E
CB C4	203 196	SET 0,H
CB C5	203 197	SET 0,L
CB C6	203 198	SET 0,(HL)
CB C7	203 199	SET 0,A
CB C8	203 200	SET 1,B
CB C9	203 201	SET 1,C
CB CA	203 202	SET 1,D
CB CB	203 203	SET 1,E
CB CC	203 204	SET 1,H
CB CD	203 205	SET 1,L
CB CE	203 206	SET 1,(HL)
CB CF	203 207	SET 1,A
CB D0	203 208	SET 2,B
CB D1	203 209	SET 2,C
CB D2	203 210	SET 2,D
CB D3	203 211	SET 2,E
CB D4	203 212	SET 2,H
CB D5	203 213	SET 2,L
CB D6	203 214	SET 2,(HL)
CB D7	203 215	SET 2,A
CB D8	203 216	SET 3,B
CB D9	203 217	SET 3,C
CB DA	203 218	SET 3,D
CB DB	203 219	SET 3,E
CB DC	203 220	SET 3,H
CB DD	203 221	SET 3,L
CB DE	203 222	SET 3,(HL)
CB DF	203 223	SET 3,A
CB E0	203 224	SET 4,B
CB E1	203 225	SET 4,C
CB E2	203 226	SET 4,D
CB E3	203 227	SET 4,E
CB E4	203 228	SET 4,H
CB E5	203 229	SET 4,L
CB E6	203 230	SET 4,(HL)

CB E7	203 231	SET 4,A
CB E8	203 232	SET 5,B
CB E9	203 233	SET 5,C
CB EA	203 234	SET 5,D
CB EB	203 235	SET 5,E
CB EC	203 236	SET 5,H
CB ED	203 237	SET 5,L
CB EE	203 238	SET 5,(HL)
CB EF	203 239	SET 5,A
CB F0	203 240	SET 6,B
CB F1	203 241	SET 6,C
CB F2	203 242	SET 6,D
CB F3	203 243	SET 6,E
CB F4	203 244	SET 6,H
CB F5	203 245	SET 6,L
CB F6	203 246	SET 6,(HL)
CB F7	203 247	SET 6,A
CB F8	203 248	SET 7,B
CB F9	203 249	SET 7,C
CB FA	203 250	SET 7,D
CB FB	203 251	SET 7,E
CB FC	203 252	SET 7,H
CB FD	203 253	SET 7,L
CB FE	203 254	SET 7,(HL)
CB FF	203 255	SET 7,A
DD 09	221 9	ADD IX,BC
DD 19	221 25	ADD IX,DE
DD 21 XXXX	221 33 XXXX	LD IX,NN
DD 22 XXXX	221 34 XXXX	LD (NN),IX
DD 23	221 35	INC IX
DD 29	221 41	ADD IX,IX
DD 2A XXXX	221 42	LD IX,(NN)
DD 2B	221 43	DEC IX
DD 34 dd	221 52 dd	INC (IX'd)
DD 35 dd	221 53 dd	DEC (IX'd)
DD 36 dd XX	221 54 dd XX	LD (IX'd),N
DD 39	221 57	ADD IX,SP
DD 46 dd	221 70 dd	LD B,(IX'd)
DD 4E dd	221 78 dd	LD C,(IX'd)
DD 56 dd	221 86 dd	LD D,(IX'd)
DD 5E dd	221 94 dd	LD E,(IX'd)
DD 66 dd	221 102 dd	LD H,(IX'd)
DD 6E dd	221 110 dd	LD L,(IX'd)
DD 70 dd	221 112 dd	LD (IX'd),B

DD 71 dd	221 113 dd	LD (IX'd),C
DD 72 dd	221 114 dd	LD (IX'd),D
DD 73 dd	221 115 dd	LD (IX'd),E
DD 74 dd	221 116 dd	LD (IX'd),H
DD 75 dd	221 117 dd	LD (IX'd),L
DD 77 dd	221 119 dd	LD (IX'd),A
DD 7E dd	221 126 dd	LD A,(IX'd)
DD 86 dd	221 134 dd	ADD A,(IX'd)
DD 8E dd	221 142 dd	ADC A,(IX'd)
DD 96 dd	221 150 dd	SUB (IX'd)
DD 9E dd	221 158 dd	SBC A,(IX'd)
DD A6 dd	221 166 dd	AND (IX'd)
DD AE dd	221 174 dd	XOR (IX'd)
DD B6 dd	221 182 dd	OR (IX'd)
DD BE dd	221 190 dd	CP (IX'd)
DD E1	221 225	POP IX
DD E3	221 227	EX (SP),IX
DD E5	221 229	PUSH IX
DD E9	221 233	JP (IX)
DD F9	221 249	LD SP,IX
DD CB dd 06	221 203 dd 6	RLC (IX'd)
DD CB dd 0E	221 203 dd 14	RRC (IX'd)
DD CB dd 16	221 203 dd 20	RL (IX'd)
DD CB dd 1E	221 203 dd 30	RR (IX'd)
DD CB dd 26	221 203 dd 38	SLA (IX'd)
DD CB dd 2E	221 203 dd 46	SRA (IX'd)
DD CB dd 3E	221 203 dd 62	SRL (IX'd)
DD CB dd 46	221 203 dd 70	BIT 0,(IX'd)
DD CB dd 4E	221 203 dd 78	BIT 1,(IX'd)
DD CB dd 56	221 203 dd 86	BIT 2,(IX'd)
DD CB dd 5E	221 203 dd 94	BIT 3,(IX'd)
DD CB dd 66	221 203 dd 102	BIT 4,(IX'd)
DD CB dd 6E	221 203 dd 110	BIT 5,(IX'd)
DD CB dd 76	221 203 dd 118	BIT 6,(IX'd)
DD CB dd 7E	221 203 dd 126	BIT 7,(IX'd)
DD CB dd 86	221 203 dd 134	RES 0,(IX'd)
DD CB dd 8E	221 203 dd 142	RES 1,(IX'd)
DD CB dd 96	221 203 dd 150	RES 2,(IX'd)
DD CB dd 9E	221 203 dd 158	RES 3,(IX'd)
DD CB dd A6	221 203 dd 166	RES 4,(IX'd)
DD CB dd AE	221 203 dd 174	RES 5,(IX'd)
DD CB dd B6	221 203 dd 182	RES 6,(IX'd)
DD CB dd BE	221 203 dd 190	RES 7,(IX'd)
DD CB dd C6	221 203 dd 198	SET 0,(IX'd)

DD CB dd CE	221 203 dd 206	SET 1,(IX'd)
DD CB dd D6	221 203 dd 214	SET 2,(IX'd)
DD CB dd DE	221 203 dd 222	SET 3,(IX'd)
DD CB dd E6	221 203 dd 230	SET 4,(IX'd)
DD CB dd EE	221 203 dd 238	SET 5,(IX'd)
DD CB dd F6	221 203 dd 246	SET 6,(IX'd)
DD CB dd FE	221 203 dd 254	SET 7,(IX'd)
ED 40	237 64	IN B,(C)
ED 41	237 65	OUT (C),B
ED 42	237 66	SBC HL,BC
ED 43 XXXX	237 67 XXXX	LD (NN),BC
ED 44	237 68	NEG
ED 45	237 69	RETN
ED 46	237 70	IM 0
ED 47	237 71	LD I,A
ED 48	237 72	IN C,(C)
ED 49	237 73	OUT (C),C
ED 4A	237 74	ADC HL,BC
ED 4B XXXX	237 75 XXXX	LD BC,(NN)
ED 4D	237 77	RETI
ED 4F	237 79	LD R,A
ED 50	237 80	IN D,(C)
ED 51	237 81	OUT (C),D
ED 52	237 82	SBC HL,DE
ED 53 XXXX	237 83 XXXX	LD (NN),DE
ED 56	237 86	IM 1
ED 57	237 87	LD A,I
ED 58	237 88	IN E,(C)
ED 59	237 89	OUT (C),E
ED 5A	237 90	ADC HL,DE
ED 5B XXXX	237 91 XXXX	LD DE,(NN)
ED 5E	237 94	IM 2
ED 5F	237 95	LD A,R
ED 60	237 96	IN H,(C)
ED 61	237 97	OUT (C),H
ED 62	237 98	SBC HL,HL
ED 63 XXXX	237 99 XXXX	LD (NN),HL
ED 67	237 103	RRD
ED 68	237 104	IN L,(C)
ED 69	237 105	OUT (C),L
ED 6A	237 106	ADC HL,HL
ED 6B XXXX	237 107 XXXX	LD HL,(NN)
ED 6F	237 111	RLD
ED 72	237 114	SBC HL,SP

ED 73	XXXX	237 115	XXXX	LD (NN),SP
ED 78		237 120		IN A,(C)
ED 79		237 121		OUT (C),A
ED 7A		237 122		ADC HL,SP
ED 7B	XXXX	237 123	XXXX	LD SP,(NN)
ED A0		237 160		LDI
ED A1		237 161		CPI
ED A2		237 162		INI
ED A3		237 163		OUTI
ED A8		237 168		LDD
ED A9		237 169		CPD
ED AA		237 170		IND
ED AB		237 171		OUTD
ED B0		237 176		LDIR
ED B1		237 177		CPIR
ED B2		237 178		INIR
ED B3		237 179		OTIR
ED B8		237 184		LDDR
ED B9		237 185		CPDR
ED BA		237 186		INDR
ED BB		237 187		OTDR
FD 09		253 9		ADD IY,BC
FD 19		253 25		ADD IY,DE
FD 21	XXXX	253 33	XXXX	LD IY,NN
FD 22	XXXX	253 34	XXXX	LD (NN),IY
FD 23		253 35		INC IY
FD 29		253 41		ADD IY,IY
FD 2A	XXXX	253 42		LD IY,(NN)
FD 2B		253 43		DEC IY
FD 34	dd	253 52	dd	INC (IY'd)
FD 35	dd	253 53	dd	DEC (IY'd)
FD 36	dd XX	253 54	dd XX	LD (IY'd),N
FD 39		253 57		ADD IY,SP
FD 46	dd	253 70	dd	LD B,(IY'd)
FD 4E	dd	253 78	dd	LD C,(IY'd)
FD 56	dd	253 86	dd	LD D,(IY'd)
FD 5E	dd	253 94	dd	LD E,(IY'd)
FD 66	dd	253 102	dd	LD H,(IY'd)
FD 6E	dd	253 110	dd	LD L,(IY'd)
FD 70	dd	253 112	dd	LD (IY'd),B
FD 71	dd	253 113	dd	LD (IY'd),C
FD 72	dd	253 114	dd	LD (IY'd),D
FD 73	dd	253 115	dd	LD (IY'd),E
FD 74	dd	253 116	dd	LD (IY'd),H

FD 75 dd	253 117 dd	LD (IY'd),L
FD 77 dd	253 119 dd	LD (IY'd),A
FD 7E dd	253 126 dd	LD A,(IY'd)
FD 86 dd	253 134 dd	ADD A,(IY'd)
FD 8E dd	253 142 dd	ADC A,(IY'd)
FD 96 dd	253 150 dd	SUB (IY'd)
FD 9E dd	253 158 dd	SBC A,(IY'd)
FD A6 dd	253 166 dd	AND (IY'd)
FD AE dd	253 174 dd	XOR (IY'd)
FD B6 dd	253 182 dd	OR (IY'd)
FD BE dd	253 190 dd	CP (IY'd)
FD E1	253 225	POP IY
FD E3	253 227	EX (SP),IY
FD E5	253 229	PUSH IY
FD E9	253 233	JP (IY)
FD F9	253 249	LD SP,IY
FD CB dd 06	253 203 dd 6	RLC (IY'd)
FD CB dd 0E	253 203 dd 14	RRC (IY'd)
FD CB dd 16	253 203 dd 20	RL (IY'd)
FD CB dd 1E	253 203 dd 30	RR (IY'd)
FD CB dd 26	253 203 dd 38	SLA (IY'd)
FD CB dd 2E	253 203 dd 46	SRA (IY'd)
FD CB dd 3E	253 203 dd 62	SRL (IY'd)
FD CB dd 46	253 203 dd 70	BIT 0,(IY'd)
FD CB dd 4E	253 203 dd 78	BIT 1,(IY'd)
FD CB dd 56	253 203 dd 86	BIT 2,(IY'd)
FD CB dd 5E	253 203 dd 94	BIT 3,(IY'd)
FD CB dd 66	253 203 dd 102	BIT 4,(IY'd)
FD CB dd 6E	253 203 dd 110	BIT 5,(IY'd)
FD CB dd 76	253 203 dd 118	BIT 6,(IY'd)
FD CB dd 7E	253 203 dd 126	BIT 7,(IY'd)
FD CB dd 86	253 203 dd 134	RES 0,(IY'd)
FD CB dd 8E	253 203 dd 142	RES 1,(IY'd)
FD CB dd 96	253 203 dd 150	RES 2,(IY'd)
FD CB dd 9E	253 203 dd 158	RES 3,(IY'd)
FD CB dd A6	253 203 dd 166	RES 4,(IY'd)
FD CB dd AE	253 203 dd 174	RES 5,(IY'd)
FD CB dd B6	253 203 dd 182	RES 6,(IY'd)
FD CB dd BE	253 203 dd 190	RES 7,(IY'd)
FD CB dd C6	253 203 dd 198	SET 0,(IY'd)
FD CB dd CE	253 203 dd 206	SET 1,(IY'd)
FD CB dd D6	253 203 dd 214	SET 2,(IY'd)
FD CB dd DE	253 203 dd 222	SET 3,(IY'd)

240 Appendix 2 – Z80 instructions listed by opcode

FD CB dd E6	253 203 dd 230	SET 4,(IY'd)
FD CB dd EE	253 203 dd 238	SET 5,(IY'd)
FD CB dd F6	253 203 dd 246	SET 6,(IY'd)
FD CB dd FE	253 203 dd 254	SET 7,(IY'd)

Appendix 3

Flag operation table

Flag table notation

Flags

- Flag is unchanged by operation.
- * Flag is affected according to result of operation.
- P P/V is set according to parity result.
- V P/V is set according to the overflow result.
- 0 Flag is set to zero
- 1 Flag is set to one.
- # Result of flag unknown.
- f Contents of the interrupt flip flop.

Addressing

- s Any 8 bit addressing mode A, B, C, D, E, H, L, (HL), (IX+dd), (IY+dd)
- r Any 8 bit register A, C, D, E, H, L
- b Bit number 0-7
- RR Any 16 bit register.
- n Any 8 bit number

Flag operation table

Instruction	C	Z	P/V	S	N	H	Instruction	C	Z	P/V	S	N	H
ADC HL,RR	*	*	V	*	0	#	CCF	*	-	-	-	0	#
ADC A,s	*	*	V	*	0	*	CPD	-	*	*	#	1	#
ADC A,n	*	*	V	*	0	*	CPDR	-	*	*	#	1	#
ADD A,s	*	*	V	*	0	*	CPI	-	*	*	#	1	#
ADD A,n	*	*	V	*	0	*	CPIR	-	*	*	#	1	#
ADD HL,RR	*	-	-	-	0	#	CP s	*	*	V	*	1	*
ADD SP,RR	*	-	-	-	0	#	CP n	*	*	V	*	1	*
ADD IX,RR	*	-	-	-	0	#	CPL	-	-	-	-	1	1
ADD IY,RR	*	-	-	-	0	#	DAA	*	*	P	*	-	*
AND s	0	*	P	*	0	1	DEC s	-	*	V	*	1	*
AND n	0	*	P	*	0	1	IN r,(C)	-	*	P	*	0	0
BIT b,s	-	*	#	#	0	1	INC s	-	*	V	*	0	*

Instruction	C	Z	P/V	S	N	H	Instruction	C	Z	P/V	S	N	H
IND	-	*	#	#	1	#	RRA	*	-	-	-	0	0
INI	-	*	#	#	1	#	RRCA	*	-	-	-	0	0
INDR	-	1	#	#	1	#	RLD (HL)	-	*	P	*	0	0
INIR	-	1	#	#	1	#	RRD (HL)	-	*	P	*	0	0
LD A,I	-	*	f	*	0	0	RL s	*	*	P	*	0	0
LD A,R	-	*	f	*	0	0	RLC s	*	*	P	*	0	0
LDD	-	#	*	#	0	0	RR s	*	*	P	*	0	0
LDI	-	#	*	#	0	0	RRC s	*	*	P	*	0	0
LDDR	-	#	0	#	0	0	SLA s	*	*	P	*	0	0
LDIR	-	#	0	#	0	0	SRA s	*	*	P	*	0	0
NEG	*	*	V	*	1	*	SRL s	*	*	P	*	0	0
OR s	0	*	P	*	0	0	SBC HL,RR	*	*	V	*	1	#
OR n	0	*	P	*	0	0	SCF	1	-	-	-	0	0
OTDR	-	1	#	#	1	#	SBC A,s	*	*	V	*	1	*
OTIR	-	1	#	#	1	#	SBC A,n	*	*	V	*	1	*
OUTD	-	*	#	#	1	#	SUB s	*	*	V	*	1	*
OUTI	-	*	#	#	1	#	SUB n	*	*	V	*	1	*
RLA	*	-	-	-	0	0	XOR s	0	*	P	*	0	0
RLCA	*	-	-	-	0	0	XOR n	0	*	P	*	0	0

Appendix 4

Spectrum monitor- assembler listing

```

                ASEG
                ORG      25500D

                JP      FSTART
ALTER EQU      39D
COMMA EQU     ','

SAVEB: LD      A, 0FFH
        ;IX POINTS TO START OF BLOCK
        ;DE CONTAINS NUMBER OF BYTES
        CALL   04C2H
        RET

LOADB: SCF
        LD      A, 0FFH
        CALL   0556H
        RET

HEADIN:
        LD      DE, 17
        LD      IX, HEADER
        XOR     A
        SCF
        CALL   0556H
        LD      A, (HEADER+11)
        LD      C, A
        LD      A, '$'
        LD      (HEADER+11), A
        CALL   CRLF

        LD      DE, HEADER+1
        CALL   PRTSTR
        LD      A, ' '
        CALL   PRTCHR
```

```

EX      DE, HL
LD      ( HL ), C
INC     HL
INC     HL
INC     HL

```

```

LD      A, ( HL )
CALL    HEX0
DEC     HL
LD      A, ( HL )
CALL    HEX0
DEC     HL
LD      A, ' '
CALL    PRTCHR

```

```

LD      A, ( HL )
CALL    HEX0
DEC     HL
LD      A, ( HL )
CALL    HEX0

```

```

CALL    CRLF
RET

```

HEADOUT:

```

LD      DE, 17
LD      IX, HEADER
XOR     A
CALL    04C2H
RET

```

```

KEY:    PUSH    HL
        PUSH    BC
        PUSH    DE
        CALL    WAITS

```

```

WAITK:  LD      A, ( 23611 )
        BIT     5, A
        JR     Z, WAITK
        RES    5, A
        LD     ( 23611 ), A
        POP    DE

```

```

; PAUSE FOR A WHILE
; LOOK AT FLAGS

```

```

; NO KEY PRESSED
; RESET FLAG

```

```

POP      BC
POP      HL
RET

GETKEY:
CALL     KEY
LD       A, (23560)      ;LOOK AT LAST-K
CALL     PRTCHR
RET

OPENCH2:
LD       A, 02           ;OPENS CHANNEL 'S'
                           ;FOR PRINTING
CALL     1601H
RET

PRTCHR:
PUSH     AF
PUSH     AF
XOR      A
LD       (23692), A
POP      AF
RST      10H
POP      AF
RET

PRTSTR:
LD       A, (DE)        ;GET CHARACTER
CP       '$'           ;IS THIS THE END
                           ;OF A STRING?
RET      Z              ;YES THEN RETURN
CALL     PRTCHR        ;PRINT CHARACTER
INC      DE             ;POINT TO NEXT CHARACTER
JR       PRTSTR

FSTART:
LD       SP, STACK
CALL     OPENCH2
LD       DE, HOME
CALL     PRTSTR
LD       DE, WELCMSS
CALL     PRTSTR
CALL     CRLF

VERYSTART:
INITV2: LD       SP, STACK
LD       A, 8
LD       (23658), A    ;CAPS ON

```

```

        LD      HL, ERRSP
        LD      (HL), LOW(VERYSTART)
        INC    HL
        LD      (HL), HIGH(VERYSTART)
        DEC    HL
        LD      (23613), HL

INITV:  LD      HL, INITV
        PUSH   HL

        CALL   CRLF
        LD      A, ' '
        CALL   PRTCHR

START:  CALL   GETKEY
        SUB    'A'                ; IS IT IN THE ALPHABET?
        RET    C                  ; NO
        CP    'S' - 'A' + 1
        RET    NC                 ; NO!
        ADD   A, A                ; *2
        LD    HL, VECTBL
        LD    E, A
        LD    D, 0
        ADD   HL, DE
        LD    E, (HL)
        INC  HL
        LD    D, (HL)
        EX   DE, HL
        JP   (HL)                ; JUMP TO COMMAND

VECTBL:
        DEFW  ERROR
        DEFW  ERROR
        DEFW  ERROR
        DEFW  DUMP                ; DUMP
        DEFW  MODIFY              ; EDIT MEMORY
        DEFW  FILL                ; FILL
        DEFW  GOTO                ; GOTO
        DEFW  HUNT                ; HUNT
        DEFW  IDENT              ; IDENTIFY FILENAME
        DEFW  ERROR
        DEFW  ERROR

```

```

DEFW LOADBYTES ;LOAD FROM TAPE
DEFW MOVE ;MOVE A BLOCK
DEFW ERROR
DEFW ERROR
DEFW PUMP ;PRINT DUMP
DEFW ERROR
DEFW CHREG ;MODIFY REGS
DEFW SAVEBYTES ;SAVE MEMORY

```

WAITS:

```

LD BC, 8000H
LD DE, 4000H
LD HL, 4000H
LDIR
RET

```

MODIFY:

```

CALL GETEXPR1 ;GET START ADDRESS
MODIFB: CALL HEXOD ;OUTPUT START ADDRESS
LD A, ' '
CALL PRTCHR
LD A, (HL)
CALL HEXD
LD A, ' '
CALL PRTCHR
PUSH HL
CALL HEXI
POP HL
LD (HL), A
INC HL
CALL CRLF
JR MODIFY ;LOOP UNTIL FORCED
;OUT BY AN ERROR

```

```

GETEXPR2: ;GET TWO WORD EXPRESSION
CALL HEXD
PUSH HL ;SAVE E1

```

```

LD      A, ' '
CALL   PRTCHR
CALL   HEXD
PUSH   HL
POP    DE      ;E2
POP    HL      ;GET E1
RET    ;E1=HL E2=DE

```

GETEXPR3:

```

CALL   GETEXPR2
PUSH   HL
PUSH   DE
CALL   GETEXPR1
POP    DE
POP    HL
RET

```

GETEXPR1:

```

LD      A, ' '
CALL   PRTCHR
CALL   HEXD
PUSH   HL
POP    BC
CALL   CRLF
RET    ;E1=HL E2=DE E3=BC

```

LOADBYTES:

```

LD      A, ' '
CALL   PRTCHR
CALL   GETEXPR2

LD      (DESTT), HL

LD      A, E
OR      D
JP      Z, ERRORS

LD      (LENTT), DE

LD      DE, LOADMESS
CALL   PRTSTR

```



```

LD      DE, FILENAME
CALL   PRTSTR

GETFILE:
CALL   HEADIN

LD      DE, HEADER+1
LD      HL, FILENAME
LD      B, 10

COMPF: LD      A, (DE)
RES     5, A
LD      C, (HL)
RES     5, C
CP      C
JR      NZ, GETFILE

INC     HL
INC     DE
DJNZ   COMPF

LOADIF: LD      DE, (LENTT)
LD      IX, (DETT)
CALL   LOADB
RET

DETT:  DEFW    0
LENTT: DEFW    0

CR     EQU    0DH          ;RETURN
LF     EQU    0AH          ;LINEFEED
WELCMESS: DEFM  CR, '*SBUG*'
        DEFM  ' (C) John Wilson'
        DEFM  ' 1984.'
        DEFM  CR, '$'
SAVEMESS: DEFM  CR, 'Press any key when ready$'

```

```

LOADMESS: DEFM CR, 'Waiting for $'
NOTMESS:  DEFM CR, 'ROUTINE NOT IMPLEMENTED$'
ERRMESS:  DEFM CR, '**ERROR**', CR, '$'
HOME:     DEFM 22, 1, 1, CR, '$'

```

SAVEBYTES:

```

LD      A, ' '
CALL    PRTCHR
CALL    GETEXPR2      ;GET 2 VALUES START
                        ;AND NUMBER OF BYTES

LD      (HEADER+13), HL
LD      A, E
OR      D
JP      Z, ERRORS
LD      (HEADER+11), DE

LD      DE, SAVEMESS
CALL    PRTSTR

CALL    WAITS
CALL    WAITS
CALL    WAITS
CALL    GETKEY

LD      A, 3
LD      DE, HEADER
LD      HL, FILENAME
LD      (DE), A
INC     DE

LD      BC, 10
LDIR
CALL    HEADOUT
CALL    WAITS
CALL    WAITS
CALL    WAITS
LD      IX, (HEADER+13)
LD      DE, (HEADER+11)
CALL    SAVEB
RET

```

```

HEXAS:  PUSH    HL
        CALL    HEXD
        LD     A, ' '
        CALL    PRTCHR
        POP    HL
        RET

LINE:   LD     B, 8
NBYTE:  LD     A, (HL)
        CALL    HEXAS
        INC   HL
        DJNZ  NBYTE
        CALL    CRLF
        RET

DUMP:   LD     A, ' '
        CALL    PRTCHR
        CALL    GETEXPR1
ALOCK:  LD     C, 8
BLOCK:  CALL    HEXDD
        LD     A, ' '
        CALL    PRTCHR
        CALL    PRTCHR
        CALL    LINE
        DEC   C
        JR    NZ, BLOCK
        CALL    CRLF
        CALL    CRLF
        CALL    GETKEY
        CP    CR
        JR    Z, ALOCK
        RET

PINE:   LD     B, 21
PBYTE:  LD     A, (HL)
        CP    32
        JR    C, SBOGGY-2
        CP    128
        JR    C, SBOGGY
        LD     A, ' '

```

```

SBOGGY: CALL    PRTCHR
         INC     HL
         DJNZ   PBYTE
         CALL   CRLF
         RET

```

```

PUMP:
        LD      A, ' '
        CALL   PRTCHR
        CALL   GETEXPR1
PALOCK: LD      C, 8
PLOCK: CALL    HEXOD
        LD      A, ' '
        CALL   PRTCHR
        CALL   PRTCHR
        CALL   PINE
        DEC    C
        JR     NZ, PLOCK
        CALL   CRLF
        CALL   CRLF
        CALL   GETKEY
        CP     CR
        JR     Z, PALOCK
        RET

```

```

ERROR:  NOP
NOTIMP: PUSH    DE
        LD     DE, NOTMESS
        CALL  PRTSTR
        POP   DE
        RET

```

;ROUTINES

```

;HEXD  OUTPUT HEX NUMBER IN ACCUMULATOR
;HEXOD OUTPUT HEX WORD IN HL
;HEXI  INPUT HEX NUMBER PUT IN ACCUMULATOR
;HEXD  INPUT HEX WORD AND PUT INTO HL

```

HEX0D:

```
LD      A, H
CALL   HEX0
LD      A, L
```

HEX0:

```
LD      E, A
SRL    A          ;GET TOP FOUR BITS
                    ;INTO LOWER NYBBLE

SRL    A
SRL    A
SRL    A
CALL   CONV      ;CONVERT TO ASCII
                    ;RETURNS ASCII VALUE IN A
LD      A, E      ;GET ORIGINAL VALUE
AND    0FH       ;MASK OFF LOWER FOUR BITS
```

; CONVERT LAST HEX DIGIT

CONV:

```
ADD    A, 30H
CP     3AH       ;IS DIGIT IN RANGE 0-9?
JP     M, DECD  ;YES THEN PRINT AND RETURN
                    ;IN THE RANGE 10-15 SO CONVERT TO A-F
ADD    A, 7
```

```
DECD:  CALL   PRTCHR      ;PRINT A HEX DIGIT
        RET
```

ERRORS: JP VERYSTART

HEXI:

```
CALL   GETKEY
CALL   CONV2
LD      E, A
CALL   GETKEY
CALL   CONV2
SLA    E          ;MOVE LOWER FOUR BITS UP
SLA    E
SLA    E
SLA    E
OR     E          ;MERGE IN SECOND DIGIT
RET
```

```
CONV2: AND    A
        SBC   A, 30H
        CP    0AH
```

```

RET      C
AND      A
SBC      A, 7
CP       10H
JR       NC, ERRORS
RET

HEXD:    CALL    HEXI
         PUSH    AF
         CALL    HEXI
         LD     L, A
         POP     AF
         LD     H, A
         RET

FILL:    ;HL POINTS TO START ADDRESS
         ;DE POINTS END ADDRESS
         ;BC =NUMBER OF BYTES

LD       A, ' '
CALL     PRTCHR
CALL     GETEXPR2
LD       A, ' '
CALL     PRTCHR

PUSH     HL
EX       DE, HL
AND      A           ;CLEAR CARRY
SBC      HL, DE
JP       C, ERRORS
JP       Z, ERRORS
PUSH     HL
POP      BC

POP      HL
PUSH     HL
POP      DE
INC      DE

PUSH     HL
PUSH     DE
CALL     HEXI

```

```

POP      DE
POP      HL

LD       (HL), A
LDIR
CALL    CRLF
RET

RETGET:  POP      DE           ;GET ORIGINAL
LD       HL, PUTREG
PUSH    HL
PUSH    DE
CALL    GETREG
RET

GOTO:

LD       A, ' '
CALL    PRTCHR
CALL    HEXD
PUSH    HL           ;GOTO
LD       A, ' '
CALL    PRTCHR

;POP ALL REGS VALUES
CALL    GETKEY
CP      CR
JR      Z, RETGET
CP      COMMA
JP      NZ, ERRORS
CALL    HEXD
PUSH    HL
LD      DE, BRKP
LD      BC, 3
LDIR                               ;SAVE BYTES

POP     HL
LD      (HL), 0CDH
INC     HL
LD      (HL), LOW(BRK)

```

```

INC      HL
LD       ( HL ), HIGH( BRK )
CALL    GETREG
RET

PUTREG: LD      ( SAVESP ), SP
LD      SP, AHLREG+2
EX      AF, AF'
EXX
PUSH    HL
PUSH    DE
PUSH    BC
PUSH    AF
EXX
EX      AF, AF'
PUSH    IX
PUSH    HL
PUSH    DE
PUSH    BC
PUSH    AF
LD      SP, ( SAVESP )
RET

GETREG: LD      ( SAVESP ), SP
LD      SP, AFREG
POP     AF
POP     BC
POP     DE
POP     HL
POP     IX
EX      AF, AF'
EXX
POP     AF
POP     BC
POP     DE
POP     HL
EX      AF, AF'
EXX
LD      SP, ( SAVESP )
RET

BRK:    ;PUSH  ALL VALUES ON STACK
CALL   PUTREG

```



```

POP      HL                      ;RET ADDR5
DEC      HL
DEC      HL
DEC      HL
;BACK SPACE 3 INSTR

CALL     CRLF
LD       A, '*'
CALL     PRTCHR
CALL     HEXOD
EX       DE, HL                  ;DEST
LD       HL, BRKP
LD       BC, 3
LDIR                    ;PUT BYTES BACK

;PUT BACK FOUR BYTES
;DISPLAY PC
;DISPR:
CALL     DISPR
RET

OUTREG:  LD       B, 4
NXTREG:  LD       E, (HL)        ;LOW
         INC     HL
         LD       D, (HL)
         PUSH    HL
         EX      DE, HL
DI1R:    CALL     HEXOD
         LD       A, ' '
         CALL     PRTCHR
         CALL     PRTCHR
         POP     HL
         INC     HL
         DJNZ   NXTREG
         RET

DISPR:   CALL     CRLF
         LD       DE, REGMESS
         CALL     PRTSTR
         LD       HL, AFREG
         CALL     OUTREG
         CALL     IXOUT          ;DO IX REG
         CALL     CRLF          ;NOW ALTERNATE

```

```

        CALL    OUTREG
        RET

IXOUT:
DOING: LD      A, (HL)
        PUSH   AF          ; SAVE LOW BYTE
        INC   HL
        LD      A, (HL)
        CALL   HEXD        ; OUT HIGH
        POP    AF          ; GET LOW
        CALL   HEXD        ; AND OUT
        INC   HL
        LD      A, ' '
        CALL   PRTCHR
        CALL   PRTCHR
        RET

CHREG: ; GETREG VALUE

        LD      A, ' '
        CALL   PRTCHR
        CALL   GETKEY
        CP     ALTER
        JR     NZ, LOD
        CALL   GETKEY

        ADD    A, ' I' -' A'

LOD:   LD      HL, LOOKUP
        LD      BC, LENTAB
        CPIR
        JP     NZ, DISPR
        DEC   HL
        LD      DE, LOOKUP
        AND   A
        SBC   HL, DE
        LD      DE, AFREG
        SLA  L
        ADD   HL, DE
        INC  HL

```

```

LD      A, ' '
CALL   PRTCHR
LD      A, (HL)           ; LOW
CALL   HEXO
DEC     HL
LD      A, (HL)           ; HIGH
CALL   HEXO
INC     HL
LD      A, ' '
CALL   PRTCHR

CALL   HEXI
LD      (HL), A
DEC     HL
CALL   HEXI
LD      (HL), A
RET

```

LOOKUP:

```

DB 'A', 'B', 'D', 'H', 'X'
DB 'I', 'J', 'L', 'P'
;      A' B' D' H'

```

LENTAB EQU 9D

```

MOVE:  LD      A, ' '
        CALL   PRTCHR
        CALL   GETEXPR3

```

```

GETBC: PUSH   HL
        AND    A
        SBC   HL, DE
        JR    NC, DSWOP
        POP   HL
        EX   DE, HL
        JR   GETBC

```

DSWOP:

```

; HL =NUMBER OF BYTES
; DE =START ADDRESS
; BC =DESTINATION
; (SP)=END ADDRESS
EX   DE, HL
; DE =NUM HL=START
; BC=DEST

```

```

PUSH BC ;SAVE DEST
PUSH DE ;SAVE COUNT
POP BC ;PUT IN DE
POP DE ;GET DEST

POP AF ;GET RID OF END
;STACK CONTAINS START
PUSH HL ;GET START
;HL CONTAINS START
;DE DESTINATION
;BC NUMBER OF BYTES
;STACK CONTAINS

AND A
SBC HL, DE
POP HL
JR C, BACKW
LDIR
RET

```

BACKW:

```

ADD HL, BC
DEC HL
EX DE, HL
ADD HL, BC
DEC HL
EX DE, HL
LDDR
RET

```

IDENT:

```

LD A, ' '
CALL PRTCHR
CALL GETKEY
CP 0D
RET Z
CP 65
JP C, ERRORS
LD HL, FILENAME
LD B, 10
LD C, 32
CLBUFF: LD (HL), C
INC HL
DJNZ CLBUFF

```

```

LD      HL, FILENAME
LD      B, 9

PUTBUF: LD      ( HL ), A
DEC     B
RET     Z

INC     HL
CALL   GETKEY

CP      0DH
RET     Z

CP      65
JP      C, ERRORS
JR      PUTBUF

HUNT:

LD      A, ' '
CALL   PRTCHR
CALL   GETEXPR2
PUSH   HL          ; SAVE START
EX     DE, HL
AND    A
SBC   HL, DE
JP    C, ERRORS
JP    Z, ERRORS
PUSH  HL
POP   BC
POP   HL
LD    A, ' '
CALL  PRTCHR

PUSH  HL
PUSH  DE
CALL  HEXI
POP   DE
POP   HL

```

```

COMP:  CP      ( HL )
       PUSH   AF
       JR     NZ, NFOUND
       CALL  CRLF
       CALL  HEXDD
       CALL  GETKEY
       CP    0DH
       JR    NZ, BHUN

NFOUND:
       INC   HL
       DEC   BC
       LD   A, B
       OR   C
       JR   Z, BHUN
       POP  AF
       JR   COMP

BHUN:
       POP  AF
       RET

CRLF:  LD   A, CR
       CALL PRTCHR
       RET

REGMESS:
       DEFM ' AF   BC   '
       ;    0000: :0000: :
       DEFM ' DE   HL   '
       DEFM ' IX'
       DB   CR, '$'

REGS:
AFREG: DEFW 0000H ; AF
BCREG: DEFW 0000H ; BC
DEREG: DEFW 0000H ; DE
HLREG: DEFW 0000H ; HL
IXREG: DEFW 0000H ; IX
AAFREG: DEFW 0000H ; AF'
ABCREG: DEFW 0000H ; BC'
ADEREG: DEFW 0000H ; DE'
AHLREG: DEFW 0000H ; HL'

BRKP:  DB    0, 0, 0
SAVESP: DB   0, 0

```

```
HEADER: DS      17
FILENAME:
          DS      10
          DB      CR, '$'

          DS      75
STACK:   DB      0
ERRSP:   DEFW    0

          END
```



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ISBN 0-330-28665-X



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